

$$\mathcal{L} = \{a, b\}$$

$\alpha$  ni dice sottostringa di  $\beta$

$\alpha \subseteq \beta$  sse  $\beta$  contiene la stringa  $\alpha$

$abaaab \subseteq ababaaabbb$

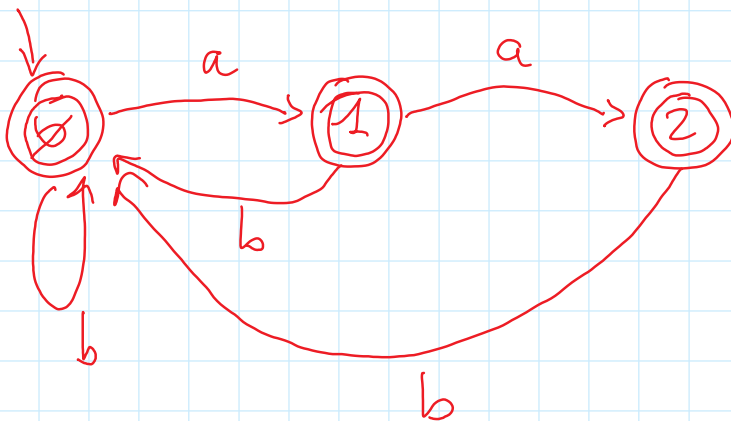
$\alpha \in \mathcal{L}^*$

$$L = \{ \alpha \mid \text{aaa} \notin \alpha \}$$

$aba \in L$

$abbaabba \in L$

$\underline{\underline{bbaaabaa}} \notin L$



$\bar{\epsilon}$  regole

$$L = \{ a^k b^m c^m \mid k+m > m > \emptyset \}$$

$\mathcal{A} = \{a, b, c\}$

$abc \in L$   
 $abbc \notin L$   
 $aabbc \in L$

$$L = \{ a^m b^m b^k c^s \mid (m \geq m \wedge s > k) \vee (m > m \wedge s \geq k) \}$$

$$S \rightarrow AB \mid CD$$

$$A \rightarrow \epsilon \mid aAb \mid aA$$

$$B \rightarrow c \mid bBc \mid Bc$$

$$C \rightarrow a \mid aCb \mid aC$$

$$D \rightarrow \epsilon \mid bDc \mid Dc$$

$$L_A = \{ a^m b^m \mid m > m > \emptyset \}$$

$$L_B = \{ b^m c^m \mid m > m \geq \emptyset \}$$

Verificare che in un array  $a$  di dimensione  $dim$ , ogni valore sia seguito da almeno un valore minore

(tra quelli che seguono ce n'è almeno uno minore)

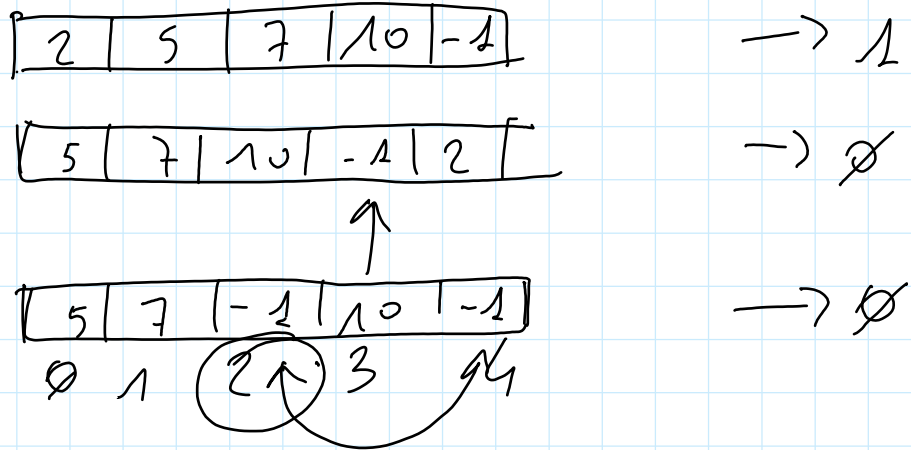
```
int trovamino (int el, int a[], int inizio,
               int fine)
```

```
{ int trovato = 0;
  int i = inizio;
  while (i < fine && ! trovato)
    if (a[i] < el) trovato = 1;
    else i++;
  return trovato;
}
```

```
int check (int a[], int dim)
```

```
{ int ok = 1;
  int i = 0;
  while (i < dim - 1 && ok)
    if (! trovamino (a[i], a, i + 1, dim))
```

```
if (!trovaminore(a[i], a, i+1, dim))
    OK = 0;
else i++;
return OK;
}
```



