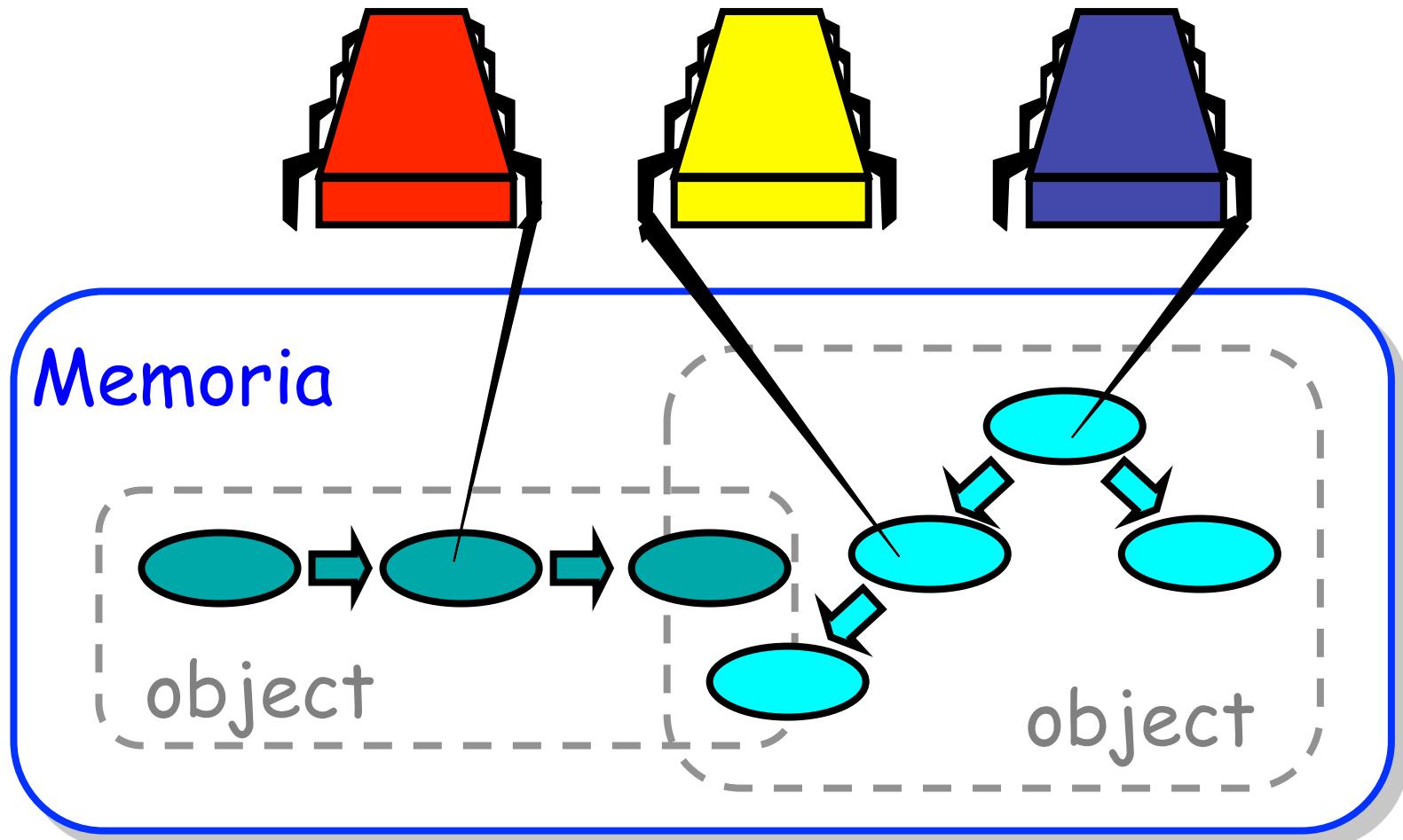


# Concurrent Objects

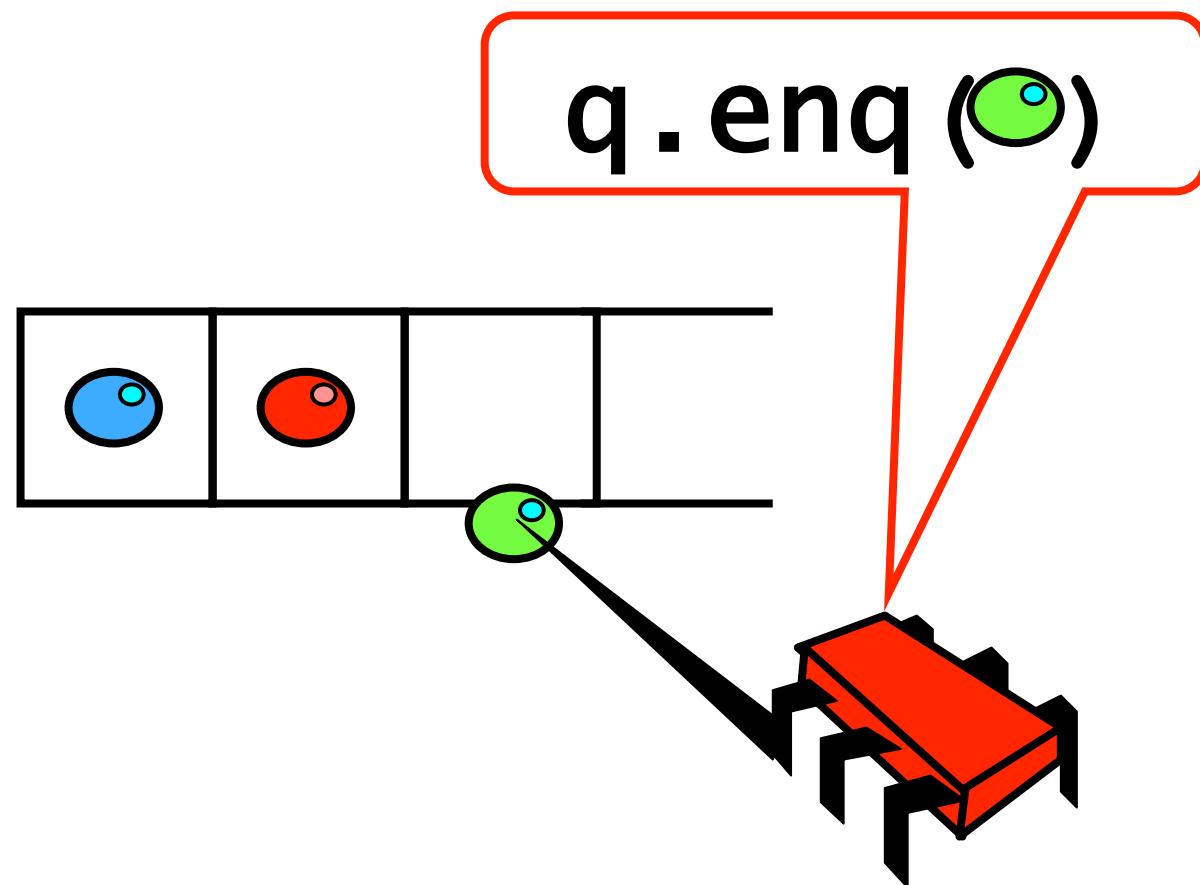
# La concorrenza



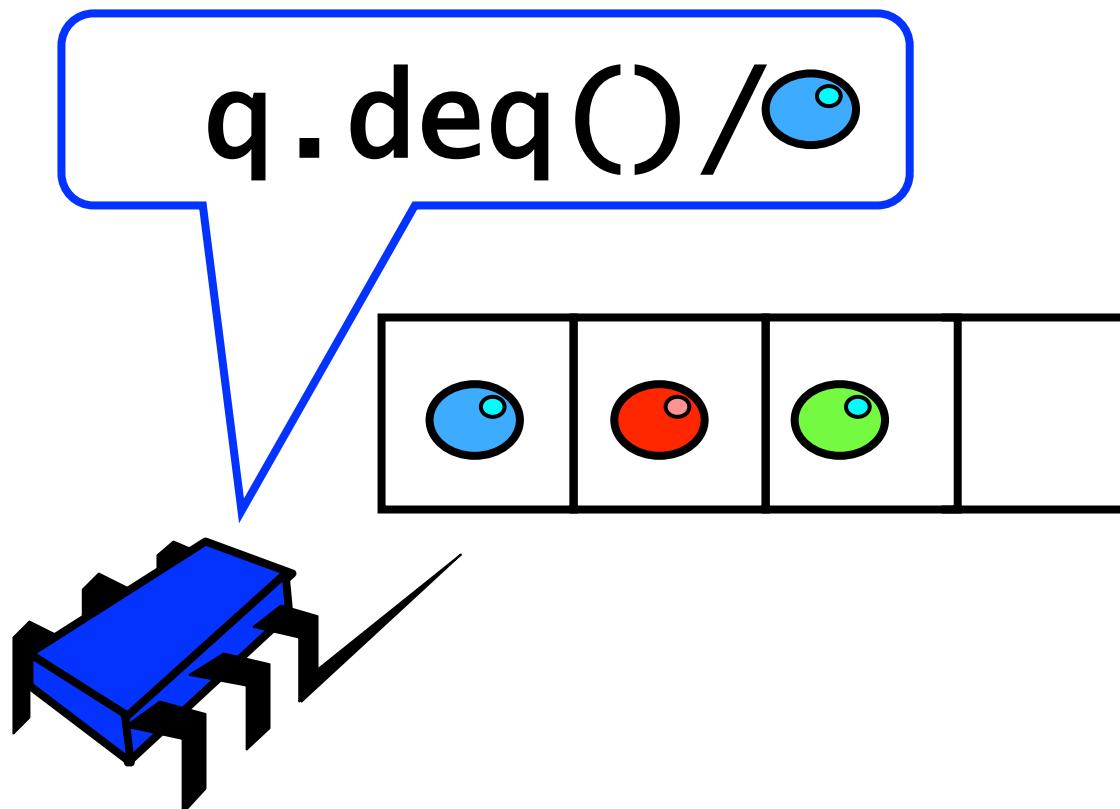
# Il problema che vogliamo affrontare

- Cosa e' un oggetto concorrente?
  - In quale modo lo descriviamo?
  - In quale modo lo implementiamo?
  - In quale modo dimostriamo la correttezza?

# FIFO Queue: Enqueue Method



# FIFO Queue: Dequeue Method

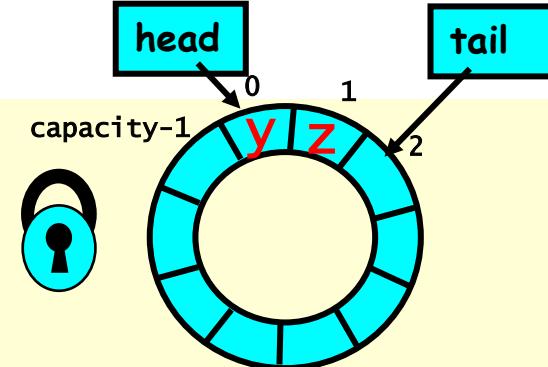


# Lock-Based Queue

```
class LockBasedQueue<T> {  
    int head, tail;  
    T[] items;  
    Lock lock;  
    public LockBasedQueue(int capacity) {  
        head = 0; tail = 0;  
        lock = new ReentrantLock();  
        items = (T[]) new Object[capacity];  
    }  
}
```

