

Principles of Programming Languages [PLP]

Exercises on Scoping, Evaluation strategies, Type checking...

1) Consider the Pascal program to the right:

- a) What is the reference environment at the location in the program indicated by `<== (*)`? That is, give the variables, arguments, and procedures that are visible (in scope) at this location.
- b) The main program calls, P1, P1 calls P3, and P3 calls P2. Draw the stack layout of the subroutine stack after these calls. Show the subroutine frames (without their details) with the static links.
- c) Draw the specific subroutine frame layout of procedure P1, indicating the relevant information that it has to contain.

```
program scopes(input, output)
  procedure P1(A1 : integer)
    var X : integer
    procedure P2(A2 : integer)
      var Y : integer
      begin (* body of P2 *)
        ... <== (*)
      end;
    procedure P3(A3 : integer)
      var X : integer;
      begin (* body of P3 *)
        P2(X)
      end
    begin (* body of P1 *)
      P3(X)
    end
  begin (* body of main program *)
    P1(0)
  end.
```

2) Consider the following outline of a program in a C-like language:

```
int add (int i) { return i + d; }
void p () { const int d = 1;
  print(add(20)); // (1)
}
void q () { const int d = 2;
  print(add(20)); // (2)
}
```

- a) If the language is dynamically scoped, what would be printed at points (1) and (2)?
 - b) If the language is statically scoped, what would happen?
- 3) Show a code fragment in which short-circuit semantics for **or** yield a different result than complete-evaluation semantics.
- 4) For each of the following mechanisms of the **C** programming language, show an example of a type error that can be caused by it:
- a) explicit deallocation of memory,
 - b) union types, and
 - c) pointer arithmetics

- 5) Show what does the program to the right prints if the programming language has:
- static scoping and deep binding
 - dynamic scoping and deep binding
 - static scoping and shallow binding
 - dynamic scoping and shallow binding

In which of the four cases above the functional parameter has to be passed as a closure?

```
int y = 2;
int function f(int function h(int)) {
    int y = 3;
    return h();
}
int function g() {
    int x = y+1;
    return x;
}
int function k() {
    int y = 4;
    return f(g);
}
write(k());
```

- 6) Describe three different ways of allocating in memory a 2-dimensional array A of dimensions $N \times M$.
- Assuming that indexes range in $\{0, \dots, N-1\}$ and $\{0, \dots, M-1\}$ respectively, give the formula for accessing an arbitrary element $A[i][j]$ for each of the three proposed allocation schemes.
 - Translate the formulas into three-address code
- 7) *Innermost* and *outermost* evaluation strategies may require a different number of steps to evaluate an expression. Show how many steps are necessary to evaluate the expression **square ((1 + 2) * 3)** using the rule **square (x) $\rightarrow x * x$** and the obvious rules for addition and multiplication, by using:
- innermost (applicative) evaluation
 - outermost (normal order) evaluation
 - outermost evaluation with memoization

8) The ML function **main** to the right computes Fibonacci numbers in a nonstandard way (see Exercise 7.3.1 page 451 of the Dragon book).

Show the stack of activation records that result from a call to **main**, up until the time that the first call (to **fib0(1)**) is about to return. Show the access link in each of the activation records on the stack.

```
fun main () {
  let
    fun fib0(n) =
      let
        fun fib1(n) =
          let
            fun fib2(n) = fib1(n-1) + fib1(n-2)
          in
            if n >= 4 then fib2(n)
            else fib0(n-1) + fib0(n-2)
          end
        in
          if n >= 2 then fib1(n)
          else 1
        end
      in
        fib0(4)
      end;
}
```