

Stein (discrete Lyapunov) equation

Note Title

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$$\rightarrow X - A^T X A = Q$$

$$X - A X A^T = G, \quad G = b b^T$$

$$x_{k+1} = A x_k + b u_k$$

$$A \in \mathbb{C}^{n \times n} \quad x_k, b \in \mathbb{C}^n$$

$$X = \sum_{k=0}^{\infty} (A^T)^k Q A^k \quad \rho(A) < 1$$

Smith method:

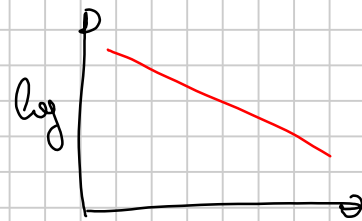
$$X_{k+1} = Q + A^T X_k A \quad X_0 = 0$$

$$X = Q - A^T X A$$

$$X_{k+1} - X = A^T (X_k - X) A$$

$$X_k - X = (A^T)^k (X_0 - X) A^k$$

$$\|X_k - X\| \sim \rho(A)^{2k}$$



A faster: "doubling variant"

Idea: construct a recurrence for $\boxed{Q_k = X_{2^k}}$

$$X_k = Q + A^T Q A + (A^T)^2 Q A^2 + \dots + (A^T)^{k-1} Q A^{k-1}$$

$$\text{vec}(X_k) = (I + M + M^2 + \dots + M^{k-1}) \text{vec } Q \quad M = A^T \otimes A^T$$

$$I + M + M^2 + \dots + M^{2^{k+1}-1} = (I + M + M^2 + \dots + M^{2^k-1}) + M^{2^k} (I + M + M^2 + \dots + M^{2^k-1})$$

$$\text{vec } Q_{k+1} = \text{vec } Q_k + M^{2^k} \text{vec } Q_k$$

$$Q_{k+1} = Q_k + (A^T)^{2^k} Q_k A^{2^k}$$

$$\begin{cases} A_0 = A, Q_0 = Q \\ A_{k+1} = A_k^2 \\ Q_{k+1} = Q_k + A_k^* Q_k A_k \end{cases}$$

$$Q_k = X_{2^k} \text{ at each step}$$

Cost: 3 mult/mol

Convergence: $\|X - Q_k\| \sim \rho(A)^{2 \cdot 2^k}$

"Squared Smith"

$$X_0 \preceq X_1 \preceq X_2 \preceq X_3 \preceq \dots \preceq X$$

Sparse A $X - A X A^* = G = \begin{bmatrix} b & b^* \end{bmatrix}$

problem: X dense, and (in general) not low-rank

"solution": in many problems, X is approximately low rank!

We want to store the solution as $X \approx Z Z^*$, Z tall thin

Smith:

$$X_k = b b^* + A b b^* A^* + A^2 b b^* (A^*)^2 + \dots + A^{k-1} b b^* (A^*)^{k-1}$$

$$X_k = Z_k Z_k^* \quad Z_k = \begin{bmatrix} b & A b & A^2 b & \dots & A^{k-1} b \end{bmatrix} \in \mathbb{C}^{n \times k}$$

$$v_0 = b \quad v_{k+1} = A v_k$$

$$Z_k = \begin{bmatrix} v_0 & \dots & v_{k-1} \end{bmatrix}$$

$$Z_1 = \begin{bmatrix} b \end{bmatrix} \quad Z_{k+1} = \begin{bmatrix} Z_k & v_k \end{bmatrix}$$

$$\|Z_k Z_k^* - X\| \sim \rho(A)^{2k}$$

$$\sigma_k(X) \sim \rho(A)^{2k}$$

$$\sigma_{k+1} = \min_{\text{rk } Y = k} \|X - Y\|$$

(Move on this algorithm with Lyapunov equations)

Compression of Z_k :

$$Z_k = \begin{bmatrix} \end{bmatrix}$$

$$Z_k = U_k S_k U_k^*$$

$\begin{bmatrix} \end{bmatrix} \triangle \square$

If any of $\text{diag}(S_k)$ are small, truncate U, S, V .

