

Search engines: ranking algorithms

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 - Ranking Algorithms
 - HITS
 - PageRank



Web Analytics

- Estimated size: up to 45 billion pages
- Indexed pages: at least 2 billion pages
- Average size of a web page is 1.3 MB.... Up 35% in last year
- Petabytes (10^{15} bytes) for storing just the indexed pages!!



Web Analytics

- Web pages are **dynamical** objects
- Most of web pages content changes **daily**
- Average lifespan of a page is around 10 days.

A huge quantity of data!!

Find a document **without** search engines it's like **looking for a needle in a haystack**.



Statistics

Market shares

- <https://www.quantcast.com/top-sites>

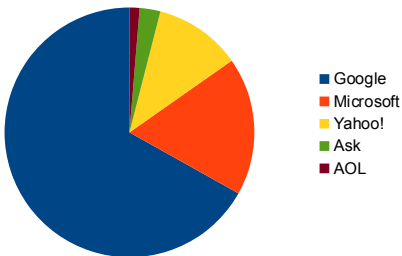


Statistics

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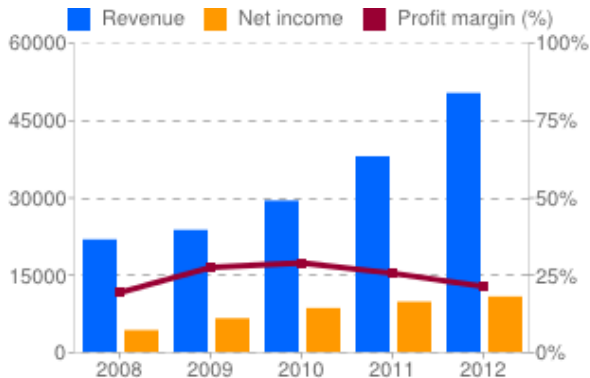
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[comScore, July 2013]



Google's Numbers

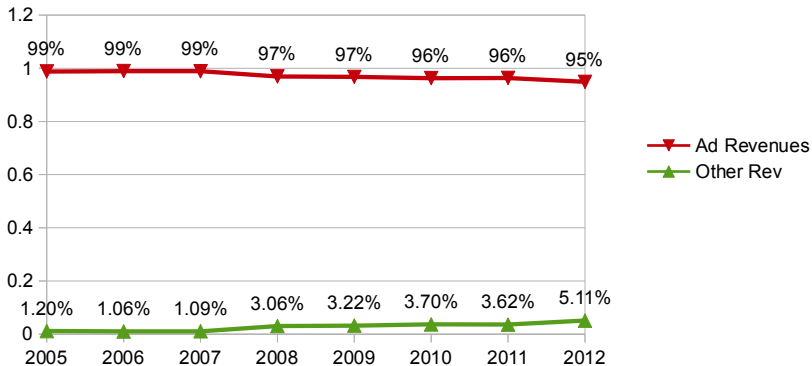


[NASDAQ:GOOG]



Google's Numbers

Revenue Sources



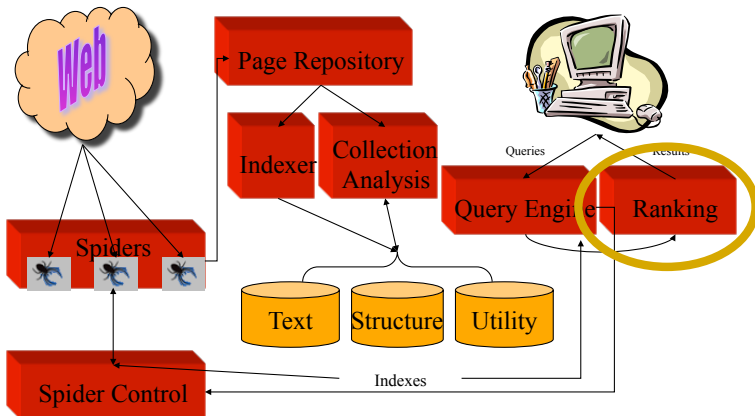
[investor.google.com/financial/tables.html]



Search Engines

- The **dream** of search engines: catalog everything published on the web
- Web content fruition by query
- Rank by relevance to the query

The Anatomy of a Search Engine



... At the very beginning (mid '90s)



The order of pages returned by the engine for a given query was controlled by the **owner** of the page, who could increase the ranking score increasing the frequency of certain terms, font color, font size etc.

