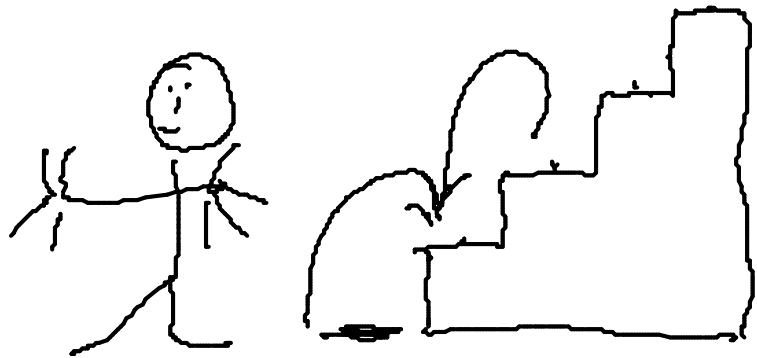


ASF

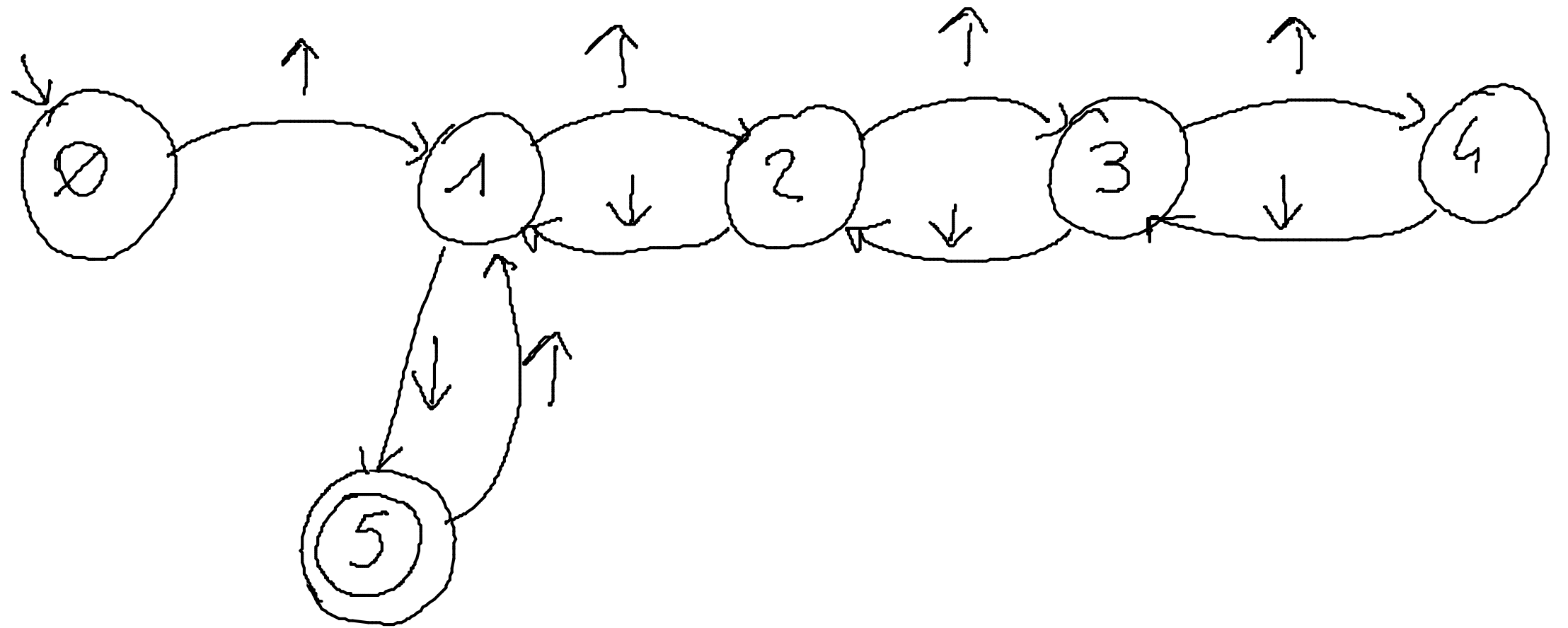
$\uparrow\downarrow \in L$ $\uparrow\uparrow\downarrow\uparrow\downarrow\downarrow \in L$

