Compilers
(Compiling Techniques)

Programs

Manipulating

Programs According to the Meaning

Programs manipulating Programs

What are meaning the following writings?

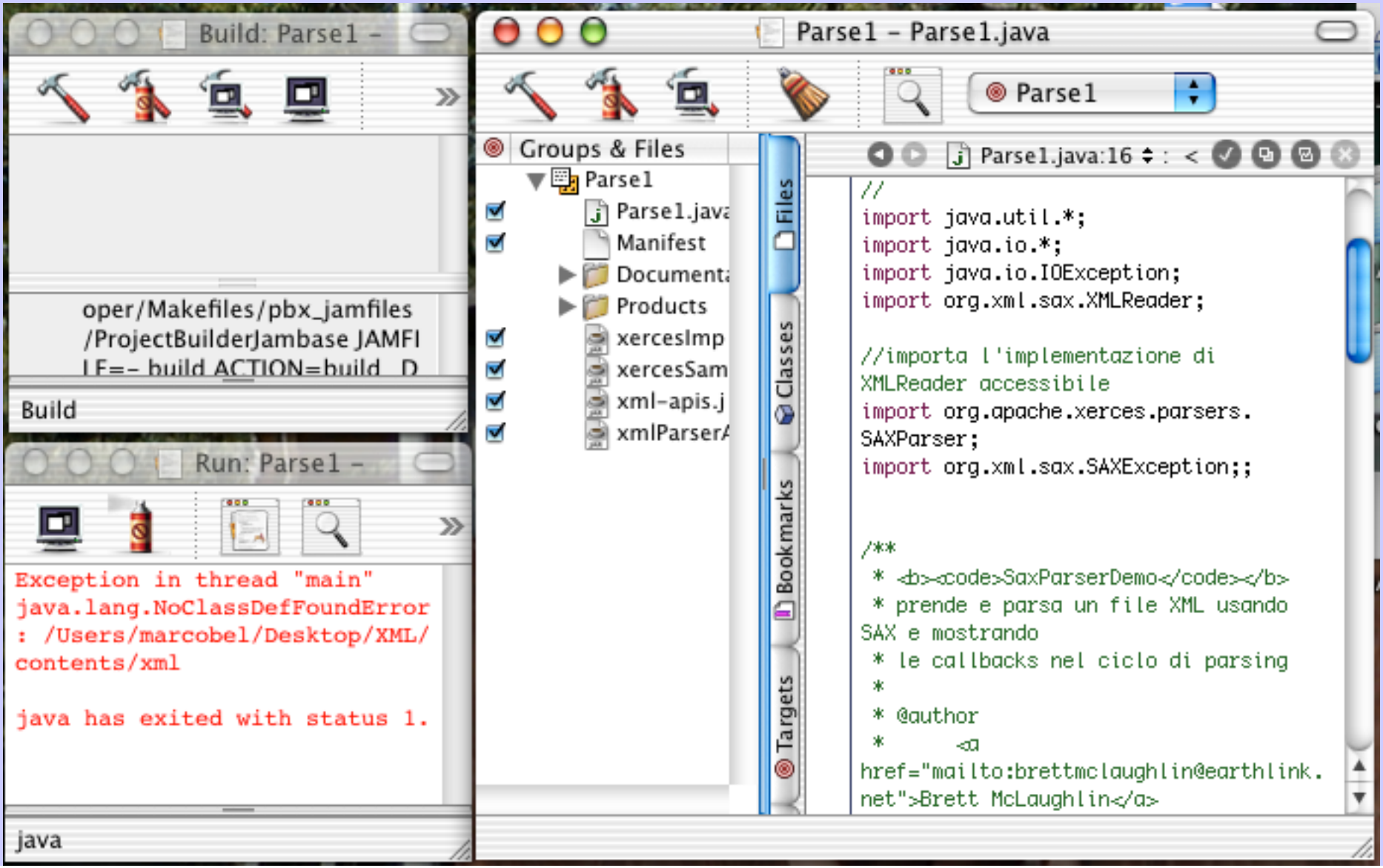
- cc -w -o code.exe ex.yy.c
- code.exe A.file B.file <In >Out

- javac -classpath /docu/XML/Xerces/Xparse.jar -v MyDoc.java
- java MyDoc

- **What are** ex.yy.c, MyDoc.java **?**
- **What are** cc, javac, java **?**

- cc, javac, java **are programs manipulating programs:**
  - **How are defined (built, obtained) ?**
  - **Where are running?**
  - **In what language are they written?**

Compiler I/0 may be inline or through a graphic interface



Many different I/0 interaction structures exist but they are unaffecting

- The (compiler) beaviour, hence:
  - its Construction Principles
  - its Construction Techniques
  - its Internal Structure

In this Course: The Techniques

Basic Techniques for the construction of **Tools for Abstract Machines**

Techniques: Automata (N/D Finite State, T/B Pushdown)
Syntax-Directed Translations
Attribute Grammars/Translation Scheme
Structure Traversal / Visit
Translation Invariants (code generation)
…

