Lightweight Reference-Free Variation Detection using the Burrows-Wheeler Transform

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Motivation We study the problem of identifying SNPs and INDELs within a reads set without aligning it against a reference sequence. Most existing tools for this problem are based on de-Bruijn graphs and share some limitations as their order k is usually small (≈ 30 bases) and they do not store k-mer coverage and k-mer adjacency in reads.

Methods We describe a new approach based on the extension of the Burrows-Wheeler transform (BWT) to a collection of strings. We show that the the output of such a transformation can be partitioned in clusters (substrings), each associated with a position of the underlying (unknown) genome; if that position exhibits a variant, then the cluster will contain more than one distinct character.

Results We compared the performance of our tool eBWT2SNP with the state-of-the-art tool DISCOSNP++ on real and simulated Human datasets. Already at 10x coverage, our tool discovers 80% of existing SNPs and 59% of the INDELs (versus 55% and 32% of DISCOSNP++). At 48x, these percentages increase to 96% and 87% (versus 75% and 46% of DISCOSNP++). DISCOSNP++, on the other hand, is more precise: on average, 89% of its output SNPs and 94% of its output INDELs are correct (versus 80% and 90% of our tool). Due to the fact that we use compressed data structures, DISCOSNP++ is also faster. We are currently developing a parallel version our algorithms to become competitive also under this metric.

DEFINITIONS

Our tool relies on the *extended Burrows-Wheeler transform* of the input reads:

Definition: eBWT

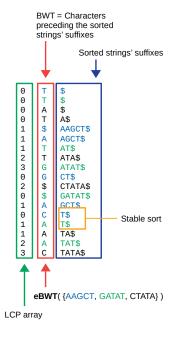
eBWT is the string containing the characters preceding the lexicographically-sorted reads' suffixes. Ties are broken by input order.

We also use the *Longest Common Prefix* array:

Definition: LCP

LCP is the array containing the lengths of the longest common prefixes between adjacent suffixes in lexicographic order.

In the following example, we show <code>ebwt</code> and <code>lcp</code> on the reads set <code>{AAGCT, GATAT, CTATA}</code>. Also the sorted reads' suffixes are shown (blue box). Note that these are shown only for illustrative purposes and are not stored in practice.



PRE-PROCESSING

eBWT can be computed with lightweight tools such as BCR (github.com/BEETL/BEETL) and EGSA (github.com/felipelouza/egsa). Currently, this is the bottleneck of our method: these tools process data at a rate of $\approx 5 \text{GB}$ per hour. We are currently looking into parallel algorithms to speed this step up.

POSITIONAL CLUSTERING

Our method relies on the *clustering* property of the eBWT. In [?] we prove the following theorem:

Theorem [?]

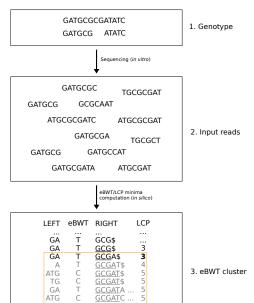
Let i and j be local minima in LCP. With highenough probability, the *cluster* $\mathtt{eBWT}[i, \ldots, j]$ contains the sequenced copies of a single position in the genome.

According to the previous theorem, a genome position contains a variation if and only if the corresponding eBWT cluster contains two different letters. This suggests the following strategy:

Our strategy: eBWT2SNP

- 1. Compute eBWT and LCP.
- 2. Detect LCP minima \Rightarrow eBWT clusters.
- 3. If the cluster contains 2 distinct letters ⇒ variation found.
- 4. Extract context, output variation.

We add INDELs detection w.r.t. the preliminary version [1]. See the following example; in bold: LCP minima. Inside orange box: eBWT cluster.



GΑ

ATGCGC

AT GC

4. Output INDEL

SPACE-EFFICIENT COMPUTATION

Our tool takes as input just eBWT. How do we find LCP minima? In [?], we show:

Theorem [?]

Given eBWT, we can find all LCP minima in linear time using just 1 Byte per base in RAM.

We moreover build a compressed index on top of the eBWT to extract the context surrounding the variations. This strategy uses 8 times less space than the preliminary version [?].

RESULTS

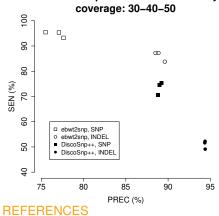
We compared eBWT2SNP with DISCOSNP++, the state-of-the-art tool for reference-free variation detection. All datasets have been downloaded from the 1000genomes website. We processed reads — both real and simulated — from a single individual in order to reconstruct its genotype:

Experiments

- Simulated. Heterozygous reads simulated from HG00096, Chr1, cov. 50x.
 Variations taken from the real VCF file.
- 2. **Real**. Reads sequenced from HG00096, Chr1, cov. 48x.

Results on simulated data:

Simulated – precision and sensitivity coverage: 30–40–50



- [1] N. Prezza, N. Pisanti, M. Sciortino, and G. Rosone. SNPs detection by eBWT positional clustering.
- [2] N. Prezza and G. Rosone. Space-Efficient Computation of the LCP Array from the Burrows-Wheeler Transform. In *CPM 2019*, Leibniz International Proceedings in Informatics (LIPIcs). Schloss Dagstuhl-Leibniz-Zentrum fuer Informatik, 2019.

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