

RECORDS

Record definition

- A **record** contains several named **fields**
- Before you can use a record, must **define** a record type:

```
type time = {hour: int; min: int; ampm: string}
```

- To *build* a record:
 - Write a record expression:
`{hour=10; min=10; ampm="am" }`
 - Order of fields doesn't matter:
`{min=10; hour=10; ampm="am" }` is equivalent
- To *access* record's field: `r.hour`

Record expressions

- **Syntax:** $\{f1 = e1; \dots; fn = en\}$
- **Evaluation:**
 - If $e1$ evaluates to $v1$, and ... en evaluates to vn
 - Then $\{f1 = e1; \dots; fn = en\}$ evaluates to $\{f1 = v1, \dots, fn = vn\}$
 - Result is a *record value*
- **Type-checking:**
 - If $e1 : t1$ and $e2 : t2$ and ... $en : tn$,
 - and if t is a defined type of the form $\{f1 : t1, \dots, fn : tn\}$
 - then $\{f1 = e1; \dots; fn = en\} : t$

Record field access

- **Syntax:** $e.f$
- **Evaluation:**
 - If e evaluates to $\{f = v, \dots\}$
 - Then $e.f$ evaluates to v
- **Type-checking:**
 - If $e : t_1$
 - and if t_1 is a defined type of the form $\{f : t_2, \dots\}$
 - then $e.f : t_2$

Evaluation notation

We keep writing statements like:

If e evaluates to $\{f = v, \dots\}$ then $e.f$ evaluates to v

Let's introduce a shorthand notation:

- Instead of " e evaluates to v "
- write " $e \Rightarrow v$ "

So we can now write:

If $e \Rightarrow \{f = v, \dots\}$ then $e.f \Rightarrow v$

By name vs. by position

- Fields of record are identified **by name**
 - order we write fields in expression is irrelevant
- Opposite choice: identify **by position**
 - e.g., “Would the student named NN. step forward?”
vs. “Would the student in seat n step forward?”
- You’re accustomed to both:
 - Java object fields accessed by name
 - Java method arguments passed by position
(but accessed in method body by name)
- OCaml has something you might not have seen:
 - A kind of data accessed by position

PAIRS AND TUPLES

Pairs

A **pair** of data: two pieces of data glued together
e.g.,

- (1 , 2)
- (**true** , "Hello")
- ([1 ; 2 ; 3] , 0.5)

We need language constructs to *build* pairs and to *access* the pieces...

Pairs: building

- Syntax: $(e1, e2)$
- Evaluation:
 - If $e1 \implies v1$ and $e2 \implies v2$
 - Then $(e1, e2) \implies (v1, v2)$
 - A pair of values is itself a value
- Type-checking:
 - If $e1 : t1$ and $e2 : t2$,
 - then $(e1, e2) : t1 * t2$
 - A new kind of type, the **product type**

Pairs: accessing

- **Syntax:** `fst e` and `snd e`
Projection functions
- **Evaluation:**
 - If `e ==> (v1, v2)`
 - then `fst e ==> v1`
 - and `snd e ==> v2`
- **Type-checking:**
 - If `e: ta*tb`,
 - then `fst e` has type `ta`
 - and `snd e` has type `tb`

Tuples

Actually, you can have *tuples* with more than two parts

- A new feature: a generalization of pairs
- Syntax, semantics are straightforward, except for projection...

- (e_1, e_2, \dots, e_n)
- $t_1 * t_2 * \dots * t_n$
- `fst e, snd e, ???`

Instead of generalizing projection functions,
use **pattern matching**...

New kind of pattern, the **tuple pattern**: (p_1, \dots, p_n)

Pattern matching tuples

```
match (1,2,3) with  
| (x,y,z) -> x+y+z
```

```
(* ==> 6 *)
```

```
let thrd t =  
  match t with  
  | (x,y,z) -> z
```

```
(* thrd : 'a*'b*'c -> 'c *)
```

Note: we never needed more than one branch in the match expression...

Pattern matching without match

```
(* OK *)  
let thrd t =  
  match t with  
  | (x,y,z) -> z
```

```
(* good *)  
let thrd t =  
  let (x,y,z) = t in z
```

```
(* better *)  
let thrd t =  
  let (_,_,z) = t in z
```

```
(* best *)  
let thrd (_,_,z) = z
```

Extended syntax for let

- Previously we had this syntax:
 - **let** $x = e1$ **in** $e2$
 - **let** [**rec**] $f\ x1 \dots xn = e1$ **in** $e2$
- Everywhere we had a variable identifier x , we can really use a pattern!
 - **let** $p = e1$ **in** $e2$
 - **let** [**rec**] $f\ p1 \dots pn = e1$ **in** $e2$
- Old syntax is just a **special case** of new syntax, since a variable identifier is a pattern

Pattern matching arguments

```
(* OK *)  
let sum_triple t =  
    let (x,y,z) = t  
    in x+y+z
```

```
(* better *)  
let sum_triple (x,y,z) = x+y+z
```

Note how that last version looks syntactically like a function in C/Java!

Unit

- Can actually have a tuple `()` with no components whatsoever
 - Think of it as a degenerate tuple
 - Or, like a Boolean that can only have one value
- “Unit” is
 - a value written `()`
 - and a type written `unit`
- Might seem dumb now; will be useful later!