# Lock-Based Queue

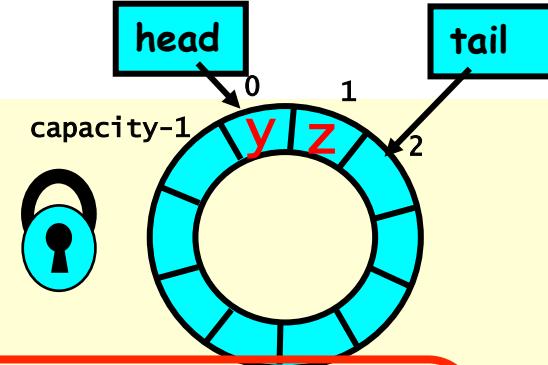
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class LockBasedQueue<T> {  
    int head, tail;  
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    Lock lock;  
    public LockBasedQueue(int capacity) {  
        head = 0; tail = 0;  
        lock = new ReentrantLock();  
        items = (T[]) new Object[capacity];  
    }
```



Queue fields  
protected by single  
shared lock

# Lock-Based Queue

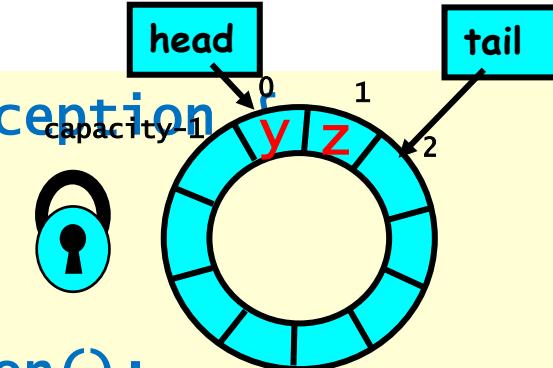
```
class LockBasedQueue<T> {  
    int head, tail;  
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    Lock lock;  
  
    public LockBasedQueue(int capacity) {  
        head = 0; tail = 0;  
        lock = new ReentrantLock();  
        items = (T[]) new Object[capacity];  
    }  
}
```



Initially `head = tail`

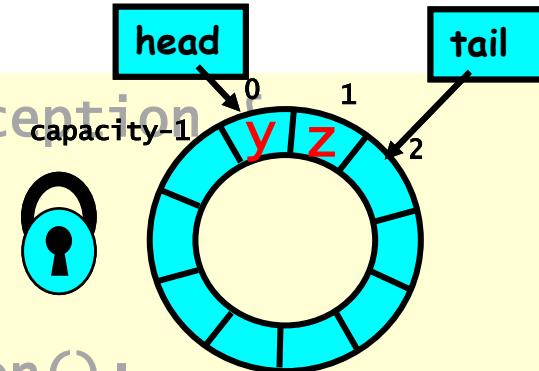
# Implementation: Deq

```
public T deq() throws EmptyException  
    lock.lock();  
    try {  
        if (tail == head)  
            throw new EmptyException();  
        T x = items[head % items.length];  
        head++;  
        return x;  
    } finally {  
        lock.unlock();  
    }  
}
```



# Implementation: Deq

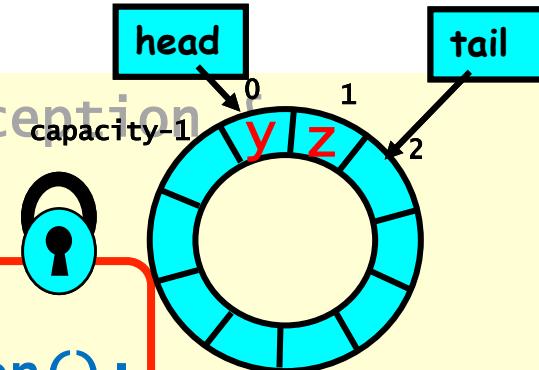
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        T x = items[head % items.length];  
        head++;  
        return x;  
    } finally {  
        lock.unlock();  
    }  
}
```



Method calls  
mutually exclusive

# Implementation: Deq

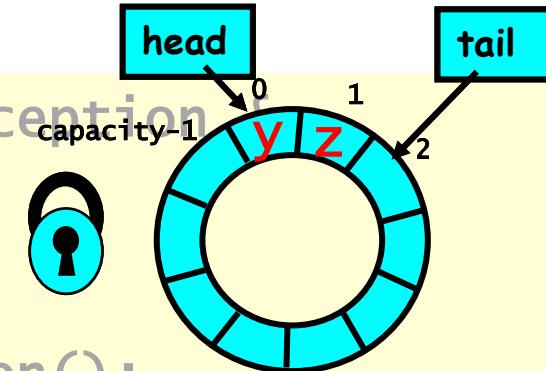
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            throw new EmptyException();  
        T x = items[head % items.length];  
        head++;  
        return x;  
    } finally {  
        lock.unlock();  
    }  
}
```



If queue empty  
throw exception

# Implementation: Deq

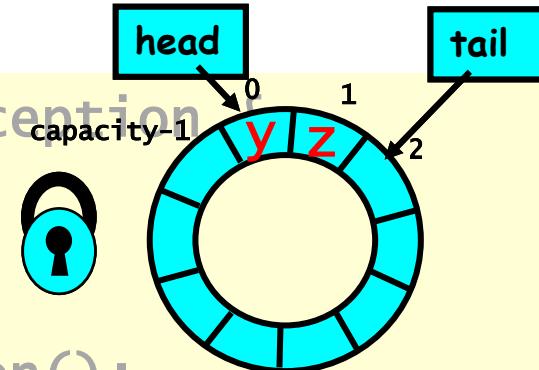
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    lock.lock();  
    try {  
        if (tail == head)  
            throw new EmptyException();  
        T x = items[head % items.length];  
        head++;  
        return x;  
    } finally {  
        lock.unlock();  
    }  
}
```



Queue not empty:  
remove item and update  
head

# Implementation: Deq

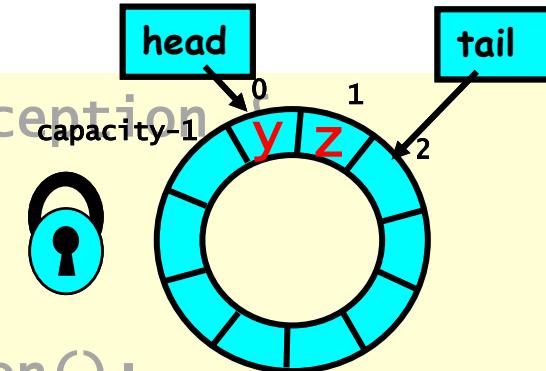
```
public T deq() throws EmptyException  
    lock.lock();  
    try {  
        if (tail == head)  
            throw new EmptyException();  
        T x = items[head % items.length];  
        head++;  
        return x;  
    } finally {  
        lock.unlock();  
    }  
}
```



Return result

# Implementation: Deq

```
public T deq() throws EmptyException  
    lock.lock();  
try {  
    if (tail == head)  
        throw new EmptyException();  
    T x = items[head % items.length];  
    head++;  
    return x;  
} finally {  
    lock.unlock();  
}  
}
```



Release lock no  
matter what!

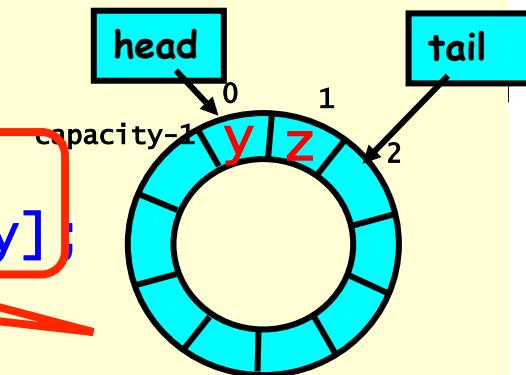
# Implementation: Deq

```
public T deq() throws EmptyException {  
    lock.lock();  
    try {  
        if (tail == head)  
            throw new EmptyException();  
        T x = items[head % items.length];  
        head++;  
        return x;  
    } finally {  
        lock.unlock();  
    }  
}
```

E' corretta? Le modifiche  
avvengono in mutua esclusione...

```
public class WaitFreeQueue {  
  
    int head = 0, tail = 0;  
    items = (T[]) new Object[capacity];  
  
    public void enq(Item x) {  
        if (tail-head == capacity) throw  
            new FullException();  
        items[tail % capacity] = x; tail++;  
    }  
    public Item deq() {  
        if (tail == head) throw  
            new EmptyException();  
        Item item = items[head % capacity]; head++;  
        return item;  
    }  
}
```

```
public class WaitFreeQueue {  
    int head = 0, tail = 0;  
    items = (T[]) new Object[capacity];  
  
    public void enq(Item x) {  
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        items[tail % capacity] = x; tail++;  
    }  
    public Item deq() {  
        if (tail == head) throw  
            new EmptyException();  
        Item item = items[head % capacity]; head++;  
        return item;  
    }  
}
```



```

public class WaitFreeQueue {
    int head = 0, tail = 0;
    items = (T[]) new Object[capacity];

    public void enq(Item x) {
        if (tail-head == capacity) throw
            new FullException();
        items[tail % capacity] = x; tail++;
    }
    public Item deq() {
        if (tail == head) throw
            new EmptyException();
        Queue is updated
        return item;
    }
}

```

The diagram shows a circular queue with 12 slots. The head pointer is at index 0, and the tail pointer is at index 2. Slots 0 and 1 are filled with 'y' and 'z' respectively. Slot 2 is empty. The label 'capacity-1' points to slot 1.

*Le modifiche non sono  
mutuamente esclusive!*

# Progresso

- In un contesto concorrente si deve specificare sia le proprie' di invarianza (safety) che le proprie' di liveness
- Correttezza
  - Quando una implementazione e' corretta
  - Le condizioni che garantiscono il progresso

# Oggetti sequenziali

- Ogni astrazione ha un proprio *state*
  - Le variabili di istanza *fields*
  - Queue: vettore di items
- Ogni astrazione possiede dei metodi *methods*
  - Descrivono come operare sullo statep
  - Queue: enq e deq

# Specifica

- (precondition)
  - Prima di invocare un metodo l'oggetto e' nello stato,
- (postcondition)
  - Il metodo modifica lo stato oppure solleva una eccezione.

# Dequeue

- Precondition:
  - Queue is non-empty
- Postcondition:
  - Returns first item in queue
- Postcondition:
  - Removes first item in queue

# Pre - Post Conditions Dequeue

- Precondition:
  - Queue is empty
- Postcondition:
  - Throws Empty exception
- Postcondition:
  - Queue state unchanged

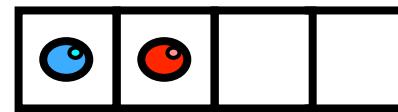
# Sequenziale

- Le interazioni tra i metodi sono catturate tramite effetti layerali sullo stato degli oggetti
  - Invariante di rappresentazione a questo serve!!!
- Ogni metodo è descritto singolarmente
- Refinement: possiamo aggiungere metodi senza modificare la descrizione dei vecchi metodi.

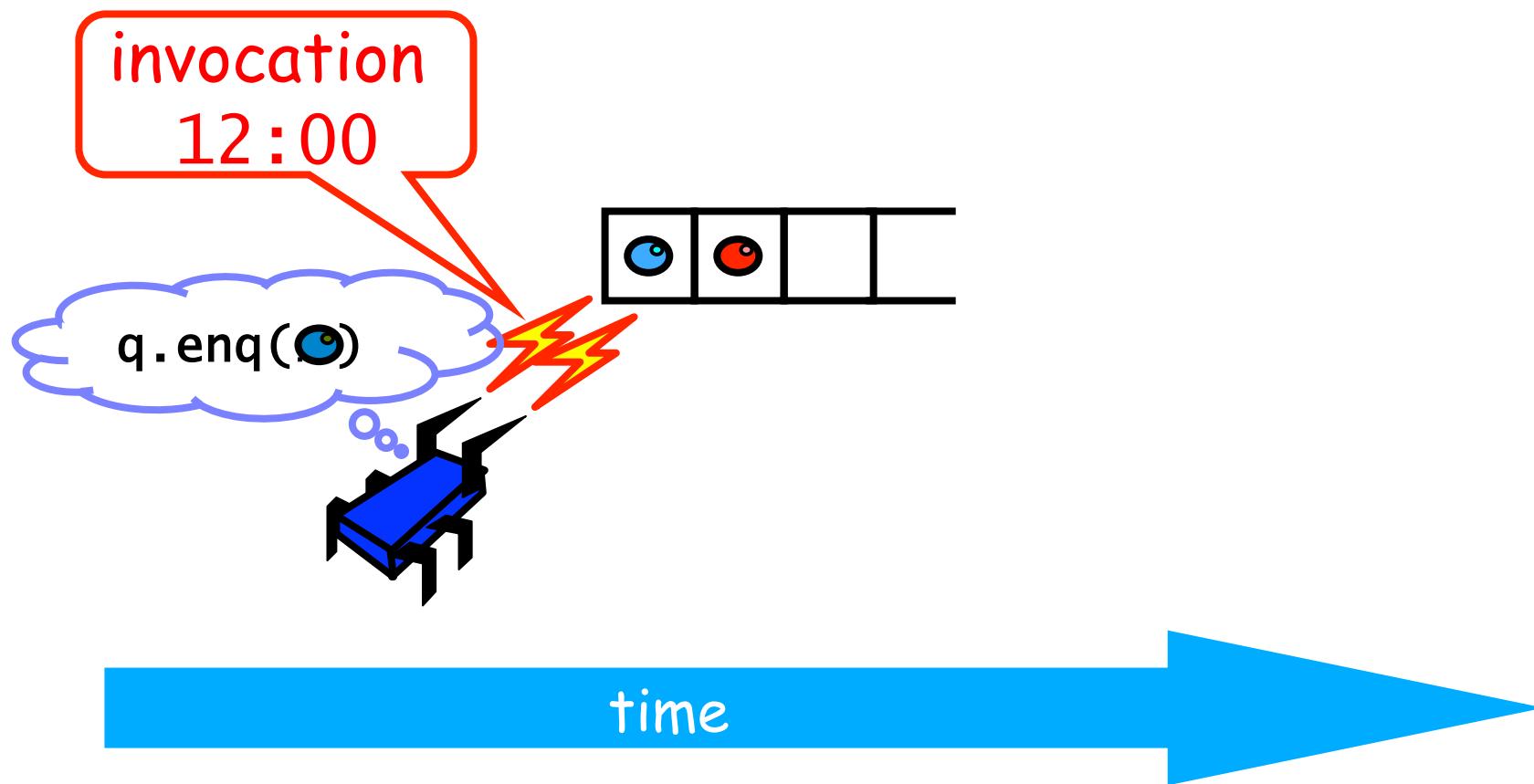
# Cosa cambia con la concorrenza?

