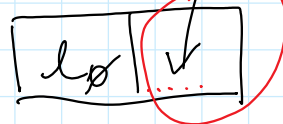


Semantica formale del linguaggio imperativo

int * x;



$$v : Loc \rightarrow Val \perp$$

$$v : Loc \rightarrow \underline{(Val \cup Loc)} \perp$$

localizzazione

$Val = \mathbb{N}$

Valori,
ma anche
indirizzi

nuovo tipo

Exp \rightarrow ... | * Ide | & Ide

Dec \rightarrow ... | ~~Type * Ide ;~~ |

Com \rightarrow ... | * Ide = Exp ;

non sense

Type \rightarrow int | float | int * | float *

$$Sem_e: Exp \rightarrow P \rightarrow M \rightarrow (Val \cup Loc)_{\perp}$$

$$Sem_e (*x) \rho \mu = \mu \left(\underbrace{\mu(\rho x)}_{e_x} \right)$$

$x \in Id$

$$Sem_e (lx) \rho \mu = \rho(x)$$

$$Sem_c: Com \rightarrow P \rightarrow M \rightarrow M$$

$$Sem_c (*x = e;) \rho \mu = \mu \left[\begin{matrix} v \\ \mu(\rho x) \end{matrix} \right]_{mod} \quad \begin{matrix} x \in Id \\ e \in EXP \end{matrix}$$

dove

$$v = Sem_e e \rho \mu$$

Cancellare gli element ripetuti più di una volta in una lista lasciando una sola occorrenza. (conc)

$$\text{conc } [2; 3; 5; 2; 7; 3; 2] = [2; 3; 5; 7]$$

Ci serve la funzione

member : 'a → 'a list → bool

pensiamole in modo ricorsivo

let rec member a l = match l with
[] → false

| x :: xs → if a = x then true
else member a xs;;

Ha senso definirle con un accumulatore?

No! perché non appena troviamo l'elemento
o arriviamo alle liste vuote

Sappiamo dove il risultato senza fare
ulteriori operazioni.

Con exists

let rec exists p l = match l with
 [] → false
 | x::xs → if p x then true
 else exists p xs ;;

let member a l =

let p x = x = a

def. ↗

↖ uguaglianza

in exists p l ;;

let member a l =

let f x y = if x = a then true
 else y

in foldl f false l ;;

Cancellare le occorrenze multiple degli element di una lista

Ricorsione pura

let rec conc l = match l with
[] -> []

| x :: xs -> if member x xs
then conc xs
else x :: conc xs ;;

conc : 'a list -> 'a list = <fun>

non è "tail recursive"

Con accumulatore

giovedì 7 dicembre 2017 16:48

let concat l =

let rec concat l a = match l with
[] → a

| x :: xs → if member x xs

then concat xs a

else concat xs (x :: a)

in concat l [];;

Conc utilizzando folbr

giovedì 7 dicembre 2017 16:52

let conc l =

let f x y = if member x y
then y
else x::y

in folbr f [] l ;

Date una lista restituire la coppia

giovedì 7 dicembre 2017 16:55

(valore massimo, numero di occorrenze di questo valore)

$\text{max } [2; 7; 3; 7; -2; -2; 7] = (7, 3)$

↑ (←) ↑ ↑
Minimo n. di occorrenze

let rec max l = match l with
[x] → (x, 1)

| x :: y :: ys → let (m, n) = max (y :: ys)

in

if x = m then (m, m+1)
else if x > m then (x, 1)
else (m, n);;

non è tail recursive

max : 'a list → 'a * int = <fun>

Con accumulatore

giovedì 7 dicembre 2017 17:03

let max l =

let maxa l (m, m) = match l with

[] → (m, m)

| x :: xs → if x > m
then maxa xs (x, 1)
else if x = m
then maxa xs (m, m+1)
else maxa xs (m, m)

in maxa (tl l) (hd l, 1) :: ; ;

max [2; -1; 2]
= { def max }
maxa [-1; 2] (2, 1) ^a
= { def max a } _{2° p}

$$\text{mexa } [2] \ (2,1) \\ = \{ \text{def mexa}, 2^{\circ} p \}$$

$$\text{mexa } [] \ (2,2) \\ = \{ \text{def mexa}, 1^{\circ} p \}$$

$$(2,2) \leftarrow$$

$$\text{mex } [2]_{ii} \\ = \{ \text{def mex} \}$$

$$\text{maxa } [] \ (2,1) \\ = \{ \text{def. mexa}, 1^{\circ} p \}$$

$$(2,1)$$

$$\text{max } [3;7] \\ = \{ \text{def. mex} \}$$

$$\text{mexa } [7] \ (3,1) \\ = \{ \text{def}$$

let ~~max~~ l =

giovedì 7 dicembre 2017 17:19

let ~~max~~ l (m, n) = match l with
[] → (m, n)

| x::xs → if x > m then max xs (x, 1)
else if x = m

then max xs (m, n+1)

else max xs (m, n)

in max (tl l) (hd l, 1) ;;

≡

match l with

x::xs → max xs (x, 1)

let max l =

let f x y =

let (m,m) = y in ...

ok

in

let max l =

let f x (m,m) =

if x > m then (x,1)

else if x = m then (m,m+1)

else (m,m)

in

foldr f (hd l, 1) (tl l)

≡ match l with

x :: xs → foldr f (x, 1) xs

$x :: xs \rightarrow \text{fold } f \left(\begin{array}{c} \swarrow \\ (x, 1) \\ \searrow \end{array} \right) xs$