

# Lists: An introduction

```
let lst = [1;2;3]
```

```
let empty = []
```

```
let longer = 5::lst
```

```
let another = 5::1::2::3::[]
```

```
let rec sum xs =
```

```
  match xs with
```

```
  | [] -> 0
```

```
  | h::t -> h + sum t
```

```
let six = sum lst
```

```
let zero = sum empty
```

# Lists: An introduction

```
let lst = [ "abc"; "def"; "ghi" ]
```

```
let rec concat ss =  
  match ss with  
  | [] -> ""  
  | s::ss' -> s ^ (concat ss')
```

```
let a_i = concat lst
```

# Building lists

## Syntax:

- `[]` is the empty list
- `e1::e2` prepends element `e1` to list `e2`
- `[e1; e2; ...; en]` is *syntactic sugar* for `e1::e2::...::en::[]`

`[]` is pronounced "nil"

`::` is pronounced "cons" (both from LISP)

**Syntactic sugar:** redundant kind of syntax that makes program "sweeter" or easier to write

# Building lists

## Evaluation:

- $[]$  is a value
- To evaluate  $e1 :: e2$ , evaluate  $e1$  to a value  $v1$ , evaluate  $e2$  to a (list) value  $v2$ , and return  $v1 :: v2$

## Consequence of the above rules:

- To evaluate  $[e1; \dots; en]$ , evaluate  $e1$  to a value  $v1$ , ..., evaluate  $en$  to a value  $vn$ , and return  $[v1; \dots; vn]$

# Building lists

## New types:

For any type  $t$ , the type  $t$  `list` describes lists where all elements have type  $t$

- `[1;2;3] : int list`
- `[true] : bool list`
- `[[1+1;2-3];[3*7]] : int list list`

## Nil:

`[] : 'a list`

i.e., empty list has type  $t$  `list` for any type  $t$

## Cons:

If  $e1 : t$  and  $e2 : t$  `list` then  $e1::e2 : t$  `list`

*With parens for clarity:*

If  $e1 : t$  and  $e2 : (t$  `list)` then  $(e1::e2) : (t$  `list)`

# Accessing lists

A list can only be:

- `nil`, or
- the cons of an element onto another list

Use **pattern matching** to access list in one of those ways:

```
let empty lst =  
  match lst with  
  | []      -> true  
  | h::t    -> false
```

*Your brain is probably exploding with AWESOME questions about pattern matching now...*

# Recursion!

Functions over lists are usually recursive: only way to “get to” all the elements

- What should the answer be for the empty list?
- What should the answer be for a non-empty list?
  - Typically in terms of the answer for the tail of the list



# Example list functions

```
let rec sum xs =  
  match xs with  
  | [] -> 0  
  | h::t -> h + sum t
```

```
let rec length xs =  
  match xs with  
  | [] -> 0  
  | h::t -> 1 + length t
```

```
let rec append lst1 lst2 =  
  match lst1 with  
  | [] -> lst2  
  | h::t -> h::(append t lst2)
```

(\* append is available as operator @ \*)



# Lists are immutable

- No way to *mutate* an element of a list
- Instead, build up new lists out of old  
e.g., `::` and `@`

# Match expressions

Syntax:

```
match e with
| p1 -> e1
| p2 -> e2
| ...
| pn -> en
```

the **pi** are *patterns*

the first pipe is optional

line breaks are optional

e.g.,

```
let empty lst =
  match lst with [] -> true | h::t -> false
```

# Patterns

Patterns have their own **syntax**

For now, a pattern can be any of these:

- a variable name (e.g., **x**)
- **[ ]**
- **p1 :: p2**
- an underscore **\_**

As we learn more data structures, we'll learn more patterns

# Patterns

Patterns **match** values

Intuition of matching is that pattern "looks like" the value, if variables in the pattern are replaced by pieces of the value

- `[ ]` looks like `[ ]`
- `h :: t` looks like `2 :: 3`
- `x` looks like `[ 1 ; 2 ; 3 ]`
- `_` looks like anything

...we'll make this precise later

# Match expressions

```
match e with
| p1 -> e1
| p2 -> e2
| ...
| pn -> en
```

## Evaluation:

- Evaluate **e** to a value **v**
- If **p1** matches **v**, then evaluate **e1** to a value **v1** and return **v1**
- Else, if **p2** matches **v**, then evaluate **e2** to a value **v2** and return **v2**
- ...
- Else, if **pn** matches **v**, then evaluate **en** to a value **vn** and return **vn**
- Else, if no patterns match, raise an exception

When evaluating branch expression **ei**, any pattern variables that matched are in scope

Type checker will warn you if you write an *inexhaustive pattern match*

...so you can **prevent exceptions** from being raised at runtime by fixing your code when compiler warns you

# Match expressions

```
match e with
| p1 -> e1
| p2 -> e2
| ...
| pn -> en
```

## Type-checking:

If  $e$  and  $p1 \dots pn$  have type  $\mathbf{ta}$   
and  $e1 \dots en$  have type  $\mathbf{tb}$   
then entire match expression has type  $\mathbf{tb}$

# Pattern matching

The pattern `[]` matches the value `[]` and nothing else

```
match [] with  
| []      -> 0  
| h::t    -> 1  (* evaluates to 0 *)
```

```
match [] with  
| h::t    -> 0  
| []      -> 1  (* evaluates to 1 *)
```

# Pattern matching

The pattern **`h::t`** matches any list with at least one element, and binds that element to **`h`**, and any remaining list to **`t`**

```
match [1;2;3] with  
| []      -> 0  
| h::t    -> h    (* evaluates to 1 *)
```

```
match [1;2;3] with  
| []      -> 0  
| h::t    -> t    (* evaluates to [2;3] *)
```



# A tricky pattern match

```
let rec drop_val v l =  
  match l with  
  | [] -> []  
  | h::t -> let t' = drop_val v t in  
             if h=v then t' else h::t'
```

# Deep pattern matching

- Pattern  $\mathbf{a} :: []$  matches all lists with exactly one element
- Pattern  $\mathbf{a} :: \mathbf{b}$  matches all lists with at least one element
- Pattern  $\mathbf{a} :: \mathbf{b} :: []$  matches all lists with exactly two elements
- Pattern  $\mathbf{a} :: \mathbf{b} :: \mathbf{c} :: \mathbf{d}$  matches all lists with at least three elements

# Accessing lists, with poor style

- Two library functions that return head and tail  
`List.hd`, `List.tl`
- **Not idiomatic** to apply directly to a list
  - Because they throw exceptions; you can easily write buggy code
  - Whereas pattern matching guarantees no exceptions when destructing list; it's hard to write buggy code!

# Tail recursion

```
# length [0; 1; ...; 1_000_000];;  
Stack overflow during evaluation  
(looping recursion?).
```

Why?

```
let rec length xs =  
  match xs with  
  | [] -> 0  
  | h::t -> 1 + length t
```

# Tail recursion

**Solution:** When recursive call is the *only thing left to do* in computation, compiler reuses the stack frame. Reduces space from  $O(n)$  to  $O(1)$ .

```
let rec length_plus_n n = function  
| [] -> n  
| h::t -> length_plus_n (n+1) t  
  
let length_tr = length_plus_n 0
```

# Lists (recap)

- **Syntax:** `[] :: [a; b; c]`
- **Semantics:** building with `nil` and `cons`, accessing with pattern matching
- **Idioms:** recursive functions with pattern for `nil` and for `cons`, **function** syntactic sugar, tail recursion
- **Library:** excellent higher-order functions in OCaml standard library (tomorrow)