$\uparrow\uparrow\downarrow\downarrow\downarrow \notin L$

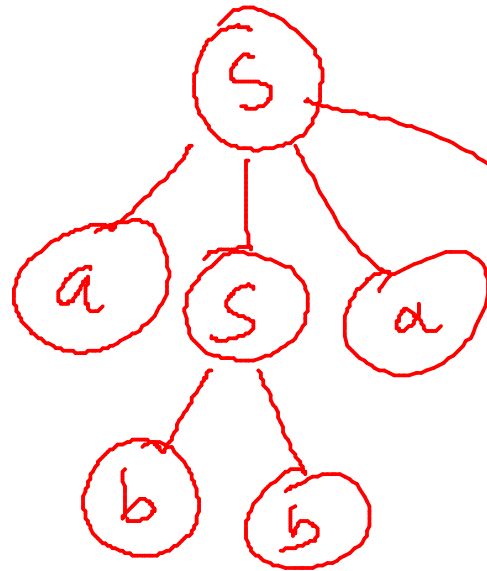
$\uparrow\downarrow\uparrow\downarrow \in L$



$\Sigma = \{\uparrow, \downarrow\}$



$\Sigma = \{a, b\}$



S palindroma

abba

aba

abaa

bab

a, b

aa

bb

palindroma

//

non //

palindroma

//

//

//

G tale che

L_G sia il linguaggio delle stringhe

palindrome

in Σ

non vuote

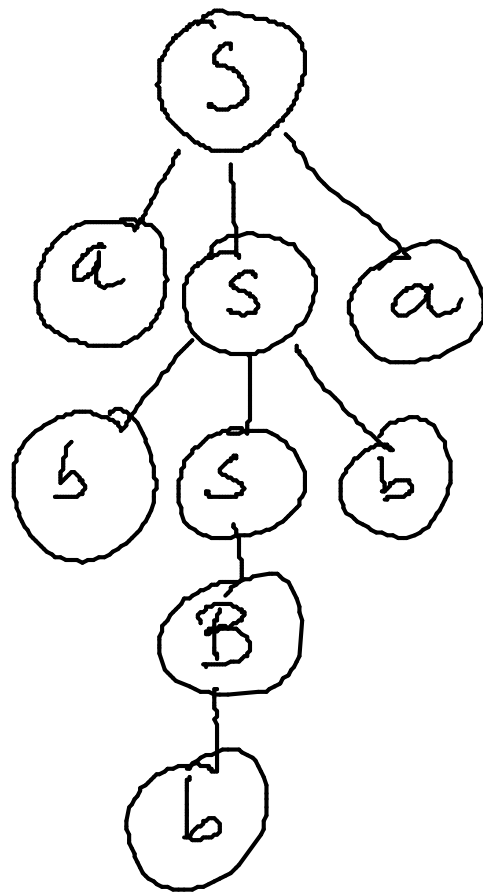
$S \rightarrow a | b | aa | bb | aSa | bSb$