- Metodi?
- La descrizione del metodo?
- Refinement?

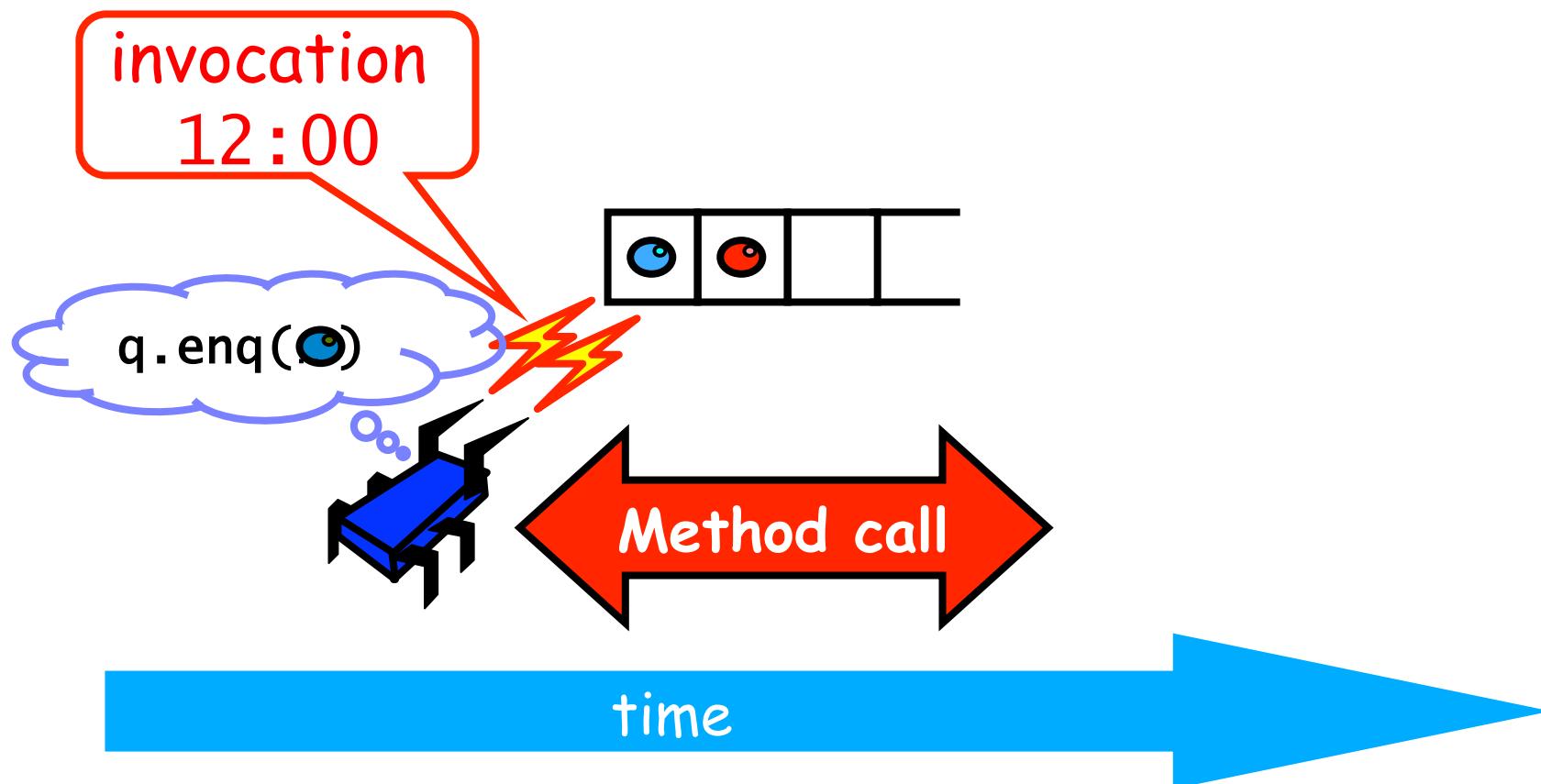
# "Methods Take Time"



