

Review: higher-order functions

- Functions are values
- Can use them **anywhere** we use values
 - Arguments, results, parts of tuples, bound to variables...
- Functions can **take** functions as arguments
- Functions can **return** functions as results

Review: anonymous functions

(aka *function expressions*)

- **Syntax:** `fun x -> e`
- **Type checking:**
 - Conclude that `fun x -> e : t1 -> t2` if `e : t2` under assumption `x : t1`
- **Evaluation:**
 - A function is already a value

Lambda

- Anonymous functions a.k.a. *lambda expressions*: $\lambda x . e$
- The lambda means “what follows is an anonymous function”
 - x is its argument
 - e is its body
 - Just like **fun** $x \rightarrow e$, but slightly different syntax
- Standard feature of any functional language (ML, Haskell, Scheme, ...)
- You’ll see “lambda” show up in many places in PL, e.g.:
 - PHP: <http://www.php.net/manual/en/function.create-function.php>
 - A popular PL blog: <http://lambda-the-ultimate.org/>
 - Lambda style: <https://www.youtube.com/watch?v=Ci48kqp11F8>

Map

```
let rec map f = function  
| [] -> []  
| x::xs -> (f x)::(map f xs)
```

```
map : ('a -> 'b) -> 'a list -> 'b list
```

Map is HUGE:

- You use it **all the time** once you know it
- Exists in standard library as **List.map**, but the idea can be used in any data structure (trees, stacks, queues...)

Question

What is value of `lst` after this code?

```
let is_even x = (x mod 2 = 0)
let lst = map is_even [1;2;3;4]
```

- A. [1;2;3;4]
- B. [2;4]
- C. [false; true; false; true]
- D. false

Question

What is value of `lst` after this code?

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Question

What is value of `lst` after this code?

```
let is_even x = (x mod 2 = 0)
let lst = filter is_even [1;2;3;4]
```

- A. [1;2;3;4]
- B. [2;4]
- C. [false; true; false; true]
- D. false

Filter

```
let rec filter f = function
  | [] -> []
  | x::xs -> if f x
              then x::(filter f xs)
              else filter f xs
```

```
val filter : ('a -> bool) -> 'a list -> 'a list = <fun>
```


Question

What is value of `lst` after this code?

```
let is_even x = (x mod 2 = 0)
let lst = filter is_even [1;2;3;4]
```

- A. [1;2;3;4]
- B. [2;4]**
- C. [false; true; false; true]
- D. false

Iterators

- Map and filter are *iterators*
 - Not built-in to the language, an idiom
- Benefit of iterators: separate recursive traversal from data processing
 - Can reuse same traversal for different data processing
 - Can reuse same data processing for different data structures
 - leads to modular, maintainable, beautiful code!
- So far: iterators that change or omit data
 - what about combining data?
 - e.g., sum all elements of list

Fold v1.0

Idea: *stick an operator between every element of list*

folding [**1 ; 2 ; 3**] with (+)

becomes

1+2+3

==>

6

Fold v2.0

Idea: *stick an operator between every element of list*
But list could have 1 element, so need an initial value

folding **[1]** with **0** and **(+)**

becomes

0+1

==>

1

Fold v2.0

Idea: *stick an operator between every element of list*
But list could have 1 element, so need an initial value

folding [**1 ; 2 ; 3**] with **0** and **(+)**

becomes

0+1+2+3

==>

6

Fold v2.0

Idea: *stick an operator between every element of list*
But list could have 1 element, so need an initial value
Or list could be empty; just return initial value

folding [] with **0** and **(+)**

becomes

0

Question #4

What should the result of folding [**1 ; 2 ; 3 ; 4**] with **1** and (*) be?

- A. 1
- B. 24
- C. 10
- D. 0

Question #4

What should the result of folding [1 ; 2 ; 3 ; 4] with 1 and (*) be?

