#### **301AA - Advanced Programming**

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**AP-15**: Functional Programming

## **Functional Programming - Outline**

- Historical origins
- Main concepts
- Languages families: LISP, ML, and Haskell
- Core concepts of Haskell
- Lazy evaluation

#### Functional Programming: Historical Origins

- The imperative and functional models grew out of work undertaken Alan Turing, Alonzo Church, Stephen Kleene, Emil Post, etc. ~1930s
  - different formalizations of the notion of an algorithm, or *effective procedure*, based on automata, symbolic manipulation, recursive function definitions, and combinatorics
- These results led Church to conjecture that any intuitively appealing model of computing would be equally powerful as well
  - this conjecture is known as *Church's thesis*

# **Historical Origins**

- Church's model of computing is called the *lambda* calculus
  - based on the notion of *parameterized expressions* (parameters introduced by letter  $\lambda$ )
  - allows one to define mathematical functions in a constructive/effective way
  - lambda calculus was the inspiration for functional programming
  - computation proceeds by substituting parameters into expressions, just as one computes in a high level functional program by passing arguments to functions
- We shall see later the basics of lambda-calculus

## **Functional Programming Concepts**

- Functional languages such as LISP, Scheme, FP, ML, Miranda, and Haskell are an attempt to realize Church's lambda calculus in practical form as a programming language
- The key idea: do everything by composing functions
  - no mutable state
  - no side effects

# Functional Programming Concepts

- Necessary features, many of which are missing in some imperative languages:
  - First class and high-order functions
    - Functions can be denoted, passed as arguments to functions, returned as result of function invocation
    - Meaningful because new functions can be defined
  - Recursion
    - Takes the place of iteration (no "control variables")
  - Powerful list facilities
    - Recursive functions exploit recursive definition of lists
  - Polymorphism (typically universal parametric implicit)
    - Relevance of Containers/Collections

## **Functional Programming Concepts**

#### - Fully general aggregates

- Wide use of tuples and records
- Data structures cannot be modified, have to be recreated
- Structured function returns
  - No side-effects, thus the only way for functions to pass information to the caller

#### – Garbage collection

- In case of static scoping, unlimited extent for:
  - locally allocated data structures
  - locally defined functions
- They cannot be allocated on the stack

# The LISP family of languages

- LISP (LISt Processor) was designed in 1958 by John McCarty (Turing award in 1971) and implemented in 1960 by Steve Russel
- Only FORTRAN is older...
- Main programming language for AI before Python
- It includes some features that are not necessary present in other functional languages:
  - Programs (S-expressions) are data (lists)
    - (func arg1 arg2 ... argn)
  - Self-definition
    - A LISP interpreter can be written in few LISP lines
  - Read-evaluate-print interactive loop

# The LISP family of languages

- Variants of LISP
  - (Original) LISP
    - purely functional
    - strong dynamic type checking
    - dynamically scoped
  - Common Lisp: current standard
    - statically scoped
    - very rich and complex
  - Scheme:
    - statically scoped
    - essential syntax
    - very elegant
    - widely used for teaching

#### Other functional languages: the ML family

- Robin Milner (Turing award in 1991, CCS, Pi-calculus, ...)
- Statically typed, general-purpose programming language
   "Meta-Language" of the LCF theorem proving system
- Type safe, with **type inference** and **formal semantics**
- Compiled language, but intended for interactive use
- Combination of Lisp and Algol-like features
  - Expression-oriented
  - Higher-order functions
  - Garbage collection
  - Abstract data types
  - Module system
  - Exceptions
- Impure: it allows side-effects
- Members of the family: Standard ML, Caml, OCaml, F<sup>#</sup>

#### Other functional languages: Haskell

- Designed by committee in 80's and 90's to unify research efforts in lazy languages
  - Evolution of Miranda, name from Haskell Curry, logician (1900-82),
  - Haskell 1.0 in 1990, Haskell '98, Haskell 2010, GHC2021 extension,...
- Several features in common with ML, but some differ:
- Types and type checking
  - Type inference
  - Implicit parametric polymorphism
  - Ad hoc polymorphism (overloading) with type classes
- Control
  - Lazy evaluation
  - Tail recursion and continuations
- Purely functional
  - Precise management of effects

### **Downloading Haskell**

https://www.haskell.org/

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Declarative, statically typed code.

```
primes = filterPrime [2..] where
filterPrime (p:xs) =
    p : filterPrime [x | x <- xs, x `mod` p /= 0]</pre>
```



For playing with Haskell now, use an online interpreter like <u>https://replit.com/</u> <u>https://play.haskell.org/</u> https://godbolt.org/

## Core Haskell

- Basic Types
  - Unit
  - Booleans
  - Integers
  - Strings
  - Reals
  - Tuples
  - Lists
  - Records

- Patterns
- Declarations
- Functions
- Polymorphism
- Type declarations
- Type Classes
- Monads
- Exceptions

#### **Overview of Haskell**

- Interactive Interpreter (ghci): read-eval-print
  - ghci infers type before compiling or executing
  - Type system does not allow casts or similar things!
- Examples