SPAM



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SPAM



Modern engines

In 1998 two similar ideas

- HITS (Jon Kleinberg)
- PageRank (S. Brin and L. Page)

The importance of a page should not depend on the owner of the page!



Basic intuition

Look at the linkage structure



Adding a link from page p to page q is a sort of vote from the author of p to q .

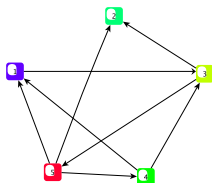
Idea If the content of a page is interesting it will receive many “votes”, i.e. it will have many **inlinks**.



Web Graph

The Web is seen as a directed graph:

- Each **page** is a **node**
- Each **hyperlink** is an **edge**



$$G = \begin{bmatrix} 0 & 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 & 1 \\ 1 & 0 & 1 & 0 & 0 \\ 1 & 1 & 0 & 1 & 0 \end{bmatrix}$$

Web Graph

- The importance of a web page is determined by the structure of the web graph
- At first approximation: Contents of pages is not used!
- Aim: The owner of a page cannot boost the importance of its pages!



HITS

An example

query=**the best automobile makers in the last 4 years**

Intention get back a list of top car brands and their official web sites

Textual ranking pages returned might be very different

Top companies might not even use the terms "automobile makers". They might use the term "car manufacturer" instead



HITS

HITS

Each page has associated two **scores**

a_i **authority** score

h_i **hub** score

A page is a **good** “authority” if it is linked by many good **hubs**

A page is a **good** “hub” if it is linked by many good **authorities**

To every page p we associate two scores:

$$\begin{cases} a_p &= \sum_{i \in \mathcal{I}(p)} h_i \\ h_p &= \sum_{i \in \mathcal{O}(p)} a_i \end{cases}$$

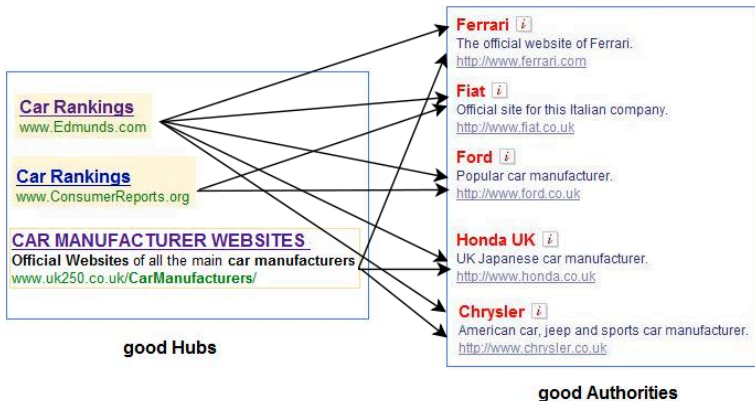
$$\begin{cases} \mathbf{a}^{(i)} &= \mathbf{G}^T \mathbf{h}^{(i-1)} \\ \mathbf{h}^{(i)} &= \mathbf{G} \mathbf{a}^{(i)} \end{cases}$$

Hub pages advertise **authoritative** pages!



HITS

HITS

Query: **Top automobile makers**

HITS

- $G^T G$ and GG^T are non negative
- $G^T G$ and GG^T are semipositive defined
- real nonnegative eigenvalues

$$\begin{cases} \mathbf{h}^{(i)} &= GG^T \mathbf{h}^{(i-1)} \\ \mathbf{a}^{(i)} &= G^T G \mathbf{a}^{(i-1)} \end{cases}$$

\mathbf{h}^* is the dominant eigenvector of GG^T

\mathbf{a}^* is the dominant eigenvector of $G^T G$

For **Perron-Frobenius Theorem** the eigenvector associated to the dominant eigenvalue has **nonnegative** entries!