In This Course:
Methodologies, Tools

Methodologies: Semantic Attachment
Abstract Interpretation
Meta-evaluation
Partial Evaluation
…

Tools: Lexical/Syntactic Analysers
Semantic Analisers
Syntactic Editors/Text Formatters
(analysis) Tools Generators
Code Generators
Interpreters / Compilers
Debuggers
Code Optimization Tools…

Foundations

- Language and Abstract Machine (MA)
- MA: Structure and Computation (States)
- Construction of MAs: Interpreter, Compiler
- Interpreter: Inside
- Compiler: Run Time Support (RTS)
- Compiler: Development Machine,
  Hierarchy Source-Host-Target (SOT)
- Intermediate Machines: Mixed Constructions

Definitions:

Language and Formalism

Language (Programming L.) =

= Formalism that Rigorously Expresses
(applications of computable functions)

Formalism = **Syntax** (*form* of the allowed sentences)
+
Semantics (*meaning* of each sentence)

Example: Programming Language

L = **<S, SEM>** is an LP

1) $\forall P \in S$, $SEM(P) \in \{N \rightarrow N\}$

2) $\forall g \in \{N \rightarrow N\}, \exists P \in S$, such that:
$g(n) = SEM(P)(n)$ $(\forall n \in N)$

[Let $\{N \rightarrow N\}$ be the set of the Computable Functions]

Definitions:

Abstract Machine

MA =

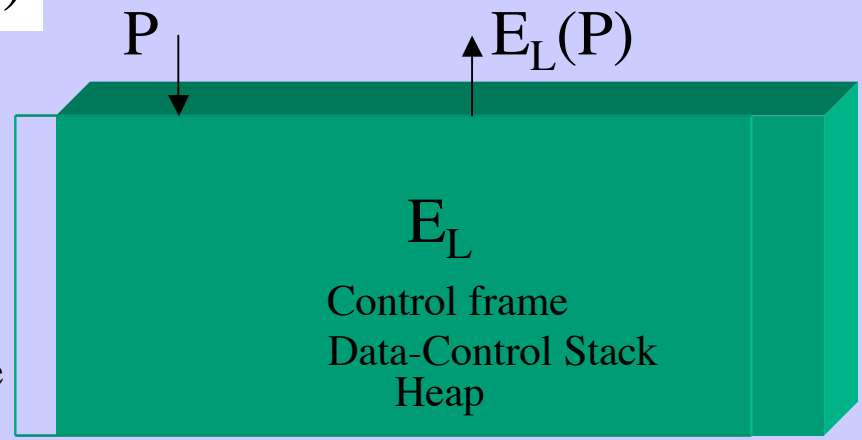
Machine Language (L=<S,SEM>)

+

Machine Executor (E_L)