$S \rightarrow aSa \mid bSb \mid A \mid B$

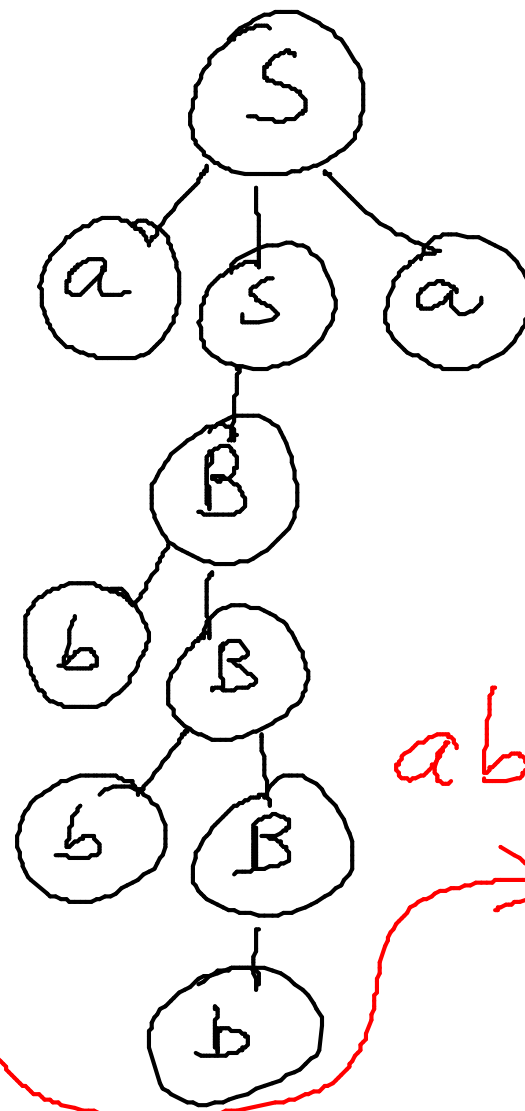
$A \rightarrow aA \mid a$

$B \rightarrow bB \mid b$

abbba



abbba



abbba

$$\Sigma = \{a, b, c\}$$

$$\alpha \in \Sigma^*$$

$$|\alpha|_{-c}$$

α dove tutte le c vengono cancellate

$$|abbcabc|_{-c} = abbab$$

$$L = \{ \alpha \mid \alpha \in \Sigma^+ \text{ e } |\alpha|_{-c} \text{ è palindroma} \}$$