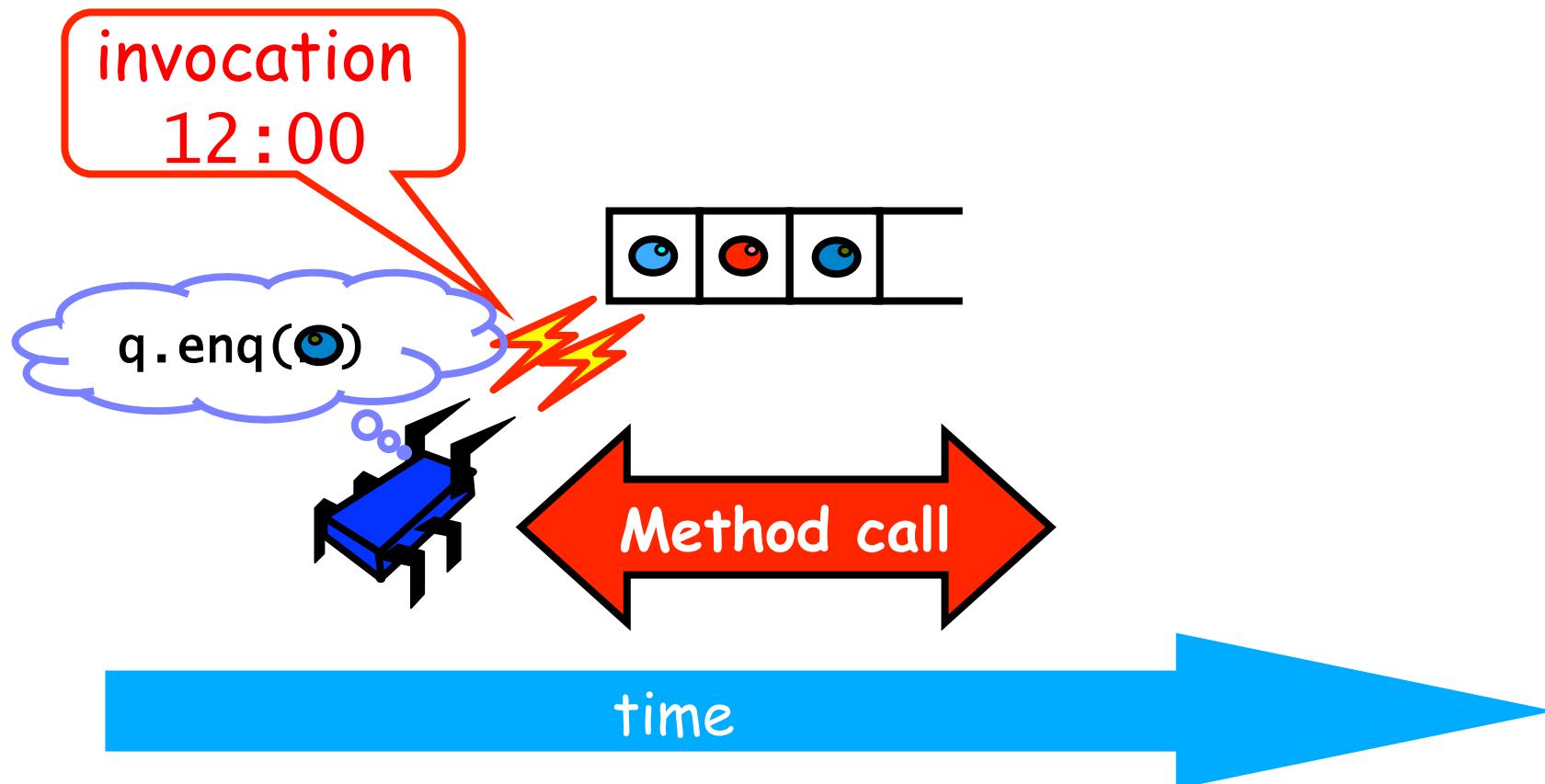
# Methods Take Time



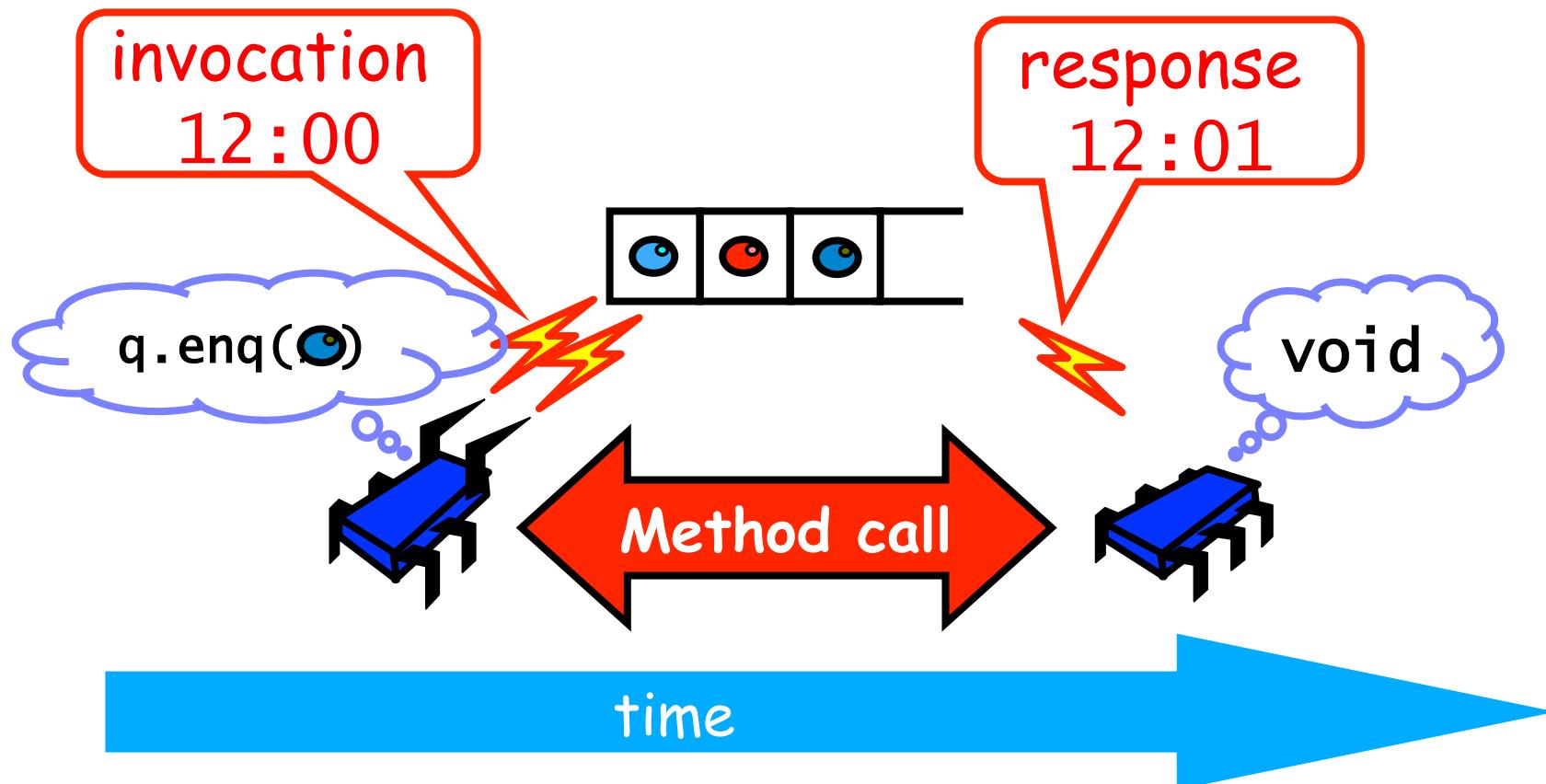
# Methods Take Time



# Methods Take Time



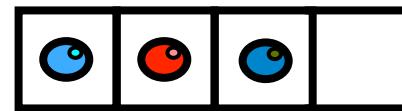
# Methods Take Time



# Sequenziale vs Concorrente

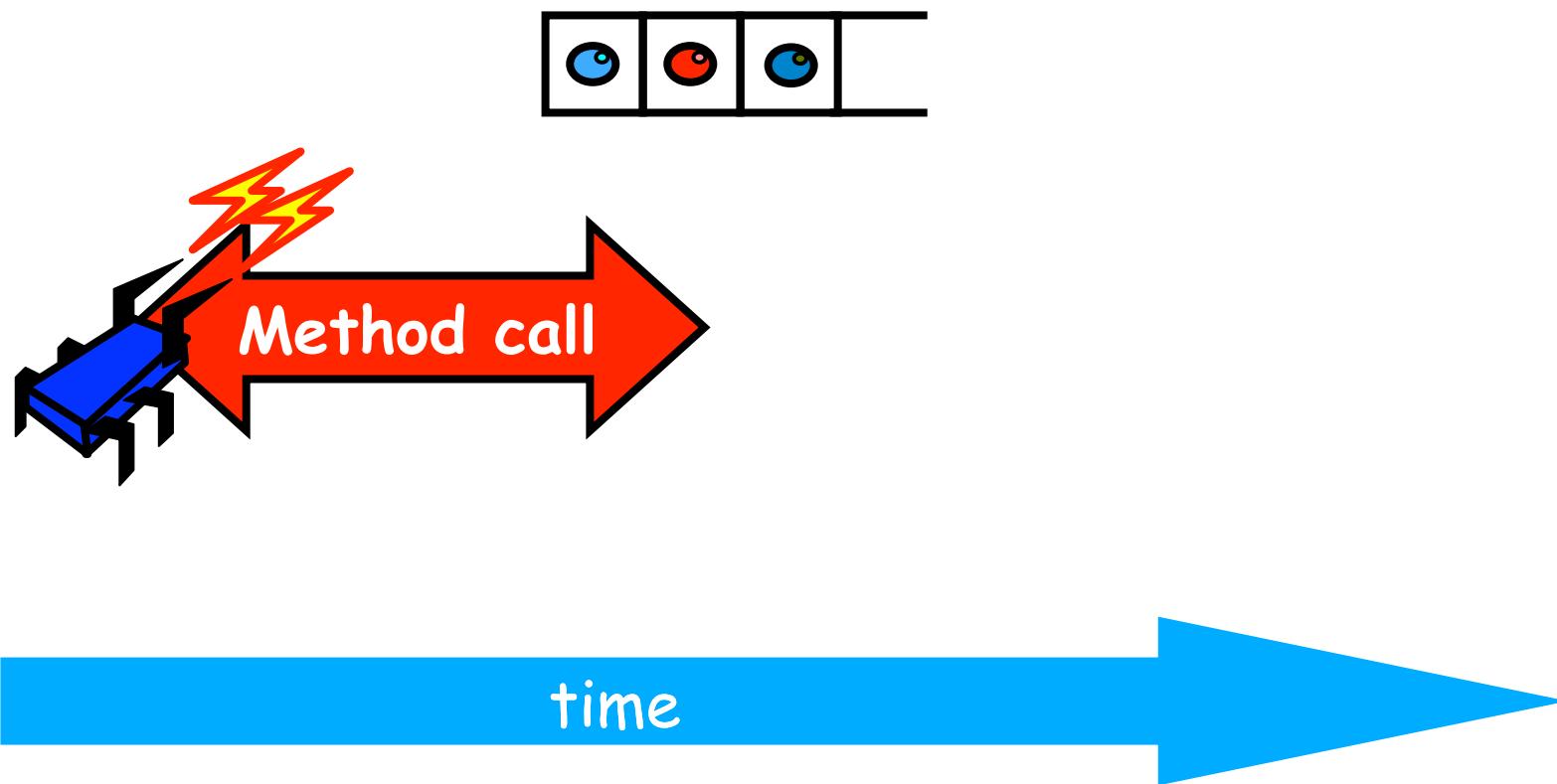
- Sequenziale
  - Methods take time?
  - Quando mai!!!
- Concorrente
  - La chiamata di un metodo e' un intervallo.

# Overlapping Time

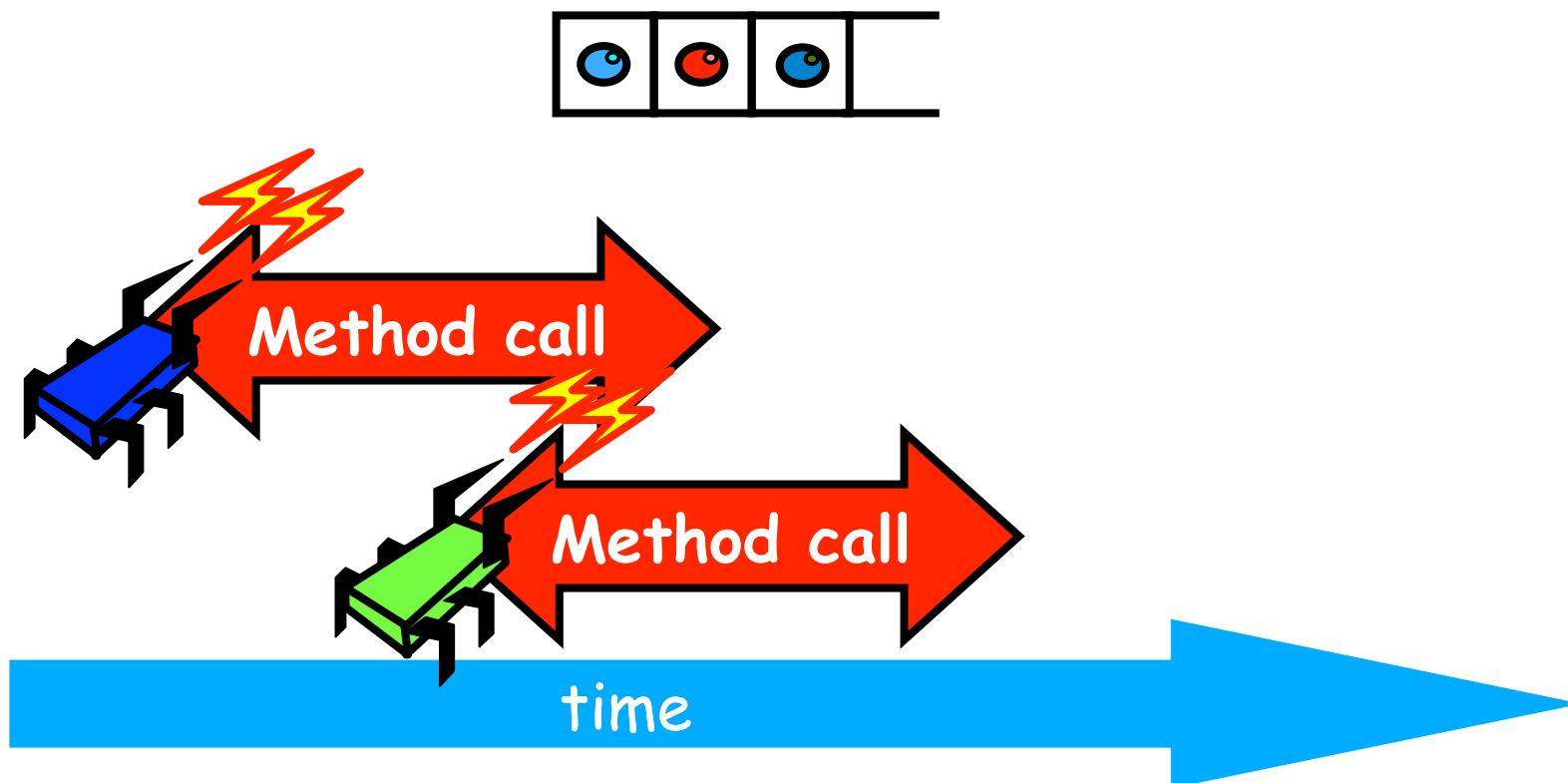


time

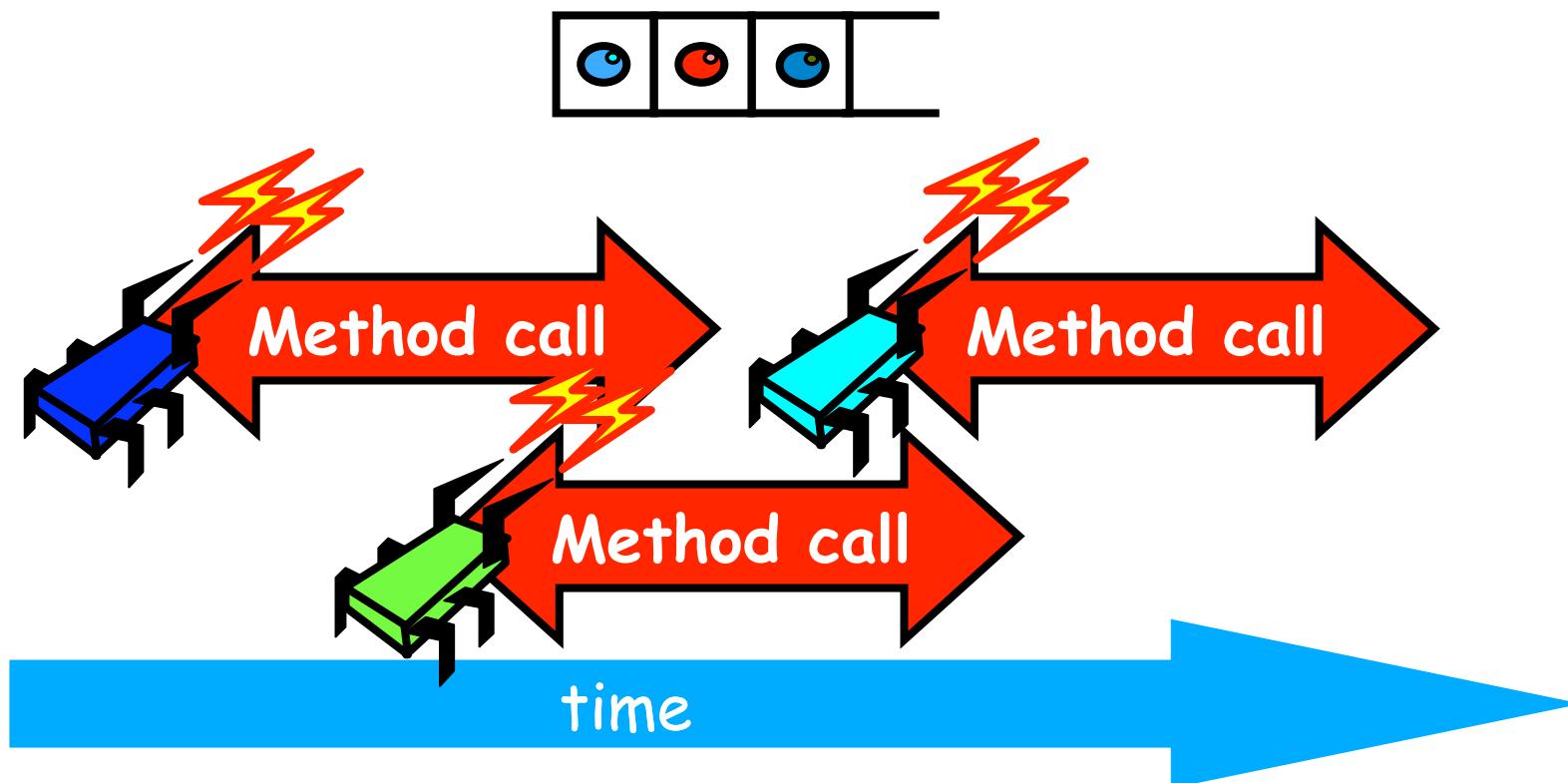
# Overlapping Time



# Overlapping Time



# Overlapping Time



# Sequenziale vs Concorrente

- Sequenziale:
  - Stato degli oggetti e' significativo solamente *tra* le invocazioni dei metodi
- Concorrente
  - Dato che le chiamate si sovrappongono la stato di un oggetto potrebbe *non essere mai* tra le invocazioni dei metodo

# Sequenziale vs Concorrente

- Sequenziale:
  - Ogni metodo e' descritto in isolamento
- Concorrente
  - Si devono comprendere *tutte* le possibili interazioni con chiamate concorrenti
    - Cosa succede se due invocazioni di enq si sovrappongono?