$\forall P \in S,$

$SEM(P) \approx E_L(P)$



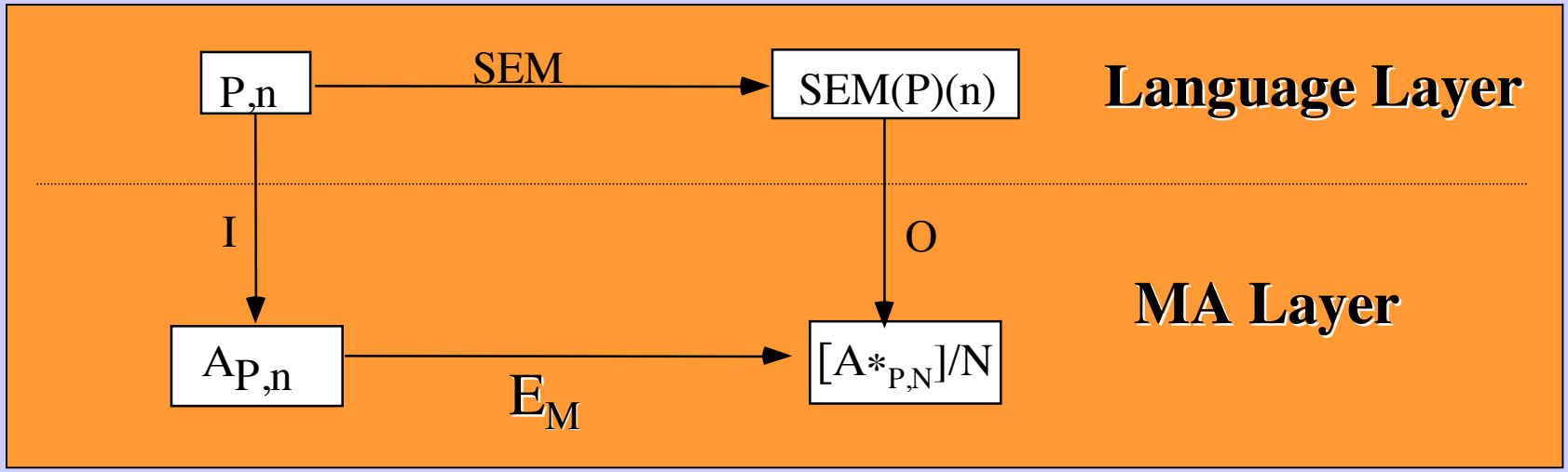
Java Virtual Machine
Landin's SECD

Example: MA, Language and Machine Executor

- L_M = <S, SEM> is an LP
- M = <L_M,E_M>
- E_M is the executor of M

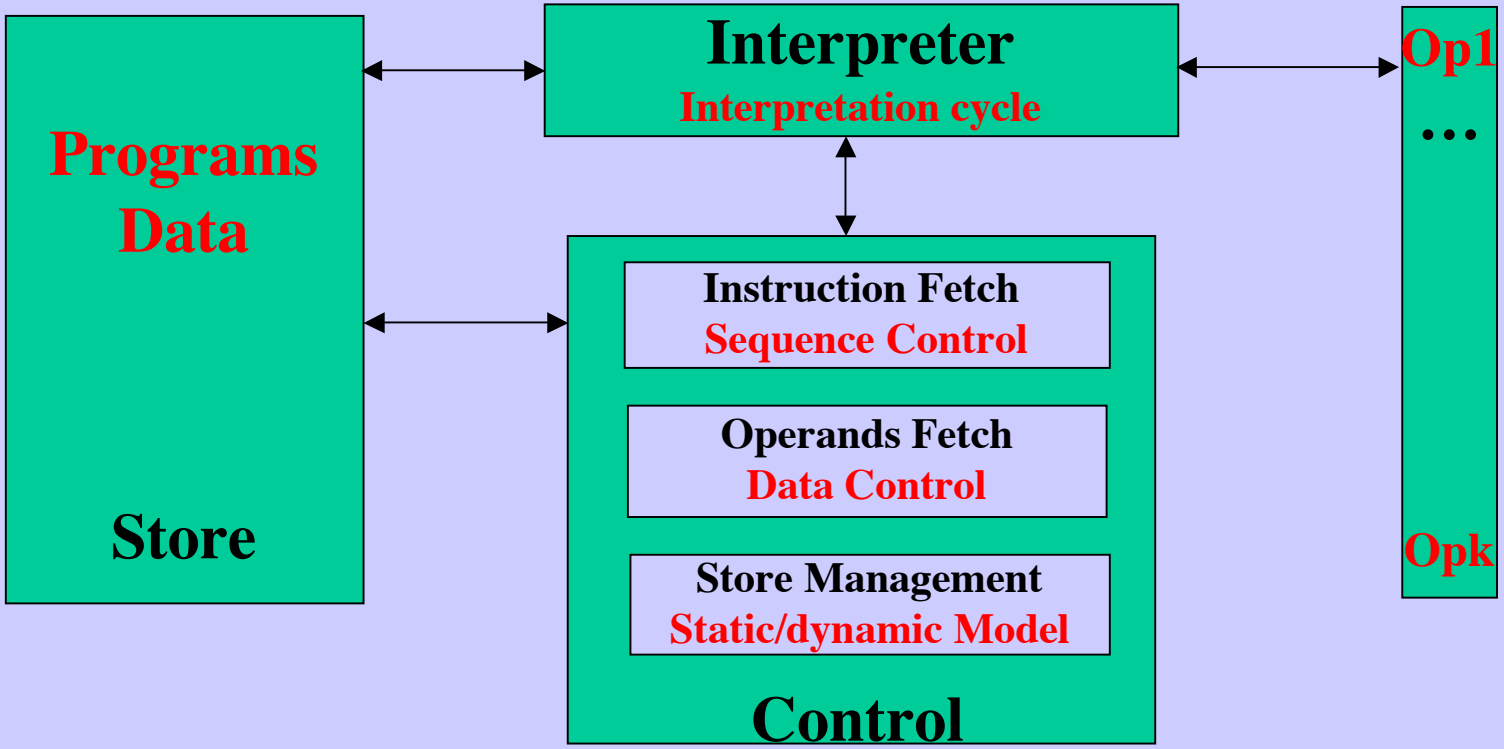
E_M: A → A*

$\forall P \in S,\ n \in N$



[where: I,O injections on A e A* risp., […]/N normalization on N]

MA: Structure e Executor States



Abstract Machine - Machine Structure

MA:
Store, Control

Store: **It is structured according to a model that relies on the specific features of the Machine Language**

- ***arrays of words, registers, stacks***
- **heap -** for dynamic allocation (Pascal, C, C++, …,Java,
- ***graph -*** for structure sharing (*functional languages*)

Control: **It handles the Executor States:**

- **finds the next *statement* or expression**
- **finds the *stat.* or *espr.* data**
- **updates store**

MA:
Elementary Execution Cycle