## Overview by Type

Booleans

```
True, False :: Bool
not :: Bool -> Bool
and, or :: Foldable t => t Bool -> Bool
if ... then ... else ...
--conditional expression: types must match
```

• Characters & Strings

'a','b',';','\t', '2', 'X' :: Char
"Ron Weasley" :: [Char] --strings are lists of chars

## Overview by Type

• Numbers

```
0,1,2,...:: Num p => p --type classes, to disambiguate
1.0, 3.1415 :: Fractional a \Rightarrow a
(45 :: Integer) :: Integer -- explicit typing
+, * , -, ... :: Num a => a -> a -> a
-- infix + becomes prefix (+)
-- prefix binary op becomes infix `op`
   :: Fractional a => a -> a -> a
div, mod :: Integral a => a -> a -> a
^:: (Num a, Integral b) => a -> b -> a
```

#### Simple Compound Types

• Tuples

("AP",2017) :: Num b => ([Char], b) -- pair
fst :: (a, b) -> a -- selector: only for pairs
snd :: (a, b) -> b -- selector: only for pairs
('4', True, "AP") :: (Char, Bool, [Char]) -- tuple

• Lists

[] :: [a] -- NIL, polymorphic type 1 : [2, 3, 4] :: Num a => [a]-- infix cons notation [1,2]++[3,4] :: Num a => [a] -- concatenation head :: [a] -> a -- first element tail :: [a] -> [a] -- rest of the list

• Records

#### More on list constructors

ghci> [1..20] -- range [1,2,3,4,5,6,7,8,9,10,11,12,13,14,15,16,17,18,19,20] ghci> ['a'..'z'] "abcdefghijklmnopqrstuvwxyz" ghci> [3,6..20] -- range with step [3,6,9,12,15,18] ghci> [7,6..1] [7,6,5,4,3,2,1]

```
ghci> [1..] -- an infinite list: runs forever
ghci> take 10 [1..] -- prefix of an infinite lists
[1,2,3,4,5,6,7,8,9,10] -- returns!
ghci> take 10 (cycle [1,2])
[1,2,1,2,1,2,1,2,1,2]
ghci> take 10 (repeat 5)
[5,5,5,5,5,5,5,5,5]
```

How does it work??? Later...

### **Binding variables**

- Variables (names) are bound to expressions, without evaluating them (because of lazy evaluation)
- The scope of the binding is the rest of the session
- Comparing OCaml and Haskell

HASKELL	OCaml
Prelude> let a = 6 no output	# let a = 6 ;;
	val a : int = 6
<pre>Prelude&gt; b = a + 2'let' optional</pre>	<b>#</b> let b = a + 2 ;;
	<pre>val b : int = 8</pre>
Prelude> b now b is evaluated	# b ;;
8	- : int = 8
Prelude> a = a + 1 no output	# let a = a + 1 ;;
<pre>Prelude&gt; a what does it print?</pre>	val a : int = 7
^CInterrupted loop broken	

#### Patterns and Declarations

- Patterns can be used in place of variables
   <pat> ::= <var> | <tuple> | <cons> | <record> ...
- Value declarations
  - General form: <pat> = <exp>
  - Examples

myTuple = ("Foo", "Bar")
(x,y) = myTuple -- x = "Foo", y = "Bar"
myList = [1, 2, 3, 4]
z:zs = myList -- z = 1, zs = [2,3,4]

Local declarations

let (x, y) = (2, "FooBar") in x \* 4

Anonymous Functions (lambda abstraction)

Anonymous functions

```
x \rightarrow x+1 --like LISP lambda, function (...) in JS

Prelude> (x \rightarrow x+1)5 => 6

Prelude> f = x \rightarrow x+1

Prelude> :t f

f :: Num a => a -> a

Prelude> f 7 => 8
```

Anonymous functions using patterns

```
Prelude> h = \langle (x, y) \rightarrow x+y
h :: Num a => (a, a) -> a
Prelude> h (3, 4) => 7
Prelude> h 3 4 => error
Prelude> k = \langle (z:zs) \rightarrow length zs
k :: [a] -> Int
Prelude> k "hello" => 4
```

#### **Function declarations**

• Function declaration form

<name> <pat<sub>1</sub>> = <exp<sub>1</sub>> <name> <pat<sub>2</sub>> = <exp<sub>2</sub>> ...

Examples

```
f (x,y) = x+y --argument must match pattern (x,y)
length [] = 0
length (x:s) = 1 + length(s)
Prelude> len (z:zs) = length zs
len :: [a] -> Int
Prelude> len [1,2,3] => 2
Prelude> len []
*** Exception: <interactive>:143:5-24: Non-
exhaustive patterns in function len
```

#### More Functions on Lists

Reverse a list

```
reverse [] = [] -- quadratic
reverse (x:xs) = (reverse xs) ++ [x]
```

```
reverse xs = -- linear, tail recursive
    let rev ( [], accum ) = accum
        rev ( y:ys, accum ) = rev ( ys, y:accum )
        in rev ( xs, [] )
```

• Other (higher-order) functions later