At query time

- Find **relevant** pages
- Construct the graph starting with those bunch of pages
- Compute dominant eigenvector of $G^T G$
- Sort the dominant eigenvalue and return the pages in accordance with the ordering induced

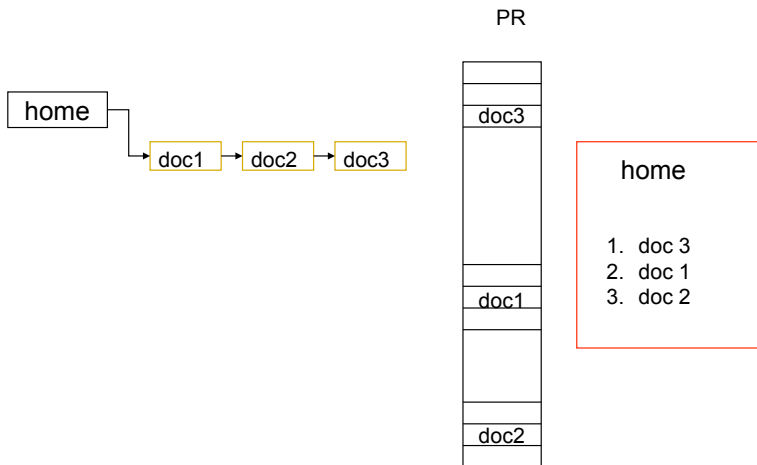


Google's PageRank

- Is a **static** ranking schema
- At query time relevant pages are retrieved neighbours
- The ranking of pages is based on the **PageRank** of pages which is precomputed



PageRank



PageRank

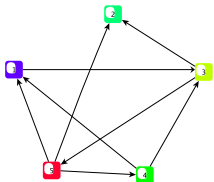
- A page is important if is **voted** by important pages
- The vote is expressed by a link



PageRank

- A page distribute its importance equally to its neighbors
- The importance of a page is the sum of the importances of pages which points to it

$$\pi_j = \sum_{i \in \mathcal{I}(j)} \frac{\pi_i}{\text{outdeg}(i)}$$



$$G = \begin{bmatrix} 0 & 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 & 1 \\ 1 & 0 & 1 & 0 & 0 \\ 1 & 1 & 0 & 1 & 0 \end{bmatrix} \quad P = \begin{bmatrix} 0 & 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 \\ 0 & 1/2 & 0 & 0 & 1/2 \\ 1/2 & 0 & 1/2 & 0 & 0 \\ 1/3 & 1/3 & 0 & 1/3 & 0 \end{bmatrix}$$



PageRank

It is called **Random surfer model**

The web surfer jumps from page to page following hyperlinks. The probability of jumping to a page depends of the number of links in that page.

Starting with a vector $\mathbf{z}^{(0)}$, compute

$$\mathbf{z}_j^{(k)} = \sum_{i \in \mathcal{I}(j)} \mathbf{z}_i^{(k-1)} p_{ij}, \quad p_{ij} = \frac{1}{\text{outdeg}(i)}$$

Equivalent to compute the stationary distribution of the Markov chain with transition matrix P .

Equivalent to compute the left eigenvector of P corresponding to eigenvalue 1.



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PageRank

Two problems:

- Presence of **dangling** nodes
 - P cannot be stochastic
 - P might not have the eigenvalue 1
- Presence of **cycles**
 - P can be reducible: the random surfer is trapped
 - Uniqueness of eigenvalue 1 not guaranteed



Perron-Frobenius Theorem

Let $A \geq 0$ be an irreducible matrix

- there exists an eigenvector equal to the spectral radius of A , with algebraic multiplicity 1
- there exists an eigenvector $\mathbf{x} > \mathbf{0}$ such that $A\mathbf{x} = \rho(A)\mathbf{x}$.
- if $A > 0$, then $\rho(A)$ is the unique eigenvalue with maximum modulo.



