

Principles of Programming Languages

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Lesson 14

- Introduction to Denotational semantics

Describing a Programming Language

- Syntax, **semantics** and pragmatics
- Semantics defines the meaning of programs
- Various kinds of semantics
 - Operational
 - Algebraic / Axiomatic
 - **Denotational**
 - Game Theoretical
 - ...
- Used for:
 - Unambiguous specification of meaning of programs
 - Correctness of implementations
 - Proving properties or equivalence of programs
 - Evaluating alternative constructs in design phase

Denotational Semantics

- Developed by Dana Scott and Christopher Strachey (~1970)
- Topic of course MOD (Models of Computations) [summer semester]
 - Mathematical foundations (Domain theory)
 - Complete semantics of simple programming languages: IMP (imperative) and HOFL (functional)
- Our presentation is orthogonal
 - Foundations: almost none and informally
 - “Descriptive” use of semantics:
 - to understand programming constructs
 - to compare them across different programming languages
 - We follow
“R.D. Tennent: The denotational semantics of programming languages, Communications of the ACM, Volume 19 Issue 8, Aug. 1976”

Basics and Syntax of LOOP

- The **denotational semantics** of a programming language map programs to mathematical objects (**denotations**) representing the *meaning* of the programs
- This is done **compositionally** on the syntax of the program
- The **abstract syntax** of a language defines
 - a collection of *syntactic domains*, corresponding to non-terminal symbols
 - Example: **Prog, Exp, Com, Var, ...**
 - a collection of *operations* on syntactic domains corresponding to productions

Abstract syntax of the LOOP language [Tennent76]

$\text{Exp} ::= 0 \mid \mathbf{succ} \text{ Exp} \mid \text{Var}$

$\text{Com} ::= \text{Var} := \text{Exp} \mid \text{Com} ; \text{Com} \mid \mathbf{to} \text{ Exp} \mathbf{do} \text{ Com}$

$\text{Prog} ::= \mathbf{read} \text{ Var} ; \text{Com} ; \mathbf{write} \text{ Exp}$

- A LOOP program computes a function on natural numbers

Productions as operations

$0: \rightarrow \text{Exp}$

$\mathit{succ}: \text{Exp} \rightarrow \text{Exp}$

$\mathit{in}: \text{Var} \rightarrow \text{Exp}$

$\mathit{seq}: \text{Com} \times \text{Com} \rightarrow \text{Com}$

$\mathit{assign}: \text{Var} \times \text{Exp} \rightarrow \text{Com}$

$\mathit{rep}: \text{Exp} \times \text{Com} \rightarrow \text{Com}$

$\mathit{prog}: \text{Var} \times \text{Com} \times \text{Exp} \rightarrow \text{Prog}$

Denotational Semantics of LOOP

For each *syntactic domain* a corresponding *semantic domain* is defined, and the meaning is given by a *semantic interpretation function*

- $P : \text{Prog} \rightarrow N \rightarrow N$ (\rightarrow associates right, read “Prog \rightarrow (N \rightarrow N)”)
 - The meaning of a program is a function from N to N
- Since $\text{Prog} ::= \text{read Var} ; \text{Com} ; \text{write Exp}$, to define P compositionally we need the semantics of Var, Com and Exp
- For evaluating variables, we introduce the domain of *states*:
 - $S = \text{Var} \rightarrow N$ thus a state $s \in S$ is a function from Var to N
 - for a state $s \in S$, $s\{v\}$ is the content of variable v
 - Def: $s[n/v]$ is a state s.t. $s[n/v]\{x\} = n$ if $v = x$, else $s[n/v]\{x\} = s\{x\}$

Note: we use $\{_ \}$ instead of $[[_]]$, the classical notation

Denotational Semantics of LOOP: Expressions and Commands

We define E and C by induction:

- $E: \text{Exp} \rightarrow S \rightarrow \mathbb{N}$
 - $E\{0\} s = 0$
 - $E\{\text{succ } e\} s = E\{e\} s + 1$
 - $E\{v\} s = s(v)$

Abstract syntax of the LOOP language [Tennent76]
Exp ::= 0 | **succ** Exp | Var
Com ::= Var := Exp | Com ; Com | **to** Exp **do** Com
Prog ::= **read** Var ; Com ; **write** Exp

Commands change the state:

- $C: \text{Com} \rightarrow S \rightarrow S$
 - $C\{v := e\} s = s[n/v]$ where $n = E\{e\} s$
 - $C\{c_1; c_2\} s = (C\{c_2\} \circ C\{c_1\})s$ [= $C\{c_2\} (C\{c_1\} s)$]
 - $C\{\text{to } e \text{ do } c\} s = ((C\{c\})^n) s$ where $n = E\{e\} s$
- Note: $f^0(x) = x$, $f^{n+1}(x) = f(f^n(x))$

Denotational Semantics of LOOP: Programs

- $P : \text{Prog} \rightarrow N \rightarrow N$
 - *The meaning of a program is a function from N to N*
- $P\{\mathbf{read } v ; c ; \mathbf{write } e\}n = E\{e\}s$
 - where $s = C\{c\}s_0[n/v]$
 - where $s_0\{w\} = 0$ for each variable w

Just a simple example to stress compositionality

- Language LOOP is total
- There is no conditionally controlled iteration
- More complex domains are needed in general