From Firewalls to Functions and Back

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What is a Firewall?



Inspects the traffic on a node of the network, for each packet

- accepts or drops it
- possibly changes the addresses (NAT)

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Based on a configuration

- List of rules
- Possibly using tags
- Procedure-like constructs
- Interaction among rules (Shadowing)

Firewalls are a basic tool for protecting network

- Widespread
- Configuration-based
- Different configuration languages (iptables, pf, ipfw)
- It's Hard to configure and manage firewalls
- Cross-platform policy porting is Harder

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Misconfigurations cause unintended behaviour **Possible Threats**

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Previous works:

Transcompilation Pipeline between firewall languages

- Supports iptables, pf, ipfw and (partially) CISCO-ios
- General approach
- Supports NAT
- Formal semantics

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Why

- For automated policy porting (first general approach!)
- For configuration refactoring
- Synthesis of a high level declarative configuration
- Basis for other policy management tasks

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Old Legacy Technology

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Old Legacy Technology	iptables ipfw pf
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Each firewall system

- Has its own configuration language
- Makes different evaluation steps to process packets
- Lots of low level details
 - First do the NAT, than filtering or vice-versa?
 - How to express complex conditions (negated)?

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$\label{eq:Firewall} \mbox{Firewall} = \mbox{set of rules} + \mbox{the evaluating procedure}$

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Control Diagram



 $\ensuremath{\mathcal{S}}$ are the addresses of the firewall

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Configuration

Assigns a rulesets to each node

Ruleset : list of rules $r = (\phi, a)$

- $\phi(p)$: condition
- *a* : action
 - ACCEPT
 - DROP
 - NAT (d_n, s_n)
 - MARK(m)
 - GOTO(*R*)
 - CALL(R)
 - RETURN





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Transcompilation Pipeline



From Firewalls to Functions and Back: The Idea

Previous implementation of the pipeline synthesis:

Compute the models of a predicate (SAT-solver) Black-box approach (no fine tuning)

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\begin{array}{ll} (\mathsf{Firewalls} \twoheadrightarrow \mathsf{Functions}): & & \\ & & \mathsf{source\ configuration} \mapsto \mathsf{function\ representing\ its\ \textbf{meaning}} \\ (\mathsf{Firewalls} \twoheadleftarrow \mathsf{Functions}): & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & \\ & & & & \\ & & & & \\ & & & & \\ & & & & \\ & & & & \\ & & & & \\ & & & & \\ & & & & \\ & & & & \\ & & & & \\ & & & & \\ & & & & \\ & & & & \\ & & & & \\ & & & & \\ & & & & \\ & & & & \\ & & & & \\ & & & & \\ & & & & \\ & & & & \\ & & & & \\ & & & & \\ & & & & \\ & & & & \\ & & & & \\ & &
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Change of domain:

Function-based redefinition of the pipeline

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(Firewalls → Functions) :
source configuration → function representing its meaning
(Firewalls ← Functions) :
functional representation → target configuration
```

Functions are an handy domain: They allow simple and general solutions

Rulesets and Firewalls as Functions

- $\ensuremath{\mathbb{P}}$ network packets
- $\mathcal{T}(\mathbb{P})$ transformations possibly applied to packets
 - $\perp\,$ discard of a packet

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New pipeline stages:

- ruleset synthesis: rulesets became functions
- composition: computes the semantics of the firewall
- generation: assign functions to the target nodes
- translation: from IFCL to pf configuration language



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Why:

- Parametric w.r.t. IFCL specification
- Support minimal control diagrams and MARK
- Translation from IFCL to target language is trivial



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Function Representation

Functions $\tau: \mathbb{P} \to \mathcal{T}(\mathbb{P}) \cup \{\bot\}$ as sets of pairs (P, t)

- $t\;$ is a transformation
- $\boldsymbol{P}~$ is a multi-cube of packets

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Cube :

Cartesian product of one segment for each dimension



Multi-cube :

Cartesian product of one **union of segments** for each dimension



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Cube :

Cartesian product of one segment for each dimension

Multi-cube :

Cartesian product of one **union of segments** for each dimension

- succinct representation
- sets of packets verifying rule conditions
- sets of packets verifying arc conditions
- closed under transformations





Synthesis



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Ruleset Synthesis

From a ruleset to a set of pairs (P, t)

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We scan the ruleset rule-by-rule, keeping track of

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Base Case: if R = [] just returns $\{(P, t)\}$

From a ruleset to a set of pairs (P, t)

We scan the ruleset rule-by-rule, keeping track of

- ${\cal P}\,$ packets not managed
- $t\,$ transformation assigned to P

Base Case: if R = [] just returns $\{(P, t)\}$

Else: take the first rule $(\phi, action)$

 $P = \begin{cases} P_s & \text{packets that verifies } \phi \\ P_n & \text{packets that do not - managed by the$ **other rules** $} \end{cases}$

if *action* terminates the packet processing then (P_s, t') else P_s also managed by the other rules (updated transformation t')

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Composition



Ideally, for each $p \in \mathbb{P}$

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 \bullet compute t in the first node

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- \bullet compute t in the first node
- compute p': (how p is when exits node q)

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- compute t in the first node
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- \bullet check $\psi(p')...$ if ti does then
 - compute t' in the second node

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Globally $p \mapsto t$ updated with t'

- compute t in the first node
- compute p': (how p is when exits node q)
- \bullet check $\psi(p')...$ if ti does then
 - $\bullet \ \mbox{compute} \ t' \ \mbox{in the second node}$
 - Overall: $p \mapsto t$ updated by t'



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- compute p': (how p is when exits node q)
- \bullet check $\psi(p')...$ if ti does then
 - compute t' in the second node
 - Overall: $p \mapsto t$ updated by t'

Composition Algorithm:

The same,

but with Multi-cubes ...

(... plus details)

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Example from ipfw to pf: formalization

ipfw -q nat 1 config ip 151.15.185.183 ipfw -q nat 2 config redirect_port tcp 9.9.8.8:17 17 ipfw -q add 0010 nat 1 tcp from 192.168.0.0/24 to not 192.168.0.0/24 ipfw -q add 0020 nat 2 tcp from 151.15.185.183 to not 192.168.0.0/24 17 ipfw -q add 0030 allow tcp from 151.15.185.183 to not 192.168.0.0/24 out ipfw -q add 0040 deny all from any to any

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Example from ipfw to pf: ruleset synthesis

$$\begin{split} R_0: (\texttt{sIP} \in 192.168.0.0/24 \land \texttt{dIP} \notin 192.168.0.0/24, \texttt{NAT}(\star: \star, 151.15.15.183: \star)); \\ (\texttt{sIP} = 151.15.185.183 \land \texttt{dIP} \notin 192.168.0.0/24 \land \texttt{dPort} = 17, \texttt{NAT}(9.9.8.8: \star, \star: \star)); \\ (true, \texttt{DROP}) \end{split}$$

Example from ipfw to pf: ruleset synthesis

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$ au_0$								
Received packets			Accepted packets					
source		destination		source		destination		
192.168.0.0/24	*	* \{ 192.168.0.0/24 }	*	151.15.185.183	-	-	-	
151.15.185.183	*	* \{ 192.168.0.0/24 }	17	-	-	9.9.8.8	-	

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Example from ipfw to pf: composition



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Example from ipfw to pf: composition



Received packets			Accepted packets				
source		destination		source		destination	
151.15.185.183	*	* \{ 151.15.185.183 192.168.0.0/24 }	* \{17}	-	-	-	-
192.168.0.0/24 \ {192.168.0.1}	*	127.0.0.1 151.15.185.183	*	151.15.185.183	-	-	-
192.168.0.0/24 \ {192.168.0.1}	*	* \ { 127.0.0.1 151.15.185.183 192.168.0.0/24 }	* \{17}	151.15.185.183	-	-	-
192.168.0.0/24 \ {192.168.0.1}	*	* \ { 127.0.0.1 151.15.185.183 192.168.0.0/24 }	17	151.15.185.183	-	9.9.8.8	-
192.168.0.1	*	* \ { 127.0.0.1 151.15.185.183 192.168.0.0/24 }	*	151.15.185.183	-	-	-
151.15.185.183	*	* \ { 192.168.0.0/24 }	17	-	-	9.9.8.8	-

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Generation



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How to generate functions

Problem: not every ruleset can be assigned to each node!

• To guarantee the final translation

- Simple targets: ACCEPT, DROP and NAT
- Assign Labels to nodes:
 - DROP
 - SNAT
 - DNAT
- Different expressive power

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Algorithm

- For each pair (P,t) with $t \neq \bot$
 - Find the path
 - For each node q
 - $\bullet~\mathsf{Preceding}~\mathsf{nodes}\to\mathbf{P_q}$
 - $\bullet~\mathsf{Labels}$ in $q\to \mathbf{t_q}$
- Special management for DROP pairs (P, \bot)
 - For each node: packets still not managed
 - Drop as many of these as possible



Conclusion

The presented transcompilation approach

- Is parametric w.r.t. the IFCL specification
- Supports the use of tags
- Supports firewalls with minimal control diagram
- Preserves the NAT
- Reveals different expressive power of firewall languages

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Ongoing and Future Works

- Coding and Testing
- Non-trivial multi-cube merging procedure
- Support for holistic network management
- High-level tools for network management, compatible with old technology

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