A. 1

B. 24

C. 10

D. 0

Fold v3.0

Idea: *stick an operator between every element of list*
But list could have 1 element, so need an initial value
Or list could be empty; just return initial value

Implementation detail: iterate left-to-right or right-to-left?

folding [**1 ; 2 ; 3**] with **0** and **(+)**

left to right becomes: $((0+1)+2)+3$

right to left becomes: $1+(2+(3+0))$

Both evaluate to 6; does it matter?

Yes: not all operators are associative, e.g. subtraction,
division, exponentiation, ...

Fold v4.0

- (+) *accumulated* a result of the same type as list itself
- What about operators that change the type?

– e.g., `::` has type `'a -> 'a list -> 'a list`

folding from the right `[1;2;3]` with `[]` and `::`

should produce

$$1 :: (2 :: (3 :: [])) = [1;2;3]$$

- So the operator needs to accept
 - the accumulated result so far, and
 - the next element of the list

...which may have different types!

Fold for real

Two versions in OCaml library:

```
List.fold_left
```

```
: ('a -> 'b -> 'a) -> 'a -> 'b list -> 'a
```

```
List.fold_right
```

```
: ('a -> 'b -> 'b) -> 'a list -> 'b -> 'b
```

Fold for real

Two versions in OCaml library:

List.fold_left

: ('a -> 'b -> 'a) -> 'a -> 'b list -> 'a

List.fold_right

: ('a -> 'b -> 'b) -> 'a list -> 'b -> 'b

Operator

Fold for real

Two versions in OCaml library:

List.fold_left

: ('a -> 'b -> 'a) -> 'a -> 'b list -> 'a

List.fold_right

: ('a -> 'b -> 'b) -> 'a list -> 'b -> 'b

Input list

Fold for real

Two versions in OCaml library:

List.fold_left

: ('a -> 'b -> 'a) -> 'a -> 'b list -> 'a

List.fold_right

: ('a -> 'b -> 'b) -> 'a list -> 'b -> 'b

Initial value of accumulator

Fold for real

Two versions in OCaml library:

```
List.fold_left
```

```
: ('a -> 'b -> 'a) -> 'a -> 'b list -> 'a
```

```
List.fold_right
```

```
: ('a -> 'b -> 'b) -> 'a list -> 'b -> 'b
```

Final value of accumulator

fold_left

```
let rec fold_left f acc xs =  
  match xs with  
  | []       -> acc  
  | x::xs'   -> fold_left f (f acc x) xs'
```

Accumulates an answer by

- repeatedly applying **f** to “answer so far”,
- starting with initial value **acc**,
- folding “from the left”

fold_left f acc [a;b;c]

computes

f (f (f acc a) b) c

fold_right

```
let rec fold_right f xs acc =  
  match xs with  
  []       -> acc  
  | x::xs' -> f x (fold_right f xs' acc)
```

Accumulates an answer by

- repeatedly applying **f** to “answer so far”,
- starting with initial value **acc**,
- folding “from the right”

fold_right f [a;b;c] acc

computes

f a (f b (f c acc))

Behold the HUGE power of fold

Implement so many other functions with fold!

```
let rev xs = fold_left (fun xs x -> x::xs) [] xs
let length xs = fold_left (fun a _ -> a+1) 0 xs
let map f xs = fold_right
  (fun x a -> (f x)::a) xs []
let filter f xs = fold_right
  (fun x a -> if f x then x::a else a) xs []
```

Beware the efficiency of fold

- **fold_left** is tail recursive, **fold_right** is not
- **fold_right** might make it easier to express computation (e.g., **map**)
- Rule of thumb: for lists with $> 10,000$ elements, use tail recursion

MapReduce

- Fold has many synonyms/cousins in various functional languages, including **scan** and **reduce**
- Google organizes large-scale data-parallel computations with MapReduce
 - open source implementation by Apache called Hadoop

"[Google's MapReduce] abstraction is inspired by the map and reduce primitives present in Lisp and many other functional languages. We realized that most of our computations involved applying a map operation to each logical record in our input in order to compute a set of intermediate key/value pairs, and then applying a reduce operation to all the values that shared the same key in order to combine the derived data appropriately."

[Dean and Ghemawat, 2008]