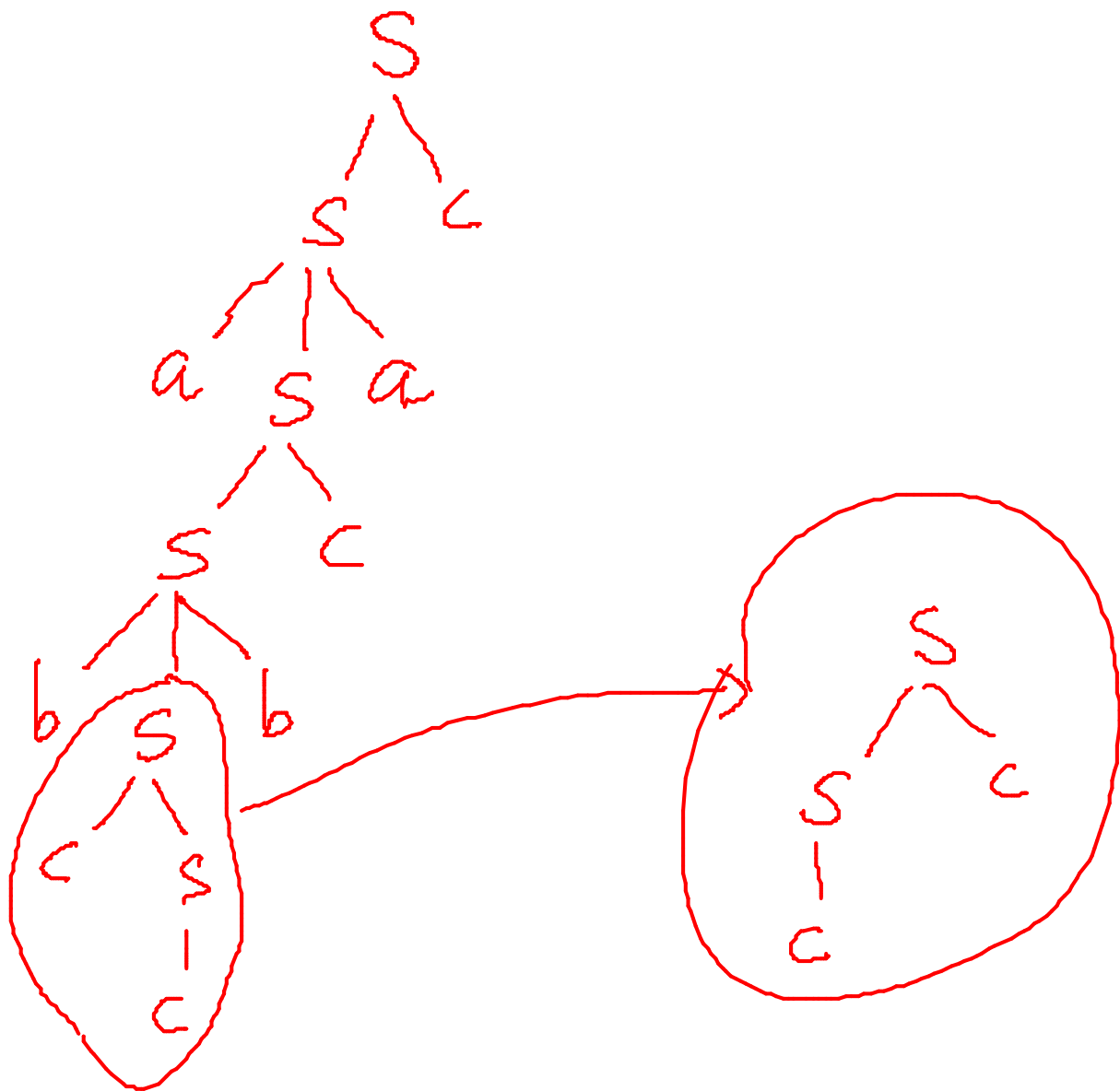
$$abccbca \in L$$

$$abaccbc \notin L$$

$$ccc \in L$$

$S \rightarrow a | b | aa | bb | aSa | bSb | cS | Sc | c$

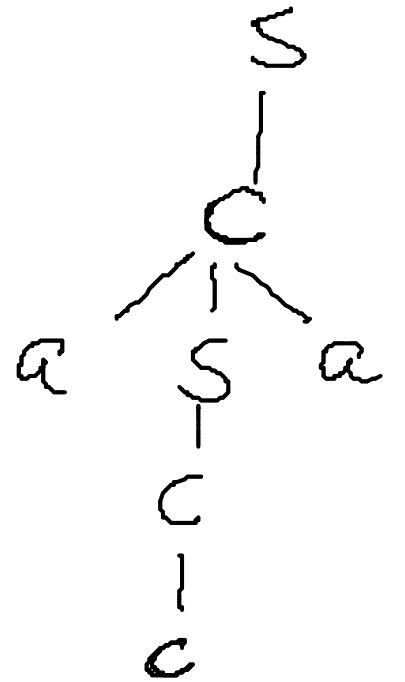
abc**cc**baac



$S \rightarrow a | b | aa | bb | c$

$C \rightarrow c | aSa | bSb$

acbb a



NO!

$$S \rightarrow \epsilon \quad S^i!$$

$$S \rightarrow CaC \mid CbC \mid CaCaC \mid CbCbC \mid \\ CaCSCaC \mid CbCSbC \mid cC$$