# Sequenziale vs Concorrente

- Sequenziale:
  - Refinement
- Concorrente:
  - Ogni metodo puo' potenzialmente interagire con tutti gli altri

# Sequenziale vs Concorrente

- Sequenziale:
  - Refinement
- Concorrente:
  - Ogni metodo puo' potenzialmente interagire con tutti gli altri

Panic!

# La domanda

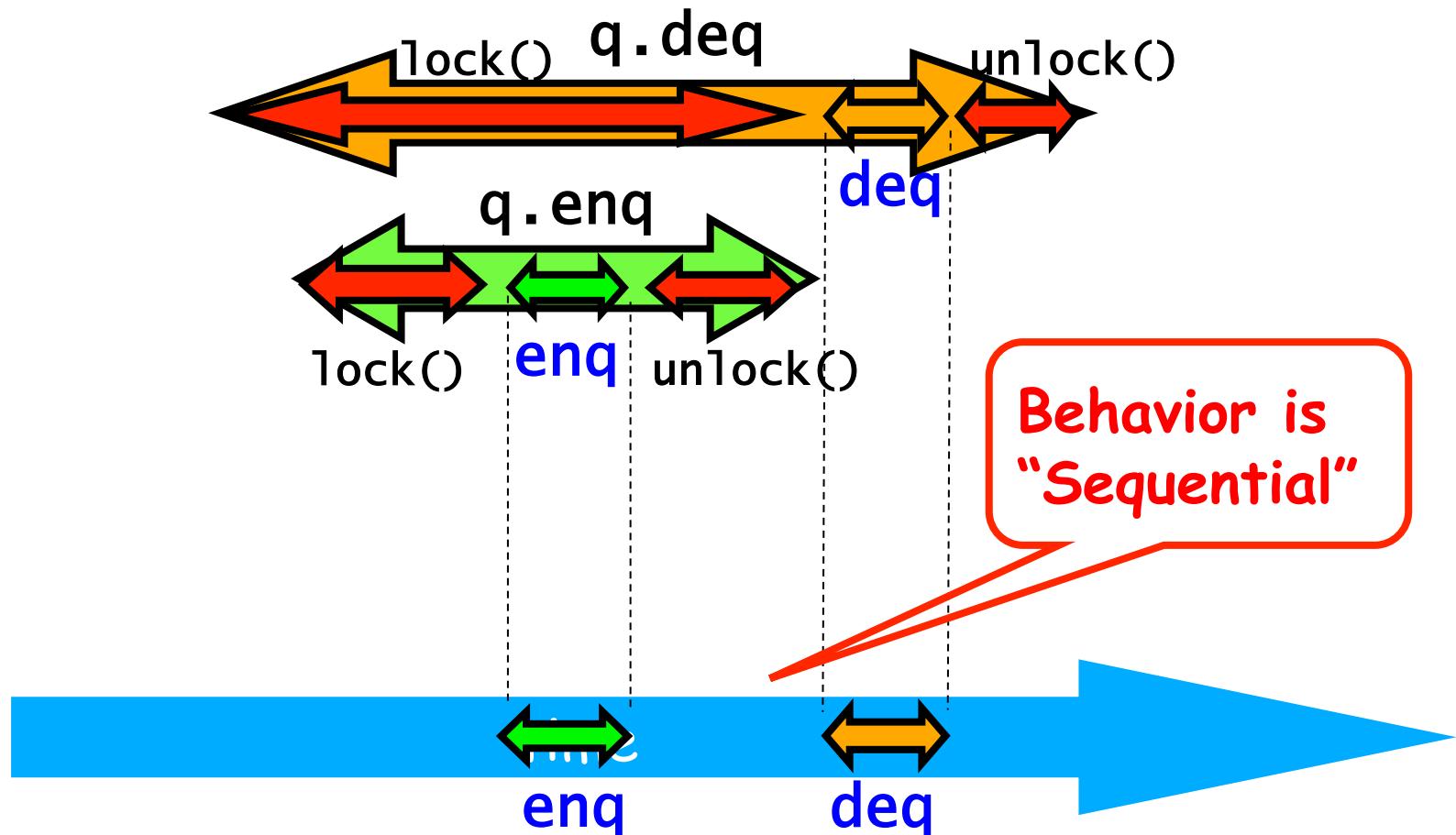
- Quale e' la nozione di correttezza nel caso concorrente?
  - FIFO implica ordine temporale stretto
  - Concorrenza implica un ordine temporale non determinato

```
public T deq() throws EmptyException {
    lock.lock();
    try {
        if (tail == head)
            throw new EmptyException();
        T x = items[head % items.length];
        head++;
        return x;
    } finally {
        lock.unlock();
    }
}
```

```
public T deq() throws EmptyException {  
    lock.lock();  
    try {  
        if (tail == head)  
            throw new EmptyException();  
        T x = items[head % items.length];  
        head++;  
        return x;  
    } finally {  
        lock.unlock();  
    }  
}
```

All modifications  
of queue are done  
mutually exclusive

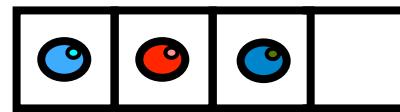
## Catturiamo la concorrenza mediante l'ordine Degli eventi



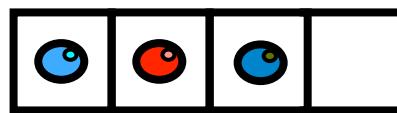
# Linearizability

- Un oggetto e' corretto se la sua proiezione sequenziale e' corretta
  - Linearizable

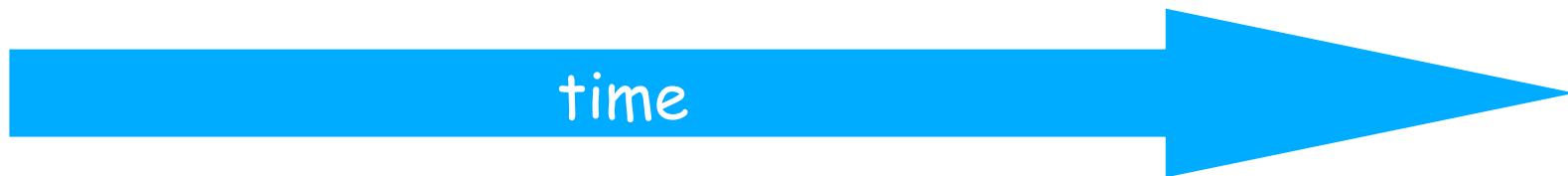
Un oggetto e'  
linearizable: se tutte  
le sue possibili  
esecuzioni sono  
linearizable

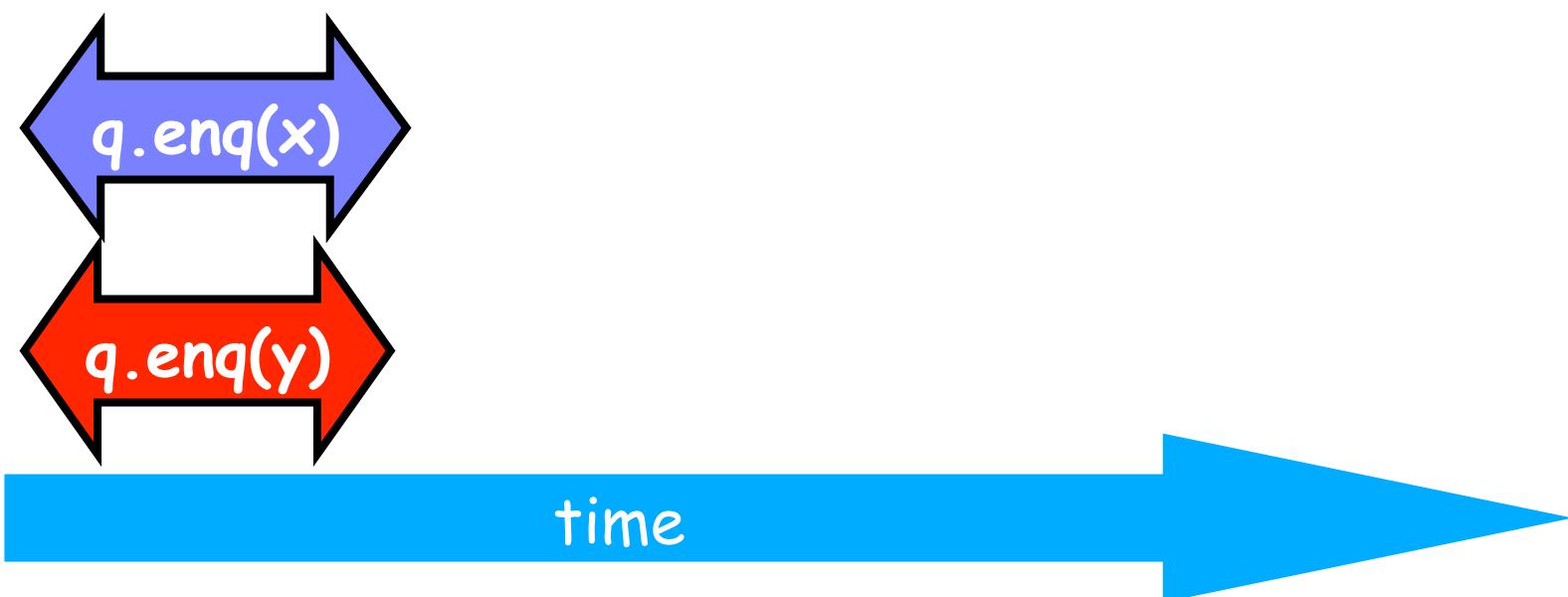
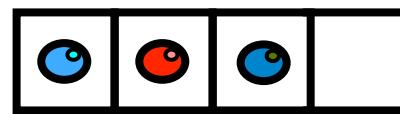