Remark: for the dense case, you do not want to solve

$$(I - M) \text{vec } X = \text{vec } Q \quad \text{with a direct } O(n^6) \text{ method.}$$

Lyapunov equations: (continuous-time equivalent of Smith)

$$Q \succ 0$$

$$A^* X + X A + Q = 0$$

$$\dot{x}(t) = A x(t)$$

$$x(t) = \exp(tA) x_0$$

$$V(x) = x^* X x$$

$$\frac{d}{dt} V(x(t)) = \frac{d}{dt} x(t)^* X x(t) = \dot{x}(t)^* X x(t) + x(t)^* X \dot{x}(t)$$

$$= x(t)^* (A^* X + X A) x(t) = -x(t)^* Q x(t) \leq 0$$

table:

cont. time

$$\dot{x}(t) = A x(t)$$

stable if $\lambda(A) \subset \text{LHP}$
left half-plane

$$x(t) = \exp(At) x_0$$

discr. time

$$x_{k+1} = A x_k$$

stable if $\lambda(A) \subset \text{disc}$

$$x_k = A^k x_0$$

Lyap

$$A^* X + X A + Q = 0$$

$$X = \int_0^\infty \exp(A^* t) Q \exp(At) dt$$

Stein

$$X - A^* X A = Q$$

$$X = \sum_{k=0}^\infty (A^*)^k Q A^k$$

(if A stable)

$$\text{Control: } \dot{x}(t) = Ax(t) + bu(t)$$

$$x(t) = \exp(At)x_0 + \int_0^t \exp(A(t-\tau))bu(\tau)d\tau$$

$$\text{since } \exp(A(t-\tau))b \in \text{Im}[b \quad Ab \quad A^2b \dots]$$

controllable if

$$\text{rank}[b \quad Ab \quad A^2b \dots] = n$$

(if A stable)

$$\text{Control: } x_{k+1} = Ax_k + bu_k$$

$$x_k = A^k x_0 + A^{k-1}bu_0 + A^{k-2}bu_1 + \dots + bu_{k-1}$$

controllable if

$$\text{rank}[b \quad Ab \quad \dots \quad A^{k-1}b] = n$$

def: A system is controllable if for given $x_F \in \mathbb{C}^n$ and t_F , I can choose $u(t)$ such that

$$x(t_F) = x_F$$

$$\min \|u\| \text{ s.t. } x(t_F) = x_F \text{ for some } t_F \text{ (including } t_F = \infty) = x_F^* X^{-1} x_F$$

$$\text{where } X \text{ solves } AX + XA^* + bb^* = 0$$

$$\min \|u\| \text{ s.t. } x_k = x_F \text{ for some } k = x_F^* X^{-1} x_F$$

$$\text{where } X \text{ solves } X - AXA^* = bb^*$$

"Controllability Gramian"

$$t_F = \infty, \quad x_F^* X_{t_F}^{-1} x_F, \text{ where}$$

$$X_{t_F} = \int_0^{t_F} \exp(At) bb^* \exp(A^*t) dt$$

$$A^*X + XA$$

$$\text{vec } X = (I \otimes A^* + A^T \otimes I)^{-1} \text{vec } Q$$

$$A = U \Lambda U^*$$

$$(U \otimes \bar{U}) (I \otimes T^* + T \otimes I) (U \otimes U)$$

orth  orth

$$\begin{bmatrix} b & Ab & \dots & A^{k-1}b \end{bmatrix} \begin{bmatrix} b^* \\ b^* A^* \\ \vdots \\ b^* A^{*k} \end{bmatrix}$$

$$\text{vec } X = (I - A^T \otimes A^*)^{-1} \text{vec } Q$$

Singular if there are two eigenvalues

$$A_i, A_j \in \Lambda(A) \text{ s.t.}$$

$$\lambda_i = -\overline{\lambda_j}$$

Singular if there are

two eigenvalues $\lambda_i, \lambda_j \in \Lambda(A)$

$$s.t. \quad \lambda_i = \frac{1}{\lambda_j}$$

$$I \otimes T^* + T^T \otimes I$$

The diagram shows two parallel paths within large square brackets. The top path contains three triangular gates (each with a diagonal line) and a circle. The bottom path contains a series of gates labeled with indices: γ_1 , γ_2 , γ_3 , γ_4 , γ_5 , γ_6 , γ_7 , γ_8 , γ_9 , γ_{10} , γ_{11} , γ_{12} , γ_{13} , γ_{14} , γ_{15} , γ_{16} , γ_{17} , γ_{18} , γ_{19} , γ_{20} , γ_{21} , γ_{22} , γ_{23} , γ_{24} , γ_{25} , γ_{26} , γ_{27} , γ_{28} , γ_{29} , γ_{30} , γ_{31} , γ_{32} , γ_{33} , γ_{34} , γ_{35} , γ_{36} , γ_{37} , γ_{38} , γ_{39} , γ_{40} , γ_{41} , γ_{42} , γ_{43} , γ_{44} , γ_{45} , γ_{46} , γ_{47} , γ_{48} , γ_{49} , γ_{50} , γ_{51} , γ_{52} , γ_{53} , γ_{54} , γ_{55} , γ_{56} , γ_{57} , γ_{58} , γ_{59} , γ_{60} , γ_{61} , γ_{62} , γ_{63} , γ_{64} , γ_{65} , γ_{66} , γ_{67} , γ_{68} , γ_{69} , γ_{70} , γ_{71} , γ_{72} , γ_{73} , γ_{74} , γ_{75} , γ_{76} , γ_{77} , γ_{78} , γ_{79} , γ_{80} , γ_{81} , γ_{82} , γ_{83} , γ_{84} , γ_{85} , γ_{86} , γ_{87} , γ_{88} , γ_{89} , γ_{90} , γ_{91} , γ_{92} , γ_{93} , γ_{94} , γ_{95} , γ_{96} , γ_{97} , γ_{98} , γ_{99} . The paths are connected by a plus sign and a dashed line.