$$C \rightarrow \epsilon \mid c \mid cC$$

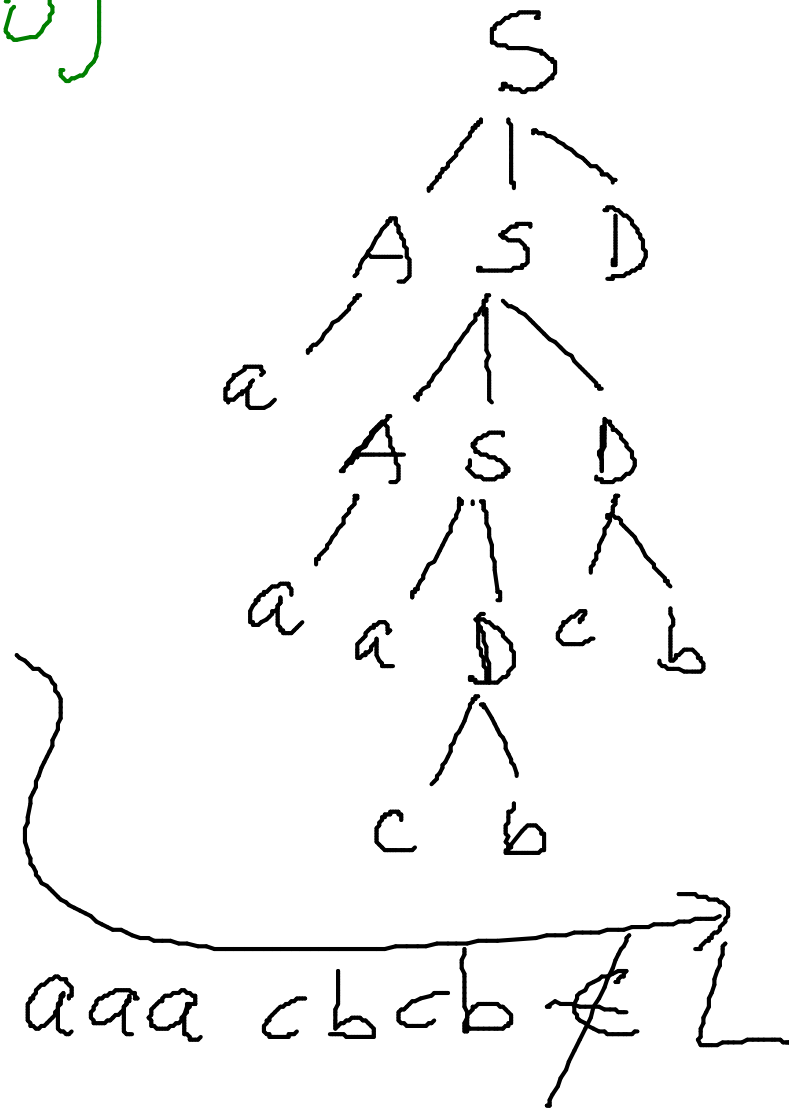
$$S \rightarrow aSa \mid bSb \mid A \mid B \mid c \mid cS \mid Sc$$

$$A \rightarrow aA \mid a$$

$$B \rightarrow bB \mid b$$

$$L = \{a^k c^m b^m \mid k, m > 0\}$$

$S \rightarrow AD \mid \textcircled{ASD}$
 $A \rightarrow a \mid aA$
 $D \rightarrow cb \mid cDb$ ~~NO!~~
 SII



Dimostrare che $L = \{ a^k c^m b^m \mid m, k > 0 \}$
non è regolare. (pumping lemma)

Se L è regolare allora

$\exists m$ tale che tutte le stringhe w
 $|w| \geq m$ possono essere divise in tre
parti $w = xyz$ con le proprietà

- $y \neq \epsilon$
- $|xy| \leq m$
- $xy^iz \in L$ per ogni $i \in \mathbb{N}$

L reg \Rightarrow

valgono le
proprietà

$L_{reg} \Rightarrow$ Valgono le proprietà

$A \Rightarrow B \equiv \neg B \Rightarrow \neg A$

Qualunque $n \in \mathbb{N}$ si trova
una stringa più lunga di n
tale che le 3 proprietà non valgono
contemporaneamente

qualunque $n \geq k$

$$a_1 \dots a_{k-m} c_1 \dots c_m b_1 \dots b_m$$

$$a^{k-m} c^m b^m$$

$$\left| a c^{\frac{m}{2}} b^{\frac{m}{2}} \right| \geq m$$

- $y \neq \varepsilon$ $v \neq \varepsilon$
- $|xy| \leq m$ $|uv| \leq m$

$$x = \varepsilon \quad \left. \begin{array}{l} y = a \\ y = a c \dots c \\ y = a c^{\frac{m}{2}} b \dots b \end{array} \right|$$

$$x = a \quad \left. \begin{array}{l} y = c \dots c \\ y = c^{\frac{m}{2}} b \dots b \end{array} \right|$$

k

$$|a^k b^k| \leq k$$

$x y z$

\uparrow
 $b \cdot b$

$$|xy| \leq k$$