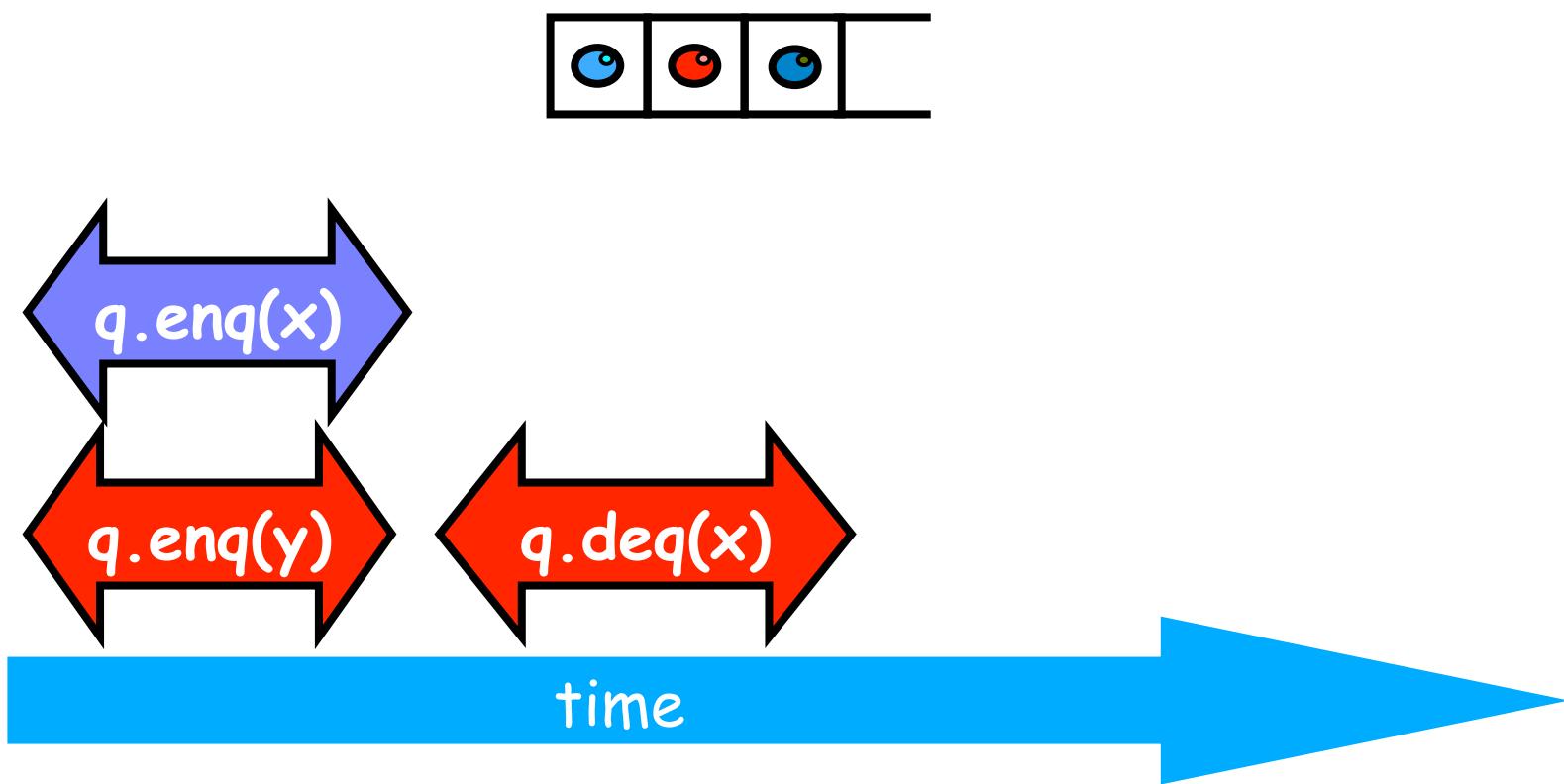
time



$\leftarrow$   
 $\text{q.enq}(x)$   
 $\rightarrow$

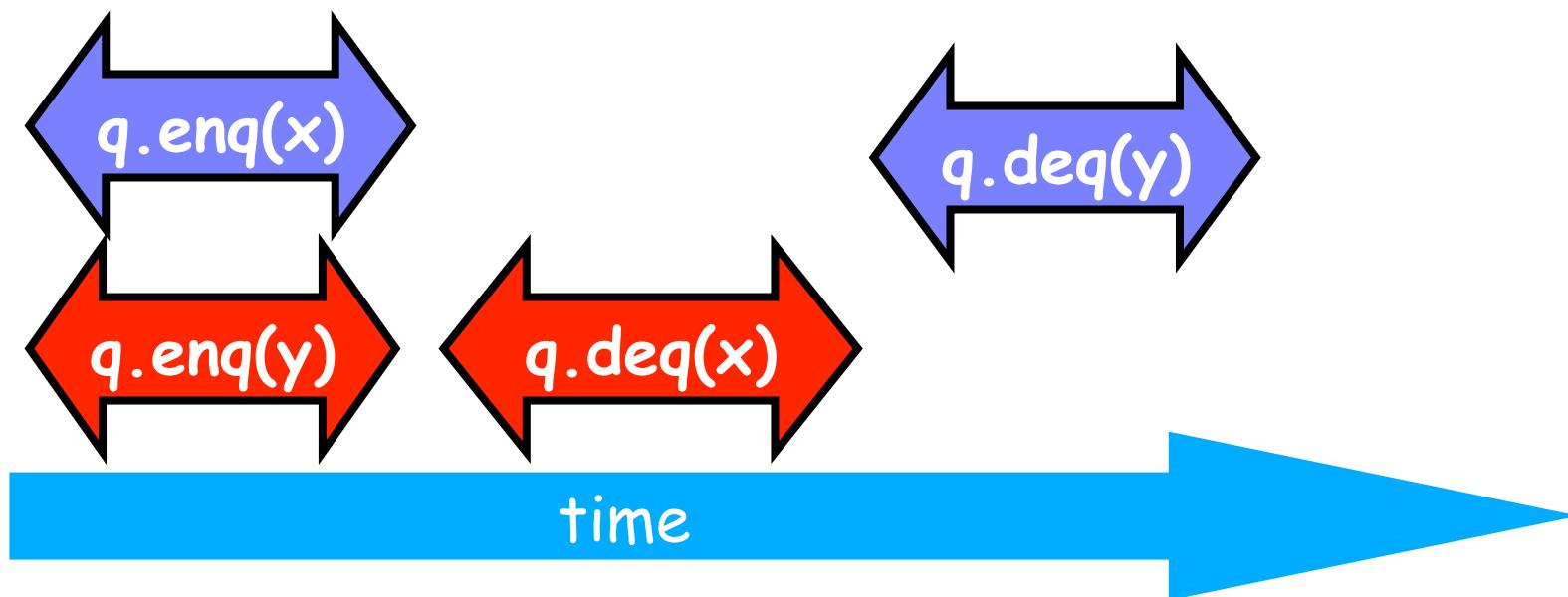
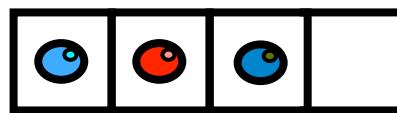


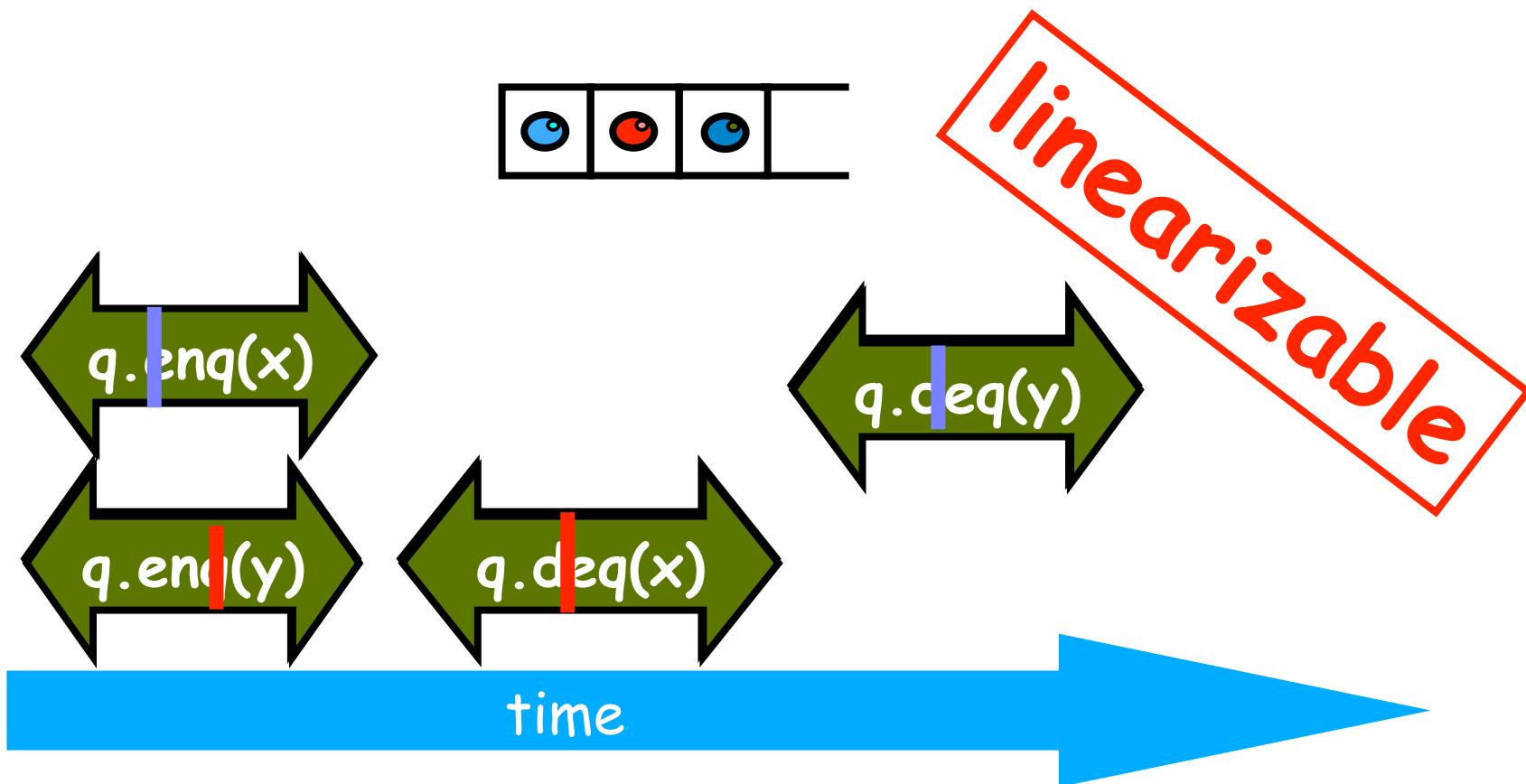




(6)

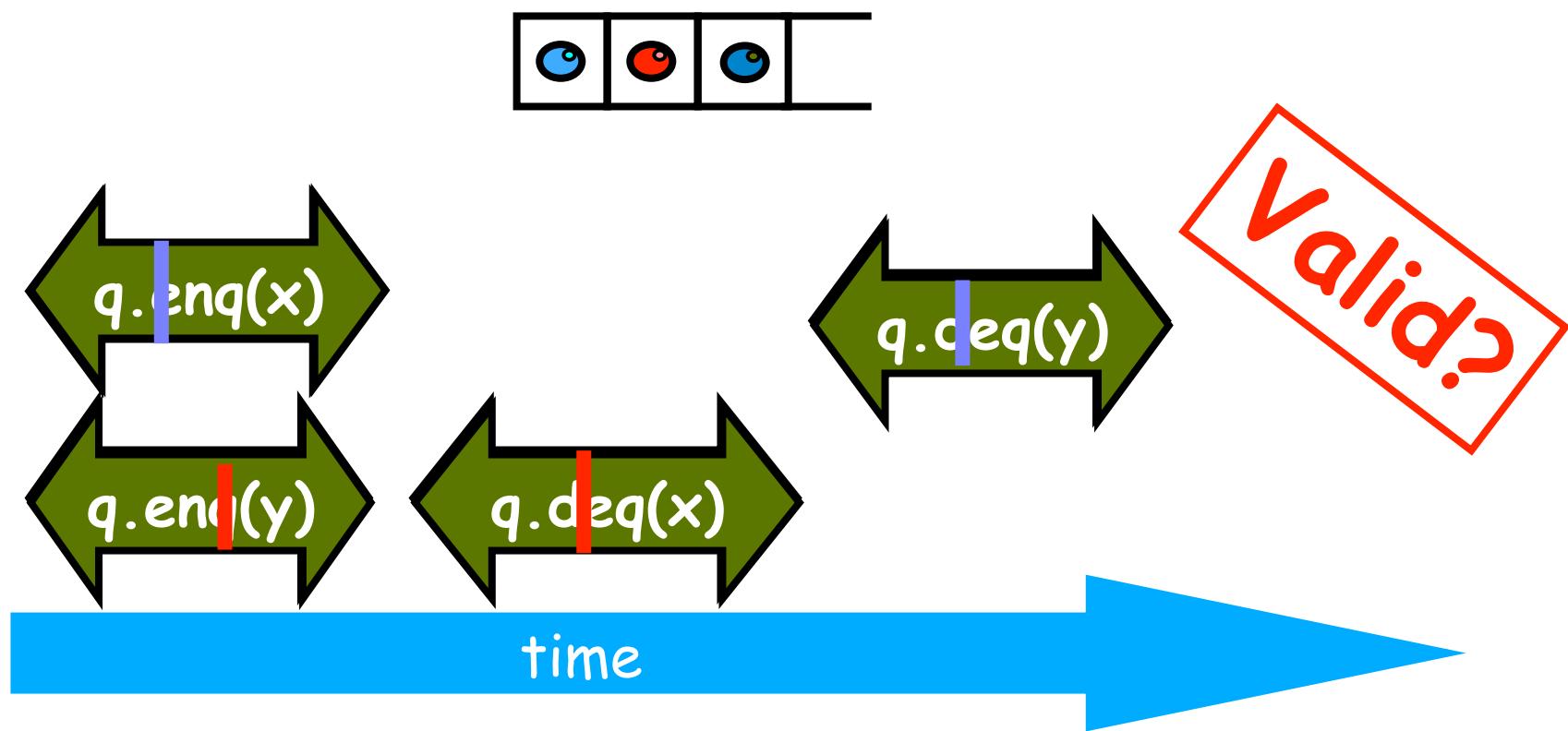
48





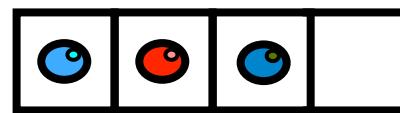
(6)

50

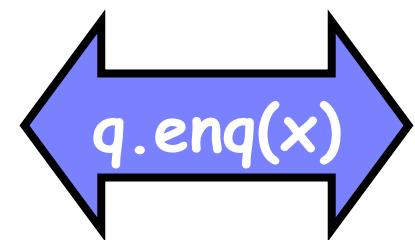
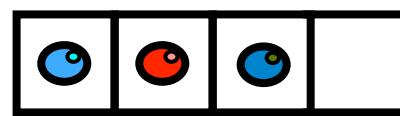


(6)

51



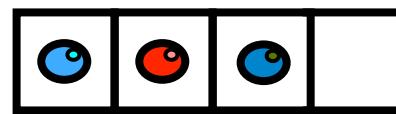
time



A diagram showing the enqueue operation `q.enq(x)`. It features a large purple double-headed arrow pointing horizontally. Inside the arrow, the text `q.enq(x)` is written in white. Below the arrow is a smaller, empty white rectangular box.