Pattern matching records

```
(* OK *)  
let get_hour t =  
  match t with  
  | {hour=h; min=m; ampm=s} -> h
```

```
(* better *)  
let get_hour t =  
  match t with  
  | {hour=h; min=_; ampm=_} -> h
```

```
(* better *)  
let get_hour t =  
  match t with  
  | {hour; min; ampm} -> hour
```

```
(* better *)  
let get_hour t =  
  match t with  
  | {hour} -> hour
```

```
(* better *)  
let get_hour t =  
  let {hour} = t in hour
```

```
(* better *)  
let get_hour {hour} = hour
```

```
(* best *)  
let get_hour t = t.hour
```

New kind of pattern, the **record pattern**:

```
{f1[=p1]; ...; fn[=pn]}
```

By name vs. by position, again

How to choose between coding $(4, 7, 9)$ and $\{f=4; g=7; h=9\}$?

- Tuples are syntactically shorter
- Records are self-documenting
- For many (4? 8? 12?) fields, a record is usually a better choice

VARIANTS

Variant

```
type day = Sun | Mon | Tue | Wed  
          | Thu | Fri | Sat
```

```
let day_to_int d =  
    match d with  
    | Sun -> 1  
    | Mon -> 2  
    | Tue -> 3  
    | Wed -> 4  
    | Thu -> 5  
    | Fri -> 6  
    | Sat -> 7
```

Building and accessing variants

Syntax: $\text{type } t = C_1 \mid \dots \mid C_n$

the C_i are called *constructors*

Evaluation: a constructor is already a value

Type checking: $C_i : t$

Accessing: use pattern matching; constructor name is a pattern

Pokémon variant



DEFENSE → ATTACK ↴	NOR	FIR	WAT	E
NORMAL				
FIRE		½	½	
WATER		2	½	

Pokémon variant



```
type ptype = TNormal | TFire | TWater
```

```
type peff = ENormal | ENotVery | ESuper
```

```
let eff_to_float = function
```

```
| ENormal   -> 1.0
```

```
| ENotVery  -> 0.5
```

```
| ESuper    -> 2.0
```

```
let eff_att_vs_def : ptype*ptype -> peff = function
```

```
| (TFire,TFire)   -> ENotVery
```

```
| (TWater,TWater) -> ENotVery
```

```
| (TFire,TWater)  -> ENotVery
```

```
| (TWater,TFire)  -> ESuper
```

```
| _ -> ENormal
```

Argument order: records

If you are worried about clients of function forgetting which order to pass arguments in tuple, use a record:

```
type att_def = {att:ptype; def:ptype}
```

```
let eff_att_vs_def : att_def -> peff = function  
  | {att=TFire;def=TFire}    -> ENotVery  
  | {att=TWater;def=TWater} -> ENotVery  
  | {att=TFire;def=TWater}  -> ENotVery  
  | {att=TWater;def=TFire}  -> ESuper  
  | _ -> ENormal
```


Argument order: labeled arguments

Or (though not quite as good) use **labeled arguments**:

```
let eff_att_vs_def ~att ~def =  
  match (att, def) with  
  | (TFire,TFire)    -> ENotVery  
  | (TWater,TWater) -> ENotVery  
  | (TFire,TWater)  -> ENotVery  
  | (TWater,TFire)  -> ESuper  
  | _                -> ENormal
```

```
let super = eff_att_vs_def ~att:TWater ~def:TFire
```

```
let super = eff_att_vs_def ~def:TFire ~att:TWater
```

```
let notvery = eff_att_vs_def TFire TWater
```

Variants vs. records vs. tuples

	Define	Build/construct	Access/destruct
Variant	type	Constructor name	Pattern matching
Record	type	Record expression with { ... }	Pattern matching OR field selection with dot operator .
Tuple	N/A	Tuple expression with (...)	Pattern matching OR fst or snd

- Variants: **one-of types** *aka* **sum types**
- Records, tuples: **each-of types** *aka* **product types**

Question

Which of the following would be better represented with records rather than variants?

- A. *Coins*, which can be pennies, nickels, dimes, or quarters
- B. *Students*, who have names and id numbers
- C. *A plated dessert*, which has a sauce, a creamy component, and a crunchy component
- D. A and C
- E. B and C

Question

Which of the following would be better represented with records rather than datatypes?

- A. Coins, which can be pennies, nickels, dimes, or quarters
- B. Students, who have names **and** NetIDs
- C. *A plated dessert*, which has a sauce, a creamy component, **and** a crunchy component
- D. A and C
- E. B and C**

OPTIONS

What is max of empty list?

```
let rec max_list = function
  | []      -> ???
  | h::t   -> max h (max_list t)
```

How to fill in the ???

- **min_int** would be a reasonable choice...
- or could raise an exception...
- in Java, might return **null**...
- but OCaml gives us another option!

Options

Options:

- `t option` is a type for any type `t`
(much like `t list` is a type for any type `t`)

Building and Type Checking and Evaluation:

- `None` has type `'a option`
 - much like `[]` has type `'a list`
 - `None` is a value
- `Some e : t option` if `e : t`
 - much like `e :: []` has type `t list` if `e : t`
 - If `e ==> v` then `Some e ==> Some v`

Accessing: `match e with`
 `None -> ...`
 `| Some x -> ...`

Again: What is max of empty list?

```
let rec max_list = function  
  | []    -> None  
  | h::t -> match max_list t with  
    | None    -> Some h  
    | Some x -> Some (max h x)  
  
(* max_list : 'a list -> 'a option *)
```

Very stylish!

...no possibility of exceptions

...no chance of programmer ignoring a "null return"

Recap: User-defined data types

- Records
- Tuples (pairs, unit)
- Variants
- Options