

Principles of Programming Languages [PLP]

Exercises on Parsing

1) Consider the following grammar G :

$$R \rightarrow (R) \mid R + R \mid RR \mid R^* \mid a$$

- Provide a leftmost and a rightmost derivation for string $(a+a)^*a$
- Draw a parse tree for string $(a+a)^*a$
- Describe the language generated by grammar G
- Draw an abstract syntax tree of string $(a+a)^*a$
- Is grammar G ambiguous? Is string $(a+a)^*a$ ambiguous?
- Transform grammar G by left-factorizing it: call $LF(G)$ the resulting grammar
- Draw a parse tree for the string $(a+a)^*a$ with respect to grammar $LF(G)$
- Transform grammar G by eliminating left recursion, obtaining grammar $LRE(G)$
- Draw a parse tree for the string $(a+a)^*a$ with respect to grammar $LRE(G)$

2) Consider the following grammar over the set of terminal symbols $\{id, ", +\}$:

$$\begin{aligned} S &\rightarrow id \mid " T " \\ T &\rightarrow S V \\ V &\rightarrow \varepsilon \mid + S V \end{aligned}$$

- Show $First(\alpha)$ for each production $X \rightarrow \alpha$ and $Follow(A)$ for each non-terminal A
- Build the LL(1) parse table
- Starting from the configuration (**stack:** $S \$$, **input:** $" id + id " \$$), show the evolution of the stack and of the input **in the first six steps** of the top-down predictive parsing algorithm using the LL(1) parse table. (Note: the top of the stack is to the left.)

3) Given grammar $A \rightarrow A A + \mid a$

- Is string $Aa+A+$ a sentential form? Is it a right-sentential form?
- Which is the handle in string $AA+a+$?

4) Given grammar $E \rightarrow E + E \mid x$

- Is the following claim true or false? Motivate your answer.
"In string $E+E+x$, both $E + E$ and x are handles."

5) Consider grammar

$$\begin{aligned} S &\rightarrow (A) \mid x \\ A &\rightarrow A + S \mid S \end{aligned}$$

- Show $\text{First}(\alpha)$ for each production $X \rightarrow \alpha$ and $\text{Follow}(X)$ for each non-terminal X
- Is the grammar LL(1)? Justify your answer
- Draw the LR(0) automaton of the grammar
- Draw the SLR parsing table of the grammar
- Starting from the configuration (**stack:** $[0]$, **input:** $[x + x]$), show the evolution of the stack and of the input **in the first six steps** of the bottom-up LR parsing algorithm using the SLR(1) parse table. (Note: the "0" in the stack represents the start state of the LR(0) automaton.)

6) Consider the following grammar, whose terminals are $\{a, ?\}$:

$$\begin{aligned} S &\rightarrow A \\ A &\rightarrow B \mid BA \\ B &\rightarrow a ? C \\ C &\rightarrow \varepsilon \mid a C \end{aligned}$$

- Left-factor the grammar,
- Compute the $\text{First}(A)$ and $\text{Follow}(a)$ sets for each production $A \rightarrow \alpha$ of the resulting grammar.
- Build the LL(1) parse table.
- Explain why the grammar is not LL(1).
- Show that the language is LL(2), arguing convincingly that the conflicts can be resolved by looking ahead one more token.

7) Consider the grammar:

$$\begin{aligned} A &\rightarrow CaBa \\ A &\rightarrow B \\ B &\rightarrow C \\ C &\rightarrow b \end{aligned}$$

- What is the language generated by the grammar? Is it ambiguous?
- Construct the LR(0) automaton
- Build the SLR parse table. Is the grammar SLR?
- Construct the LR(1) automaton and the LR(1) parsing table. Is the grammar LR(1)?