Cerchiamo l'indice del più piccolo

```
int indmin (int a[], int dim);
{
  int ind = ∅;
```

```
  int i;
```

```
  for (i = 1; i < dim; i++)
```

```
    if (a[i] < a[ind]) ind = i;
```

```
  return ind;
```

```
}
int check (int a[], int dim)
```

```
{
  return indmin (a, dim) == dim - 1;
}
```

$$L = \{ a^p b^m a^m b^q \mid m, p, m, q > 0 \wedge m > m \wedge p > q \}$$

$\Sigma = \{a, b\}$

def. una grammatica di genere  $L$  dire se è regolare o libero e giustificare le risposte (dare la dimostrazione)

- AST o G.R. se il lang. è regolare
- Pumping lemma se non è regolare

$$S \rightarrow a S b \mid a S \mid a A b$$

$$A \rightarrow b a a \mid b A a \mid A a$$

$L \text{ reg} \Rightarrow$  Proprietà

~~$L \text{ non reg} \Rightarrow$  no Proprietà~~

$n$  divisibile per 10  $\Rightarrow$   $n$  pari

~~$n$  non div. per 10  $\Rightarrow$   $n$  dispari~~

$n$  non è pari  $\Rightarrow$   $n$  non è div. per 10

$$A \Rightarrow B \equiv \neg B \Rightarrow \neg A$$

$$\left( \forall m \in \mathbb{N}. \left( \exists w \in L. |w| \geq m \wedge \right. \right.$$

$$\left. \left( \forall x, y, z. w = xyz \wedge |xy| \leq m \wedge y \neq \varepsilon \right. \right. \\ \left. \left. \Rightarrow \left( \exists i \in \mathbb{N}. xy^i z \notin L \right) \right) \right)$$

$\Rightarrow$  L non è regolare

$$L = \left\{ a^p b^s a^t b^k \mid p, s, t, k > 0 \wedge p > k \wedge t > s \right\}$$

Qualunque  $n$  a  $n > 0$  prendo la stringa  $w = a^{n+1} b a a b^n \in L$

$$|w| \geq n$$

$$a^n b a a b^n \notin L$$

$$\forall x, y, z. w = xyz \wedge |xy| \leq n \wedge y \neq \varepsilon$$

$$\begin{array}{c} a^{n+1} b a a b^n \\ \underbrace{\hspace{1.5cm}}_{xy} \quad z \end{array}$$

$$\forall x, y, z.$$

$$\left( \begin{array}{l} x = a^s \quad \emptyset \leq s < n \\ y = a^t \quad \emptyset < t \leq n - s \\ z = a^k b a a b^n \quad k = n + 1 - t - s \end{array} \right)$$



$\Rightarrow \left( \mathbb{N} \cdot x y^i z \notin L \right)$

$$xy^0z = a^s a^k b a a b^m \notin L$$

perché  $s+k < m+1$  dato che  $t > 0$

Si dice quel  $\bar{\epsilon}$  il linguaggio generato dalle seq. grammatiche su  $\Sigma = \{a, b, c\}$

lunedì 23 ottobre 2017 10:40

$$S \rightarrow \underline{aSb} \mid aANAb$$

$$A \rightarrow \underline{aA} \mid a \quad \uparrow \uparrow \uparrow$$

$$N \rightarrow \underline{bNa} \mid ba$$

$$L = \left\{ \underbrace{a^m a^p}_{\text{---}} \underbrace{b^m}_{\text{---}} \underbrace{a^m a^s}_{\text{---}} \underbrace{b^m}_{\text{---}} \mid \begin{matrix} m, p, s > 0 \\ m, \end{matrix} \right.$$

$$L = \left\{ a^k b^m a^s b^m \mid \begin{matrix} k \geq 2 \wedge m > 0 \\ s \geq 2 \wedge m > 0 \end{matrix} \right.$$

non regolare  
(dim precedente)