Bartels-Stewart algorithm (1970s)

$O(n^3)$ direct algorithm to solve $A^*X + XA + Q = 0$:

1. Compute $A = U T U^*$, Schur factorization

2. Change of variables in the equation:

$$\cancel{U^*} U^T \cancel{U^*} X U + U^* X U T \cancel{U^*} + U^* Q U = 0$$

$$T^* \hat{X} + \hat{X} T + \hat{Q} = 0$$

$$\hat{X} = U^* X U$$

$$\hat{Q} = U^* Q U$$

3. Solve the triangular version entry by entry:

$$\begin{bmatrix} \bar{x} \\ \textcircled{x} & x \\ x & x & x \\ x & x & x & x \end{bmatrix} \begin{bmatrix} x & x & x & x \\ x & x & x & x \\ x & x & x & x \\ x & x & x & x \end{bmatrix} + \begin{bmatrix} x & x & x & x \\ \textcircled{x} & x & x & x \\ x & x & x & x \\ x & x & x & x \end{bmatrix} \begin{bmatrix} x & x & x & x \\ x & x & x & x \\ x & x & x & x \\ x & x & x & x \end{bmatrix} + \begin{bmatrix} x & x & x & x \\ x & x & x & x \\ x & x & x & x \\ \textcircled{x} & x & x & x \end{bmatrix} = 0$$

T^* $\hat{X}$ $\hat{X}$ T $\hat{Q}$

(1,1) entry:

$$\bar{t}_{11} x_{11} + x_{11} t_{11} + q_{11} = 0 \rightarrow x_{11} = \frac{-q_{11}}{\bar{t}_{11} + t_{11}}$$

$$(2,1) \quad \bar{t}_{21} x_{11} + \bar{t}_{22} x_{21} + x_{21} t_{11} + q_{21} = 0 \rightarrow x_{21} = \frac{-q_{21} - \bar{t}_{21} x_{11}}{\bar{t}_{22} + t_{11}}$$

$$(i,j) \quad \sum_{k=1}^i \bar{t}_{ik} x_{kj} + \sum_{k=1}^j x_{ik} t_{kj} + q_{ij} = 0$$

$$x_{ij} = \frac{-q_{ij} - \sum_{k=1}^{i-1} \bar{t}_{ik} x_{kj} - \sum_{k=1}^{j-1} x_{ik} t_{kj}}{\bar{t}_{ii} + t_{jj}}$$



We can compute $\hat{X}$ entry by entry

$$\hookrightarrow \boxed{X = U \hat{X} U^*}$$

Idea: this is inverting $I \otimes A^* + A^T \otimes I$ using the Schur factorization:

$$(I \otimes A^* + A^T \otimes I)^{-1} \text{vec } Q = (\bar{U} \otimes U) (I \otimes T^* + T^T \otimes I)^{-1} (\underbrace{U^T \otimes U^*}_{\text{vec}(\hat{Q})}) \text{vec } Q$$

forward substitution to compute $\text{vec}(\hat{X})$

$$\text{vec}(X) = \text{vec}(U \hat{X} U^*)$$

$$O(n^3) \text{ beats } O(n^6)$$

Same idea for Stein equations:

$$\hat{X} - T^* \hat{X} T = \hat{Q}$$

$$\square - \triangle \square \nabla = \square$$

entries can be computed one by one, because $I - T^T \otimes T^*$ is lower triangular.

The same idea works for $AX + XB = C$ (Sylvester equations)

and for

$$AXB + CXD = E \quad (\text{generalized Sylvester equations})$$

with a generalization of the Schur factorization $(q_2(A, C), q_2(B, D))$

The Bartels-Stewart idea does not work for 3-term equations:

$$AX + XB + DXE = C$$

(no cubic algorithms known for those)

$$\begin{cases} A_1^* X + X A_1 \preceq 0 \\ A_2^* X + X A_2 \preceq 0 \end{cases} \quad V(x) = x^* X x$$