PageRank

Presence of **dangling** nodes

$$\bar{P} = P + \mathbf{d}\mathbf{v}^T$$

$$d_i = \begin{cases} 1 & \text{if page } i \text{ is dangling} \\ 0 & \text{otherwise} \end{cases} \quad v_i = 1/n;$$

$$P = \begin{bmatrix} 0 & 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 \\ 0 & 1/2 & 0 & 0 & 1/2 \\ 1/2 & 0 & 1/2 & 0 & 0 \\ 1/3 & 1/3 & 0 & 1/3 & 0 \end{bmatrix} \quad \bar{P} = \begin{bmatrix} 0 & 0 & 1 & 0 & 0 \\ 1/5 & 1/5 & 1/5 & 1/5 & 1/5 \\ 0 & 1/2 & 0 & 0 & 1/2 \\ 1/2 & 0 & 1/2 & 0 & 0 \\ 1/3 & 1/3 & 0 & 1/3 & 0 \end{bmatrix}$$



PageRank

Presence of **cycles**

Force irreducibility by adding artificial arcs chosen by the random surfer with “small probability” α .

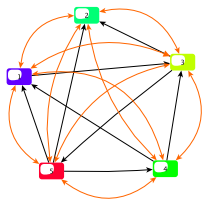
$$\bar{\bar{P}} = (1 - \alpha)\bar{P} + \alpha\mathbf{e}\mathbf{v}^T,$$

$$\bar{\bar{P}} = (1 - \alpha) \begin{bmatrix} 0 & 0 & 1 & 0 & 0 \\ 1/5 & 1/5 & 1/5 & 1/5 & 1/5 \\ 0 & 1/2 & 0 & 0 & 1/2 \\ 1/2 & 0 & 1/2 & 0 & 0 \\ 1/3 & 1/3 & 0 & 1/3 & 0 \end{bmatrix} + \alpha \begin{bmatrix} 1/5 & 1/5 & 1/5 & 1/5 & 1/5 \\ 1/5 & 1/5 & 1/5 & 1/5 & 1/5 \\ 1/5 & 1/5 & 1/5 & 1/5 & 1/5 \\ 1/5 & 1/5 & 1/5 & 1/5 & 1/5 \\ 1/5 & 1/5 & 1/5 & 1/5 & 1/5 \end{bmatrix}.$$

Typical values of α is 0.15.



PageRank: a toy example



$$\bar{P} = \begin{bmatrix} 0.03 & 0.03 & 0.88 & 0.03 & 0.03 \\ 0.2 & 0.2 & 0.2 & 0.2 & 0.2 \\ 0.03 & 0.45 & 0.03 & 0.03 & 0.45 \\ 0.45 & 0.03 & 0.45 & 0.03 & 0.03 \\ 0.31 & 0.31 & 0.03 & 0.31 & 0.03 \end{bmatrix}$$

Computing the largest eigenvector of \bar{P} we get

$$\mathbf{z}^T \approx [0.38, 0.52, 0.59, 0.27, 0.40],$$

which corresponds to the following order of importance of pages

$$[3, 2, 5, 1, 4].$$



PageRank

- P is sparse, $\bar{\bar{P}}$ is full.
- The vector $\mathbf{y} = \bar{\bar{P}}^T \mathbf{x}$, for $\mathbf{x} \geq 0$, such that $\|\mathbf{x}\|_1 = 1$ can be computed as follows

$$\begin{aligned} \mathbf{y} &= (1 - \alpha)P^T \mathbf{x} \\ \gamma &= \|\mathbf{x}\| - \|\mathbf{y}\|, \\ \mathbf{y} &= \mathbf{y} + \gamma \mathbf{v}. \end{aligned}$$

- The eigenvalues of \bar{P} and $\bar{\bar{P}}$ are interrelated:

$$\lambda_1(\bar{P}) = \lambda_1(\bar{\bar{P}}) = 1, \quad \lambda_j(\bar{\bar{P}}) = (1 - \alpha) \lambda_j(\bar{P}), j > 1.$$

- For the web graph $\lambda_2(\bar{\bar{P}}) = (1 - \alpha)$.