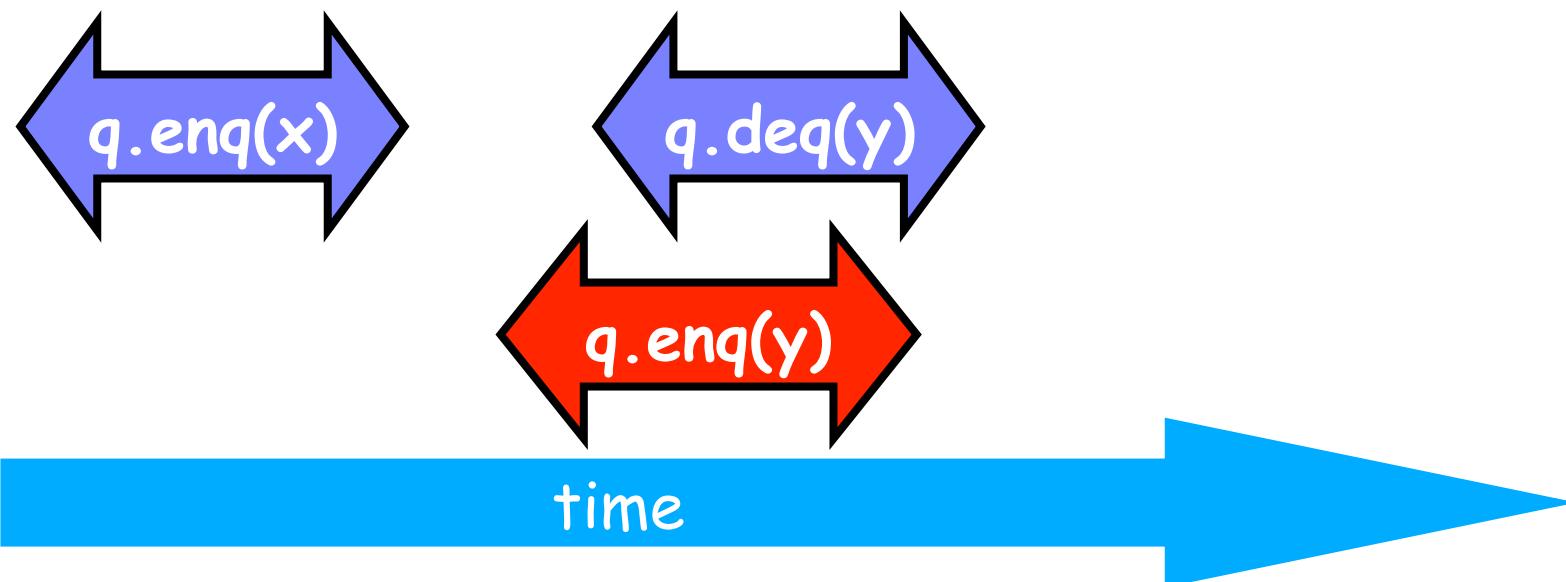
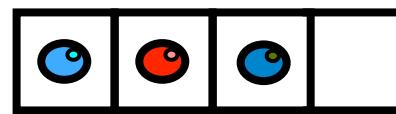
A diagram showing the progression of time. A long, thick blue arrow points to the right. The word "time" is written in white in the center of the arrow.

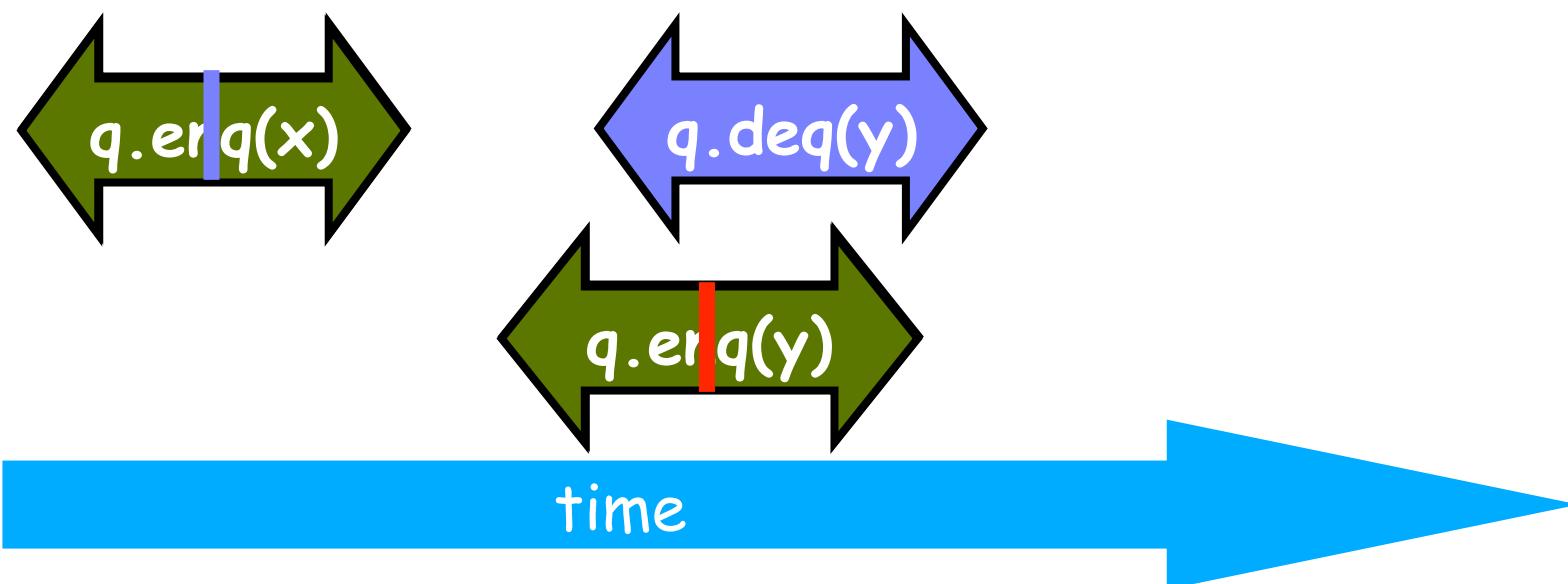
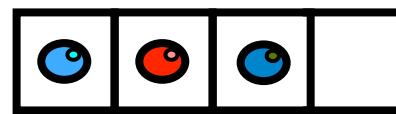


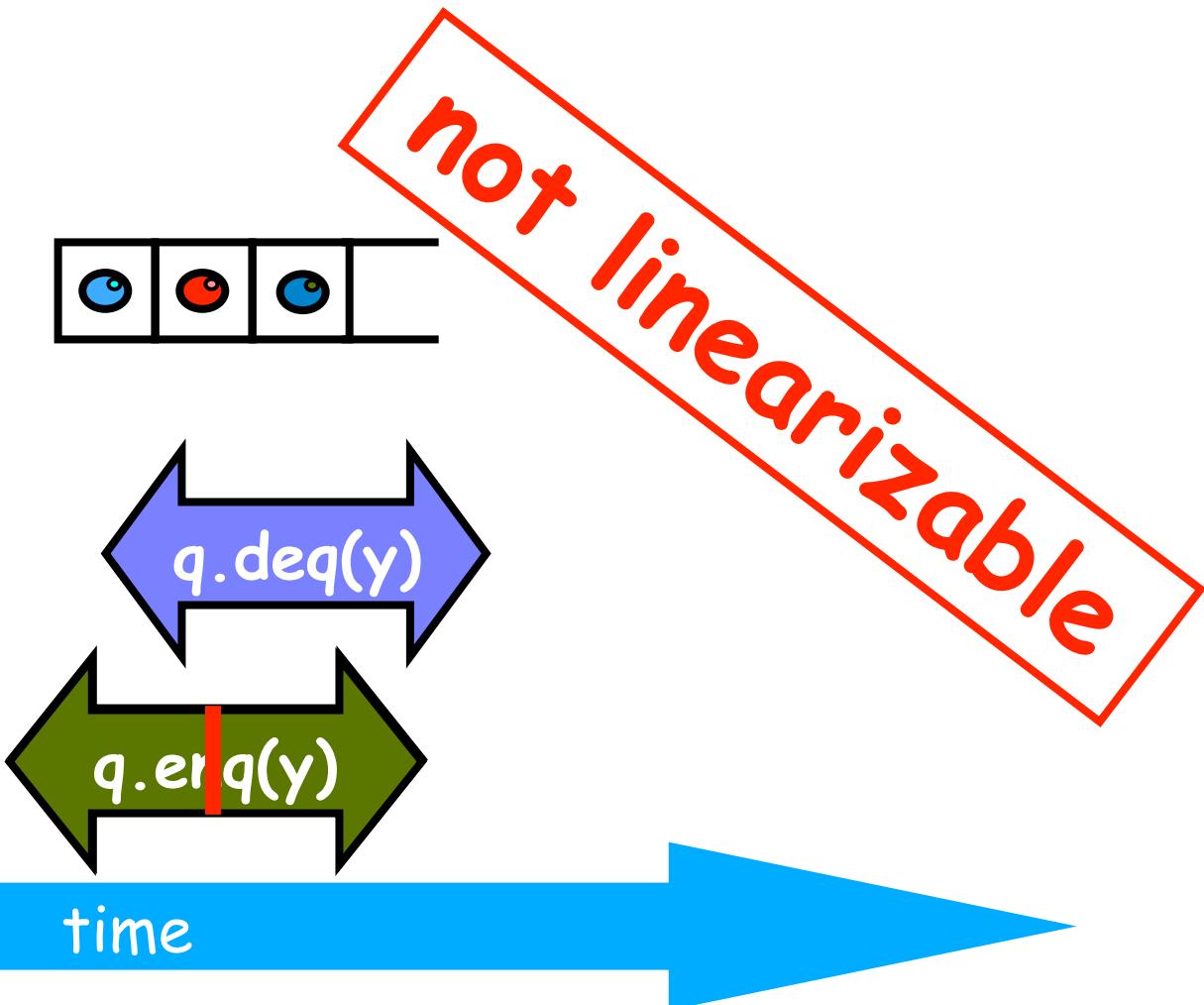
$\longleftrightarrow$   
q.enq(x)

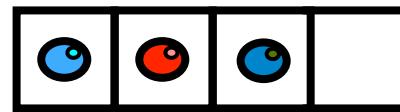
$\longleftrightarrow$   
q.deq(y)

time

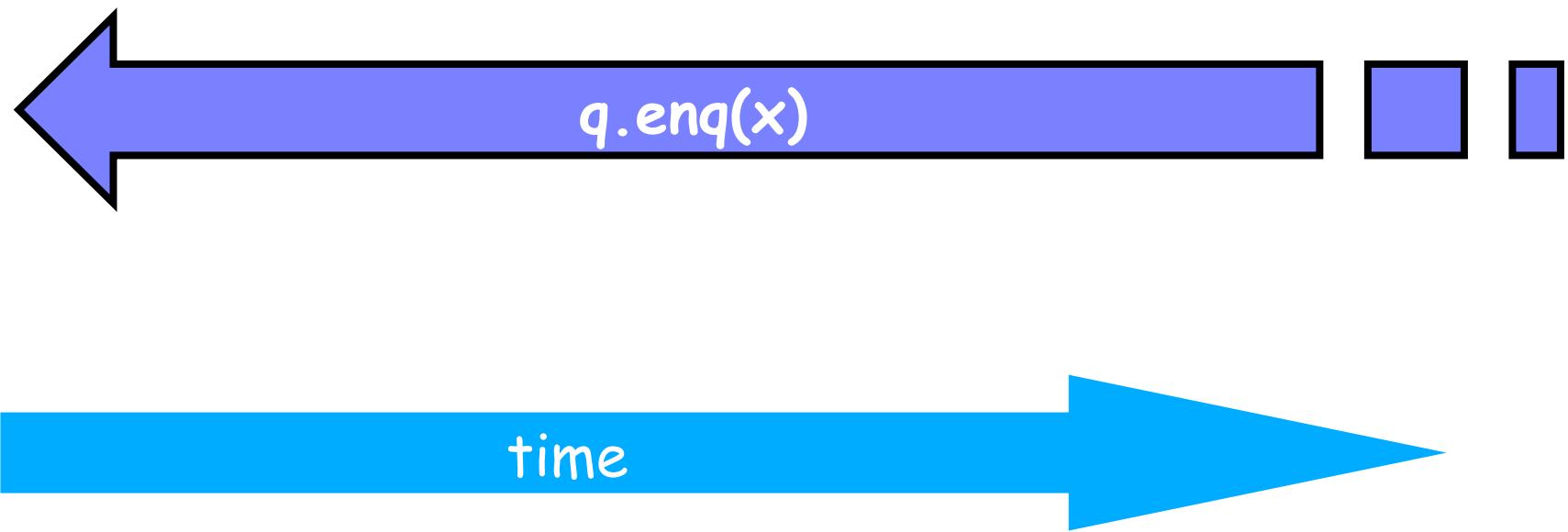
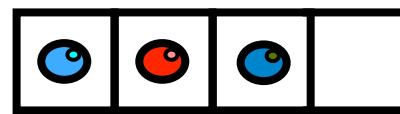