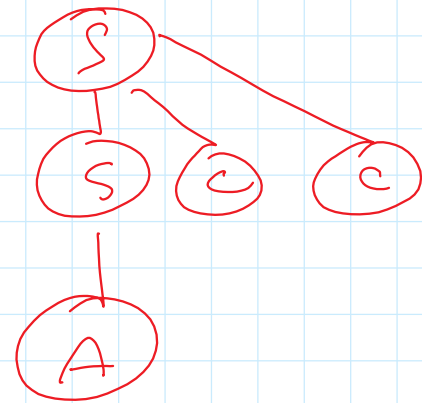
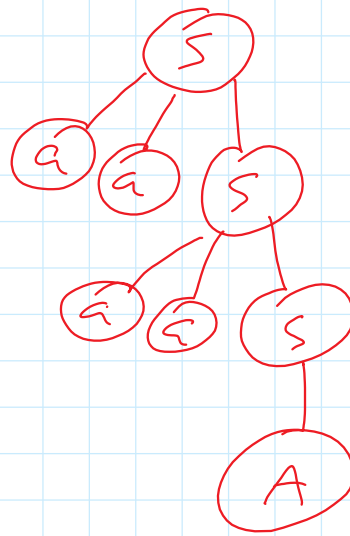
$$\left. \begin{matrix} k > m \wedge \\ s > m \end{matrix} \right\}$$

$$S \rightarrow \underline{aaS} \mid \underline{ScC} \mid \underline{A}$$

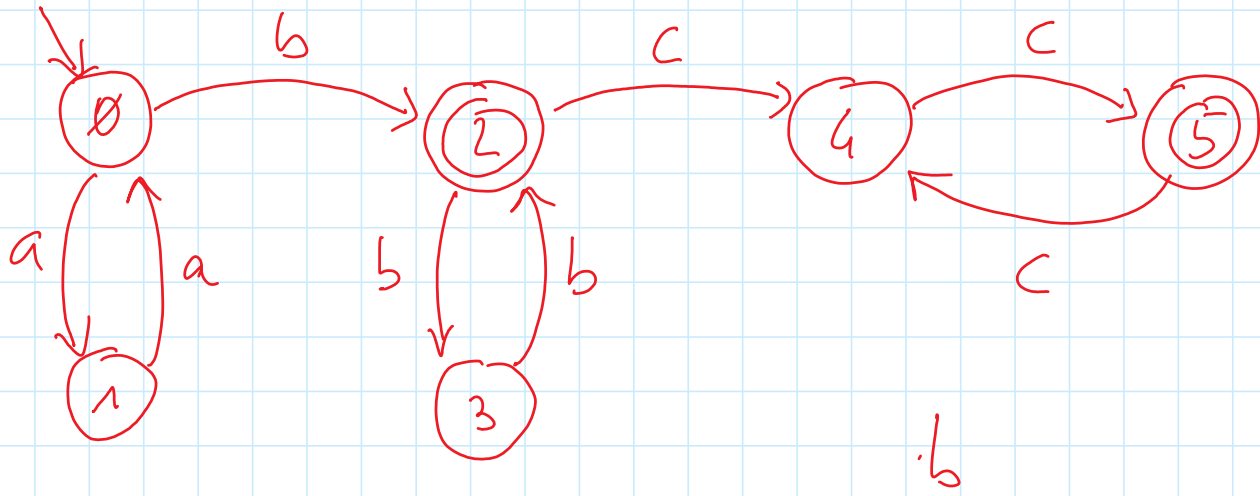
$$A \rightarrow \underline{AbA} \mid \underline{b}$$

$$\mathcal{L} = \{a, b, c\}$$

$$L = \left\{ a^{2m} b^{2k+1} c^{2m} \mid m \geq 0 \wedge m \geq 0 \wedge k \geq 0 \right\}$$



$$L = \left\{ a^{2m} b^{2k+1} c^{2m} \mid m, k \geq 0 \right\}$$



$$\begin{aligned} S &\rightarrow aSb \mid aAb \\ A &\rightarrow \underline{aaA} \mid aBb \\ B &\rightarrow \underline{Bbb} \mid c \end{aligned}$$

$$\begin{aligned} S &\rightarrow aSb \mid aAb \\ A &\rightarrow aA \mid aBb \\ B &\rightarrow Bb \mid c \end{aligned}$$

$$L = \left\{ a^m c b^m \mid m, m \geq 2 \wedge m \% 2 = m \% 2 \right\}$$

↑  
risult delle div. per 2