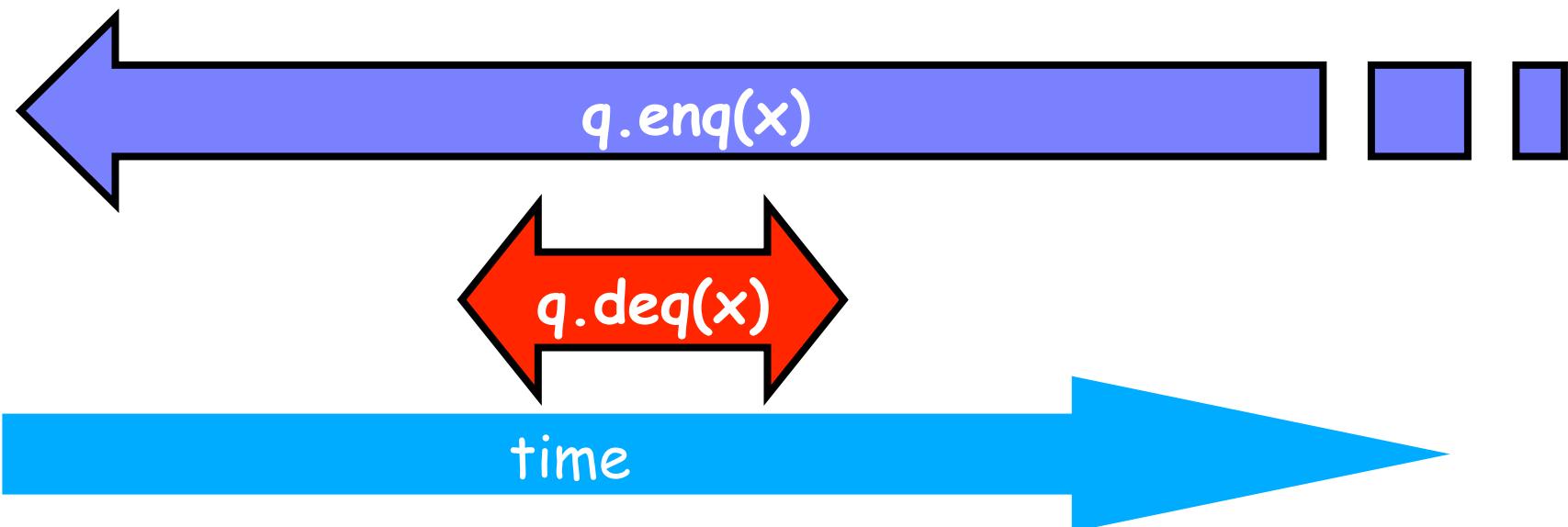
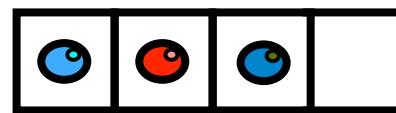


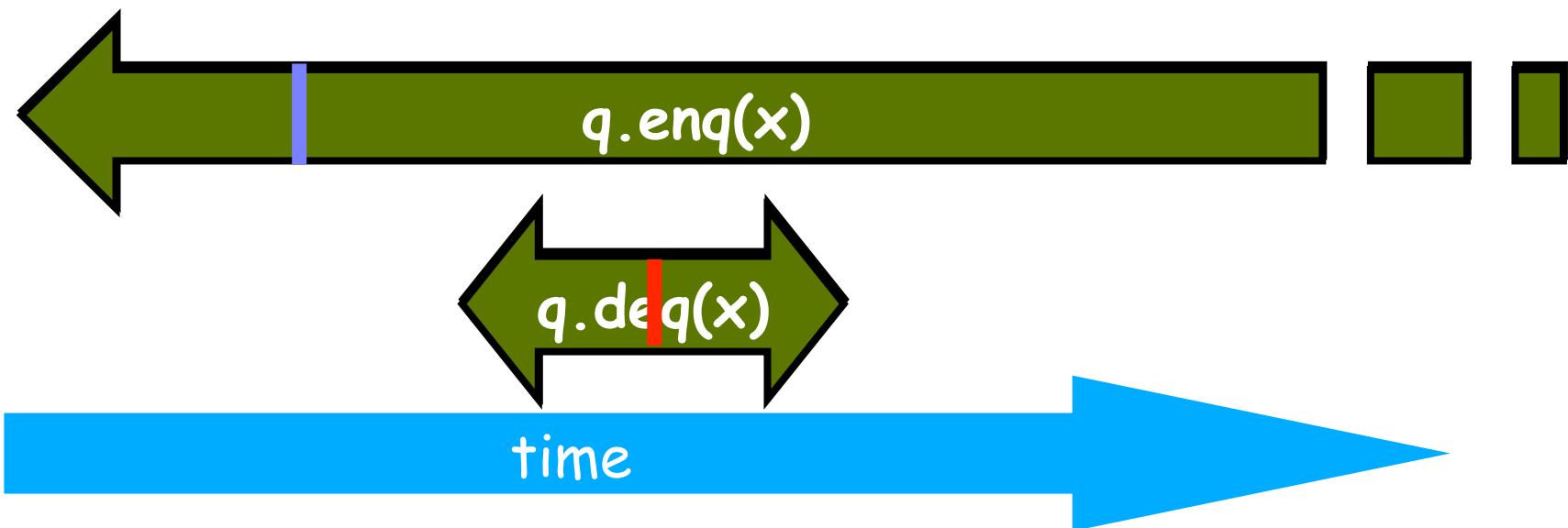
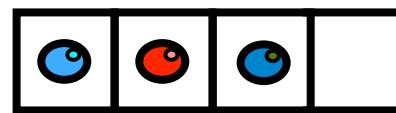


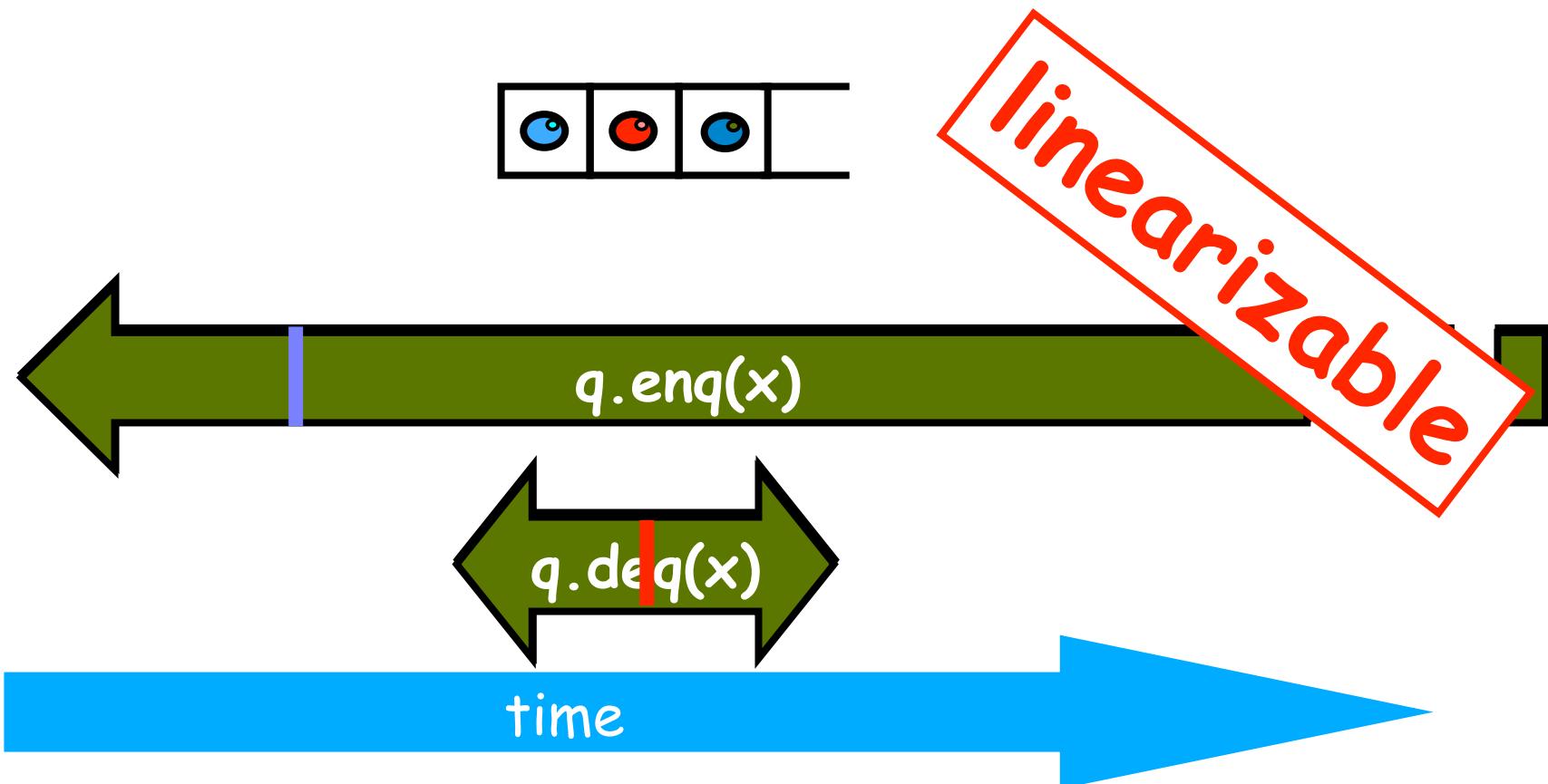


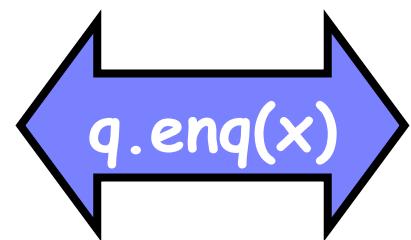
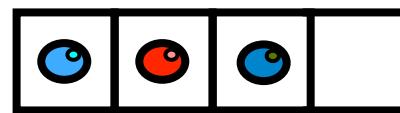
time



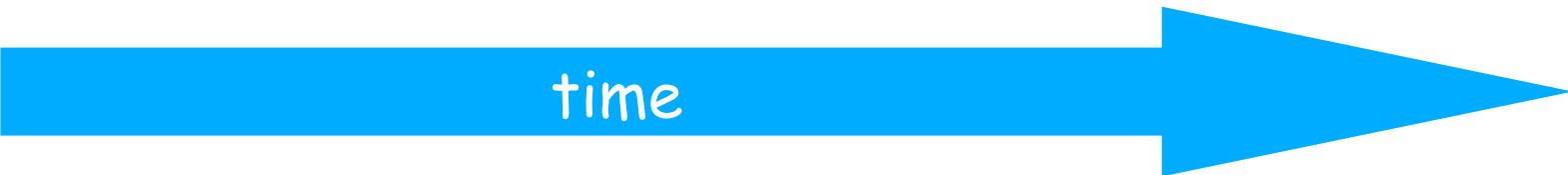








A diagram showing the enqueue operation `q.enq(x)`. It features a large purple double-headed arrow pointing horizontally. Inside the arrow, the text `q.enq(x)` is written in white. Below the arrow is a smaller, empty white rectangular box, likely representing the position where the new element will be inserted.



A diagram showing the progression of time. A long, thick blue arrow points to the right, with the word `time` written in white in the center of the arrow's body.

